

<i>Z, len</i>	integer
<i>f</i>	float
<i>n</i>	int
<i>ident, id</i>	identifier
<i>label, l</i>	label
<i>dcls</i>	global variable declarations
<i>fndefns</i>	function declarations
<i>opt_tid</i>	optional thread id
<i>ef_sig</i>	external signature
<i>p</i>	pointer
<i>typ</i>	
<i>fundef, fd</i>	
<i>fn_body</i>	
<i>fn</i>	
<i>ge</i>	
$\delta$	

<i>signedness</i>	<pre> ::=   Signed   Unsigned </pre>	Signedness
<i>intsize</i>	<pre> ::=   I8   I16   I32 </pre>	Integer sizes
<i>floatsize</i>	<pre> ::=   F32   F64 </pre>	Float sizes
<i>type, ty</i>	<pre> ::=   void   int (<i>intsize</i>, <i>signedness</i>)   float (<i>floatsize</i>)   pointer (<i>ty</i>)   array (<i>ty</i>, <i>len</i>)   function (<i>ty*</i>, <i>ty</i>)   struct (<i>id</i>, <i>φ</i>)   union (<i>id</i>, <i>φ</i>)   comp_pointer(<i>id</i>)   (<i>ty</i>) </pre>	Types the void type integer types floating-point types pointer types ( <i>*ty</i> ) array types ( <i>ty[<i>len</i>]</i> ) function types struct types union types pointer to named struct or union S
<i>typelist, ty*</i>	<pre> ::=   nil   <i>ty</i>:::<i>ty*</i> </pre>	Type list
<i>fieldlist, φ</i>	<pre> ::=   nil   (<i>id</i>, <i>ty</i>)::<i>φ</i> </pre>	Field lists
<i>unary_operation, op<sub>1</sub></i>	<pre> ::=   !   ~   - </pre>	unary Boolean negation Integer complement opposite
<i>binary_operation, op<sub>2</sub></i>	<pre> ::=   +   -   *   /   %   &amp;       ^ </pre>	binary addition subtraction multiplication division modulo bitwise and bitwise or bitwise xor

		<<	left shift
		>>	right shift
		==	equality
		!=	not equal
		<	less than
		>	greater than
		<=	less than equal
		>=	greater than equal
<i>expr, e</i>	::=		typed expression
		<i>a<sup>ty</sup></i>	expression
<i>expr_descr, a</i>	::=		basic expressions
		<i>n</i>	integer literal
		<i>f</i>	float literal
		<i>id</i>	variable
		<i>*e</i>	unary pointer dereference
		<i>&amp;e</i>	address-of
		<i>op<sub>1</sub> e</i>	unary operation
		<i>e<sub>1</sub> op<sub>2</sub> e<sub>2</sub></i>	binary operation
		<i>(ty) e</i>	type cast
		<i>e<sub>1</sub>?e<sub>2</sub>:e<sub>3</sub></i>	conditional
		<i>e<sub>1</sub>&amp;&amp;e<sub>2</sub></i>	sequential and
		<i>e<sub>1</sub>  e<sub>2</sub></i>	sequential or
		<i>sizeof (ty)</i>	size of a type
		<i>e.id</i>	access to a member of a struct or union
<i>opt_lhs</i>	::=		optional lhs expression
		<i>(id:ty)=</i>	
<i>opt_e</i>	::=		optional expression
		<i>e</i>	
<i>e*</i>	::=		expression list
<i>atomic_statement, astmt</i>	::=		atomic
		CAS	compare and swap
		ATOMIC_INC	locked inc
<i>statement, s</i>	::=		statements
		skip	do nothing
		<i>e<sub>1</sub>=e<sub>2</sub></i>	assignment [lvalue = rvalue]
		<i>opt_lhs e' (e*)</i>	function or procedure call
		<i>s<sub>1</sub>; s<sub>2</sub></i>	sequence
		<i>if (e<sub>1</sub>) then s<sub>1</sub> else s<sub>2</sub></i>	conditional

		<code>while (e) do s</code>	while
		<code>do s while (e)</code>	do while
		<code>for (s<sub>1</sub>; e<sub>2</sub>; s<sub>3</sub>) s</code>	for loop
		<code>break</code>	break
		<code>continue</code>	continue
		<code>return opt_e</code>	return
		<code>switch (e) ls</code>	switch
		<code>l:s</code>	labelled statement
		<code>goto l</code>	goto
		<code>thread_create(e<sub>1</sub>, e<sub>2</sub>)</code>	thread creation
		<code>opt_lhs astmt(e*)</code>	atomic operation
		<code>mfence</code>	mfence
<i>labeled_statements, ls</i>	::=		labeled statements
		<code>default : s</code>	default
		<code>case n : s ; ls</code>	labeled case
<i>arglist</i>	::=		Argument lists
		<code>ty id</code>	
		<code>ty id, arglist</code>	
<i>varlist</i>	::=		Local variable lists
		<code>ty id ; varlist</code>	
<i>fndefn_internal</i>	::=		function definition
		<code>ty id (arglist) { varlist s }</code>	
<i>program</i>	::=		programs
		<code>dcls fndefns main = id</code>	
<i>val, v</i>	::=		untyped values
		<code>n</code>	integer value
		<code>f</code>	floating point value
		<code>p</code>	pointer
		<code>undef</code>	undef
<i>extval, evl</i>	::=		external values
		<code>extint n</code>	external integer value
		<code>extfloat f</code>	external floating point value
<i>vs</i>	::=		value list
		<code>nil</code>	
		<code>v :: vs</code>	
		<code>vs@[v]</code>	

<i>arg_list, args</i>	::=   nil   <i>id<sup>ty</sup>::args</i>	argument lists
<i>evs</i>	::=   nil   <i>evl::evs</i>	eventval list
<i>memory_chunk, c</i>	::=   Mint32	
<i>mobject_kind</i>	::=   MObjStack	
<i>rmw_instr, rmwi</i>	::=   ADD <i>v</i>   CAS <i>v v'</i>   SET <i>v</i>	
<i>mem_event, me</i>	::=   write <i>p memory_chunk v</i>   read <i>p memory_chunk v</i>   alloc <i>p n mobject_kind</i>   free <i>p mobject_kind</i>   rmw <i>p memory_chunk v rmwi</i>   fence	
<i>event, ev</i>	::=   call <i>id evs</i>   return <i>typ evl</i>   exit <i>n</i>   fail	
<i>thread_event, te</i>	::=   ext <i>event</i>   mem <i>mem_event</i>     exit   start <i>p vs</i>	thread events externally observable event memory event thread-local event normal exit thread start (bootstrap)
<i>opt_pty</i>	::=	optional pointer/type pair
<i>opt_v</i>	::=	optional value
<i>ps</i>	::=	pointer list
$\rho, \rho', \rho''$	::=	environment

$cont, \kappa_S$	$::=$   <code>stop</code>   <code>[- ; <math>s_2</math>]</code> · $\kappa_S$   <code>[while (<math>e</math>) do <math>s</math>]</code> · $\kappa_S$   <code>[do <math>s</math> while (<math>e</math>)]</code> · $\kappa_S$   <code>[for( ; <math>e_2</math>; <math>\diamond s_3</math>) <math>s</math>]</code> · $\kappa_S$   <code>[for( ; <math>\diamond e_2</math>; <math>s_3</math>) <math>s</math>]</code> · $\kappa_S$   <code>[<math>opt\_lhs</math> <math>fd</math>(<math>\_</math>)  <math>\rho</math>]</code> · $\kappa_S$   <code>[switch <math>\kappa_S</math>]</code>   <code>[free <math>ps</math>; return <math>opt\_v</math>]</code> · $\kappa_S$	statement continuation  sequence while do while for loop, pending increment for loop, pending condition evaluation call awaiting args switch protecting break
$expr\_cont, \kappa_e$	$::=$   <code>[<math>op_1^{ty}</math> <math>\_</math>]</code> · $\kappa_e$   <code>[- <math>op_2^{ty_1 * ty_2 \rightarrow ty}</math> <math>e_2</math>]</code> · $\kappa_e$   <code>[<math>v</math> <math>op_2^{ty_1 * ty_2 \rightarrow ty}</math> <math>\_</math>]</code> · $\kappa_e$   <code>[(<math>ty</math>) <math>\_</math> <math>ty'</math>]</code> · $\kappa_e$   <code>[- <math>ty</math> ? <math>e_2</math> : <math>e_3</math>]</code> · $\kappa_e$   <code>[- <math>\cdot</math> <math>\delta</math>]</code> · $\kappa_e$   <code>[* <math>ty</math>]</code> · $\kappa_e$   <code>[- <math>ty = e_2</math>]</code> · $\kappa_S$   <code>[<math>v</math> <math>ty = \_</math>]</code> · $\kappa_S$   <code>[<math>opt\_lhs</math> <math>\_</math> <math>ty</math> (<math>e^*</math>)]</code> · $\kappa_S$   <code>[<math>opt\_lhs</math> <math>v</math> <math>ty</math> (<math>vs, e^*</math>)]</code> · $\kappa_S$   <code>[<math>opt\_lhs</math> <math>astmt</math> (<math>vs, e^*</math>)]</code> · $\kappa_S$   <code>[if (<math>\_</math> <math>ty</math>) then <math>s_1</math> else <math>s_2</math>]</code> · $\kappa_S$   <code>[while (<math>\_e</math>) do <math>s</math>]</code> · $\kappa_S$   <code>[do <math>s</math> while (<math>\_e</math>)]</code> · $\kappa_S$   <code>[for ( ; <math>\_e_2</math>; <math>s_3</math>) <math>s</math>]</code> · $\kappa_S$   <code>[return <math>\_</math>]</code> · $\kappa_S$   <code>[switch (<math>\_</math>) <math>ls</math>]</code> · $\kappa_S$   <code>[thread_create(<math>\_, e_2</math>)]</code> · $\kappa_S$   <code>[thread_create(<math>p, \_</math>)]</code> · $\kappa_S$	expression continuations  unary operation binary operation binary operation   access to member of struct load assignment assignment call function call args atomic args if while dowhile for loop, pending test function return switch thread creation thread creation
$state, \sigma$	$::=$   <code>lval (<math>e</math>)</code> · $\kappa_e$   $\rho$   <code><math>e</math></code> · $\kappa_e$   $\rho$   <code><math>v</math></code> · $\kappa_e$   $\rho$   <code><math>s</math></code> · $\kappa_S$   $\rho$   <code><math>vs</math></code> · $\kappa_S$   <code>bind (<math>fn, vs, args</math>)</code> · $\kappa_S$   $\rho$   <code>alloc (<math>vs, args</math>)</code> · $\kappa_S$   $\rho$   <code><math>opt\_lhs</math> ext (<math>\_^{typ}</math>)</code> · $\kappa_S$   $\rho$   <code><math>opt\_lhs</math> <math>v</math></code> · $\kappa_S$   $\rho$	states

$\sigma \xrightarrow{te} \sigma'$     Labelled Transitions (parameterised over  $ge$ )

$$\frac{}{n^{ty} \cdot \kappa_e | \rho \longrightarrow n \cdot \kappa_e | \rho} \text{ STEPCONSTINT}$$

$$\frac{}{f^{ty} \cdot \kappa_e | \rho \longrightarrow f \cdot \kappa_e | \rho} \text{ STEPCONSTFLOAT}$$

$\frac{id^{ty} \cdot \kappa_e  _\rho \longrightarrow \text{lval}(id^{ty}) \cdot [*_{-ty}] \cdot \kappa_e  _\rho}{}$	STEPVAREXPRBYVALUE
$\frac{\rho!id = \text{Some } p}{\text{lval}(id^{ty}) \cdot \kappa_e  _\rho \longrightarrow p \cdot \kappa_e  _\rho}$	STEPVARLOCAL
$\frac{\rho!id = \text{None} \quad \text{Genv.find\_symbol } ge \ id = \text{Some } p}{\text{lval}(id^{ty}) \cdot \kappa_e  _\rho \longrightarrow p \cdot \kappa_e  _\rho}$	STEPVARGLOBAL
$\frac{\text{access\_mode } ty' = \text{By\_value } c \quad ty' = \text{type\_of\_chunk } c \quad \text{Val.has\_type } v \ ty'}{p \cdot [*_{-ty'}] \cdot \kappa_e  _\rho \xrightarrow{\text{mem(read } p \ c \ v)} v \cdot \kappa_e  _\rho}$	STEPLOADBYVALUE
$\frac{\text{access\_mode } ty' = \text{By\_reference} \ \backslash / \ \text{access\_mode } ty' = \text{By\_nothing}}{p \cdot [*_{-ty'}] \cdot \kappa_e  _\rho \longrightarrow p \cdot \kappa_e  _\rho}$	STEPLOADNOTBYVALUE
$\frac{}{\&e^{ty} \cdot \kappa_e  _\rho \longrightarrow \text{lval}(e) \cdot \kappa_e  _\rho}$	STEPADDR
$\frac{}{e_1?e_2:e_3^{ty} \cdot \kappa_e  _\rho \longrightarrow e_1 \cdot [-^{typeof \ e_1?e_2:e_3}] \cdot \kappa_e  _\rho}$	STEPCONDITION
$\frac{\text{is\_true } v \ ty}{v \cdot [-^{ty?e_2:e_3}] \cdot \kappa_e  _\rho \longrightarrow e_2 \cdot \kappa_e  _\rho}$	STEPCONDITIONTRUE
$\frac{\text{is\_false } v \ ty}{v \cdot [-^{ty?e_2:e_3}] \cdot \kappa_e  _\rho \longrightarrow e_3 \cdot \kappa_e  _\rho}$	STEPCONDITIONFALSE
$\frac{}{*e^{ty} \cdot \kappa_e  _\rho \longrightarrow e \cdot [*_{-ty}] \cdot \kappa_e  _\rho}$	STEPDEREF
$\frac{}{\text{lval}(*e^{ty}) \cdot \kappa_e  _\rho \longrightarrow e \cdot \kappa_e  _\rho}$	STEPDEREFLVAL
$\frac{}{e.id^{ty} \cdot \kappa_e  _\rho \longrightarrow \text{lval}(e.id^{ty}) \cdot [*_{-ty}] \cdot \kappa_e  _\rho}$	STEPFIELD
$\frac{\text{typeof } e = \text{struct}(id', \phi) \quad \text{field\_offset } id \ \phi = \text{OK } \delta}{\text{lval}(e.id^{ty}) \cdot \kappa_e  _\rho \longrightarrow \text{lval}(e) \cdot [- \cdot \delta] \cdot \kappa_e  _\rho}$	STEPFSTRUCT1
$\frac{p' = \text{Ptr.add } p \ (\text{Int.repr } \delta)}{p \cdot [- \cdot \delta] \cdot \kappa_e  _\rho \longrightarrow p' \cdot \kappa_e  _\rho}$	STEPFSTRUCT2
$\frac{\text{typeof } e = \text{union}(id', \phi)}{\text{lval}(e.id^{ty}) \cdot \kappa_e  _\rho \longrightarrow \text{lval}(e) \cdot \kappa_e  _\rho}$	STEPFUNION
$\frac{v = \text{Vint}(\text{Int.repr}(\text{sizeof } ty'))}{\text{sizeof}(ty')^{ty} \cdot \kappa_e  _\rho \longrightarrow v \cdot \kappa_e  _\rho}$	STEPSIZEOF
$\frac{}{op_1 e^{ty} \cdot \kappa_e  _\rho \longrightarrow e \cdot [op_1^{typeof \ e} \_ ] \cdot \kappa_e  _\rho}$	STEPUNOP1
$\frac{\text{sem\_unary\_operation } op_1 \ v \ ty = \text{Some } v'}{v \cdot [op_1^{ty} \_ ] \cdot \kappa_e  _\rho \longrightarrow v' \cdot \kappa_e  _\rho}$	STEPUNOP
$\frac{}{(e_1 \ op_2 \ e_2)^{ty} \cdot \kappa_e  _\rho \longrightarrow e_1 \cdot [- \ op_2^{typeof \ e_1 * \text{typeof } e_2 \rightarrow ty} \ e_2] \cdot \kappa_e  _\rho}$	STEPBINOP1
$\frac{\text{valid\_arg } op_2 \ ty_1 \ ty_2 \ v = \text{true}}{v \cdot [- \ op_2^{ty_1 * ty_2 \rightarrow ty} \ e_2] \cdot \kappa_e  _\rho \longrightarrow e_2 \cdot [v \ op_2^{ty_1 * ty_2 \rightarrow ty} \_ ] \cdot \kappa_e  _\rho}$	STEPBINOP2

$$\begin{array}{c}
\frac{\text{sem\_binary\_operation } op_2 \ v_1 \ ty_1 \ v_2 \ ty_2 = \text{Some } v}{v_2 \cdot [v_1 \ op_2^{ty_1 * ty_2 \rightarrow ty} \ ] \cdot \kappa_e \mid \rho \longrightarrow v \cdot \kappa_e \mid \rho} \text{STEPBINOP} \\
\frac{}{(ty) \ e^{ty'} \cdot \kappa_e \mid \rho \longrightarrow e \cdot [(ty) \ \text{typeof } e] \cdot \kappa_e \mid \rho} \text{STEPCAST1} \\
\frac{\text{cast } v \ ty' \ ty \ v'}{v \cdot [(ty) \ \_ {ty'}] \cdot \kappa_e \mid \rho \longrightarrow v' \cdot \kappa_e \mid \rho} \text{STEPCAST2} \\
\frac{n_0 = \text{Int.repr } 0 \\ n_1 = \text{Int.repr } 1}{e_1 \ \&\& \ e_2^{ty} \cdot \kappa_e \mid \rho \longrightarrow e_1 ? (e_2 ? (n_1 \ ty) : (n_0 \ ty)^{ty}) : n_0 \ ty^{ty} \cdot \kappa_e \mid \rho} \text{STEPANDBOOL} \\
\frac{n_0 = \text{Int.repr } 0 \\ n_1 = \text{Int.repr } 1}{e_1 \mid \mid \ e_2^{ty} \cdot \kappa_e \mid \rho \longrightarrow e_1 ? (n_1 \ ty) : (e_2 ? (n_1 \ ty) : (n_0 \ ty)^{ty})^{ty} \cdot \kappa_e \mid \rho} \text{STEPORBOOL} \\
\frac{}{\text{thread\_create}(e_1, e_2) \cdot \kappa_s \mid \rho \longrightarrow e_1 \cdot [\text{thread\_create}(\_, e_2)] \cdot \kappa_s \mid \rho} \text{STEPTHREAD} \\
\frac{}{p \cdot [\text{thread\_create}(\_, e_2)] \cdot \kappa_s \mid \rho \longrightarrow e_2 \cdot [\text{thread\_create}(p, \_)] \cdot \kappa_s \mid \rho} \text{STEPTHREADFN} \\
\frac{}{v \cdot [\text{thread\_create}(p, \_)] \cdot \kappa_s \mid \rho \xrightarrow{\text{start } p \ v : \text{nil}} \text{skip} \cdot \kappa_s \mid \rho} \text{STEPTHREADEVT} \\
\frac{}{e_1 = e_2 \cdot \kappa_s \mid \rho \longrightarrow \text{lval}(e_1) \cdot [\_ \text{typeof } e_1 = e_2] \cdot \kappa_s \mid \rho} \text{STEPASSIGN1} \\
\frac{}{v_1 \cdot [\_ = e_2] \cdot \kappa_s \mid \rho \longrightarrow e_2 \cdot [v_1 \ ty = \_] \cdot \kappa_s \mid \rho} \text{STEPASSIGN2} \\
\frac{\text{type\_to\_chunk } ty_1 = \text{Some } c \\ \text{cast\_value\_to\_chunk } c \ v_1 = v_2}{v_1 \cdot [p_1 \ ty_1 = \_] \cdot \kappa_s \mid \rho \xrightarrow{\text{mem}(\text{write } p_1 \ c \ v_2)} \text{skip} \cdot \kappa_s \mid \rho} \text{STEPASSIGN} \\
\frac{}{s_1 ; s_2 \cdot \kappa_s \mid \rho \longrightarrow s_1 \cdot [\_ ; s_2] \cdot \kappa_s \mid \rho} \text{STEPSEQ} \\
\frac{}{\text{opt\_lhs } e' (e^*) \cdot \kappa_s \mid \rho \longrightarrow e' \cdot [\text{opt\_lhs } \_ \text{typeof } e' (e^*)] \cdot \kappa_s \mid \rho} \text{STEPCALL} \\
\frac{\text{Genv.find\_funct } ge \ v = \text{Some } fd \\ \text{type\_of\_fundef } fd = ty}{v \cdot [\text{opt\_lhs } \_ \ ty (\text{nil})] \cdot \kappa_s \mid \rho \longrightarrow \text{nil} \cdot [\text{opt\_lhs } fd(\_) \mid \rho] \cdot \kappa_s} \text{STEPCALLARGSNONE} \\
\frac{}{v \cdot [\text{opt\_lhs } \_ \ ty (e :: e^*)] \cdot \kappa_s \mid \rho \longrightarrow e \cdot [\text{opt\_lhs } v^{ty}(\text{nil}, e^*)] \cdot \kappa_s \mid \rho} \text{STEPCALLARGS1} \\
\frac{}{v_1 \cdot [\text{opt\_lhs } v^{ty}(vs, e :: e^*)] \cdot \kappa_s \mid \rho \longrightarrow e \cdot [\text{opt\_lhs } v^{ty}(vs@[v_1], e^*)] \cdot \kappa_s \mid \rho} \text{STEPCALLARGS2} \\
\frac{\text{Genv.find\_funct } ge \ v = \text{Some } fd \\ \text{type\_of\_fundef } fd = ty}{v' \cdot [\text{opt\_lhs } v^{ty}(vs, \text{nil})] \cdot \kappa_s \mid \rho \longrightarrow vs@[v'] \cdot [\text{opt\_lhs } fd(\_) \mid \rho] \cdot \kappa_s} \text{STEPCALLFINISH} \\
\frac{}{\text{opt\_lhs } astmt(e :: e^*) \cdot \kappa_s \mid \rho \longrightarrow e \cdot [\text{opt\_lhs } astmt(\text{nil}, e^*)] \cdot \kappa_s \mid \rho} \text{STEPATOMIC} \\
\frac{}{v \cdot [\text{opt\_lhs } astmt(vs, e :: e^*)] \cdot \kappa_s \mid \rho \longrightarrow e \cdot [\text{opt\_lhs } astmt(vs@[v], e^*)] \cdot \kappa_s \mid \rho} \text{STEPATOMICARGS} \\
\frac{\text{sem\_atomic\_statement } astmt (vs ++ v :: \text{nil}) = \text{Some} (p, rmwi) \\ \text{Val.has\_type } v' (type\_of\_chunk \text{Mint32})}{v \cdot [astmt(vs, \text{nil})] \cdot \kappa_s \mid \rho \xrightarrow{\text{mem}(\text{rmw } p \ \text{Mint32 } v' \ \text{rmwi})} \text{skip} \cdot \kappa_s \mid \rho} \text{STEPATOMICFINISHNONE}
\end{array}$$



$$\begin{array}{c}
\text{sem\_atomic\_statement } \mathit{astmt} \ (vs \ ++ \ v \ :: \ \text{nil}) = \text{Some} \ (p, \ \mathit{rmwi}) \\
\text{Val.has\_type } v' \ (\text{type\_of\_chunk } \text{Mint32}) \\
\hline
v \cdot [(id:ty) = \mathit{astmt}(vs, \text{nil})] \cdot \kappa_s \mid \rho \xrightarrow{\text{mem}(\text{rmw } p \text{Mint32 } v' \ \mathit{rmwi})} (id:ty) = v' \cdot \kappa_s \mid \rho \quad \text{STEPATOMICFINISHSOME} \\
\\
\frac{}{\text{mfence} \cdot \kappa_s \mid \rho \xrightarrow{\text{mem fence}} \text{skip} \cdot \kappa_s \mid \rho} \quad \text{STEPFENCE} \\
\\
\frac{}{\text{continue} \cdot [-; s] \cdot \kappa_s \mid \rho \longrightarrow \text{continue} \cdot \kappa_s \mid \rho} \quad \text{STEPCONTINUE} \\
\\
\frac{}{\text{break} \cdot [-; s] \cdot \kappa_s \mid \rho \longrightarrow \text{break} \cdot \kappa_s \mid \rho} \quad \text{STEPBREAK} \\
\\
\frac{}{\text{if} \ (e) \ \text{then } s_1 \ \text{else } s_2 \cdot \kappa_s \mid \rho \longrightarrow e \cdot [\text{if} \ (\_ \text{typeof } e) \ \text{then } s_1 \ \text{else } s_2] \cdot \kappa_s \mid \rho} \quad \text{STEPIFTHENELSE} \\
\\
\frac{\text{is\_true } v \ ty}{v \cdot [\text{if} \ (\_ \text{ty}) \ \text{then } s_1 \ \text{else } s_2] \cdot \kappa_s \mid \rho \longrightarrow s_1 \cdot \kappa_s \mid \rho} \quad \text{STEPIFTHENELSETRUE} \\
\\
\frac{\text{is\_false } v \ ty}{v \cdot [\text{if} \ (\_ \text{ty}) \ \text{then } s_1 \ \text{else } s_2] \cdot \kappa_s \mid \rho \longrightarrow s_2 \cdot \kappa_s \mid \rho} \quad \text{STEPIFTHENELSEFALSE} \\
\\
\frac{}{\text{while} \ (e) \ \text{do } s \cdot \kappa_s \mid \rho \longrightarrow e \cdot [\text{while} \ (\_ e) \ \text{do } s] \cdot \kappa_s \mid \rho} \quad \text{STEPWHILE} \\
\\
\frac{\text{is\_true } v \ (\text{typeof } e)}{v \cdot [\text{while} \ (\_ e) \ \text{do } s] \cdot \kappa_s \mid \rho \longrightarrow s \cdot [\text{while} \ (e) \ \text{do } s] \cdot \kappa_s \mid \rho} \quad \text{STEPWHILETRUE} \\
\\
\frac{\text{is\_false } v \ (\text{typeof } e)}{v \cdot [\text{while} \ (\_ e) \ \text{do } s] \cdot \kappa_s \mid \rho \longrightarrow \text{skip} \cdot \kappa_s \mid \rho} \quad \text{STEPWHILEFALSE} \\
\\
\frac{}{\text{continue} \cdot [\text{while} \ (e) \ \text{do } s] \cdot \kappa_s \mid \rho \longrightarrow \text{while} \ (e) \ \text{do } s \cdot \kappa_s \mid \rho} \quad \text{STEPCONTINUEWHILE} \\
\\
\frac{}{\text{break} \cdot [\text{while} \ (e) \ \text{do } s] \cdot \kappa_s \mid \rho \longrightarrow \text{skip} \cdot \kappa_s \mid \rho} \quad \text{STEPBREAKWHILE} \\
\\
\frac{}{\text{do } s \ \text{while} \ (e) \cdot \kappa_s \mid \rho \longrightarrow s \cdot [\text{do } s \ \text{while} \ (e)] \cdot \kappa_s \mid \rho} \quad \text{STEPDOWHILE} \\
\\
\frac{\text{is\_true } v \ (\text{typeof } e)}{v \cdot [\text{do } s \ \text{while} \ (\_ e)] \cdot \kappa_s \mid \rho \longrightarrow \text{do } s \ \text{while} \ (e) \cdot \kappa_s \mid \rho} \quad \text{STEPDOWHILETRUE} \\
\\
\frac{\text{is\_false } v \ (\text{typeof } e)}{v \cdot [\text{do } s \ \text{while} \ (\_ e)] \cdot \kappa_s \mid \rho \longrightarrow \text{skip} \cdot \kappa_s \mid \rho} \quad \text{STEPDOWHILEFALSE} \\
\\
\frac{}{\text{continue} \cdot [\text{do } s \ \text{while} \ (e)] \cdot \kappa_s \mid \rho \longrightarrow e \cdot [\text{do } s \ \text{while} \ (\_ e)] \cdot \kappa_s \mid \rho} \quad \text{STEPDOCONTINUEWHILE} \\
\\
\frac{}{\text{break} \cdot [\text{do } s \ \text{while} \ (e)] \cdot \kappa_s \mid \rho \longrightarrow \text{skip} \cdot \kappa_s \mid \rho} \quad \text{STEPDOBREAKWHILE} \\
\\
\frac{}{\text{for} \ (s_1; e_2; s_3) s \cdot \kappa_s \mid \rho \longrightarrow s_1 \cdot [\text{for} \ (; \diamond e_2; s_3) s] \cdot \kappa_s \mid \rho} \quad \text{STEPFORINIT} \\
\\
\frac{}{\text{skip} \cdot [\text{for} \ (; \diamond e_2; s_3) s] \cdot \kappa_s \mid \rho \longrightarrow e_2 \cdot [\text{for} \ (; \_ e_2; s_3) s] \cdot \kappa_s \mid \rho} \quad \text{STEPFORCOND} \\
\\
\frac{\text{is\_true } v \ (\text{typeof } e_2)}{v \cdot [\text{for} \ (; \_ e_2; s_3) s] \cdot \kappa_s \mid \rho \longrightarrow s \cdot [\text{for} \ (; e_2; \diamond s_3) s] \cdot \kappa_s \mid \rho} \quad \text{STEPFORTRUE} \\
\\
\frac{\text{is\_false } v \ (\text{typeof } e_2)}{v \cdot [\text{for} \ (; \_ e_2; s_3) s] \cdot \kappa_s \mid \rho \longrightarrow \text{skip} \cdot \kappa_s \mid \rho} \quad \text{STEPFORFALSE} \\
\\
\frac{}{\text{skip} \cdot [\text{for} \ (; e_2; \diamond s_3) s] \cdot \kappa_s \mid \rho \longrightarrow s_3 \cdot [\text{for} \ (; \diamond e_2; s_3) s] \cdot \kappa_s \mid \rho} \quad \text{STEPFORINCR}
\end{array}$$

$$\begin{array}{c}
\frac{}{\text{break} \cdot [\text{for}(\ ; e_2; \diamond s_3) s] \cdot \kappa_s |_\rho \longrightarrow \text{skip} \cdot \kappa_s |_\rho} \text{STEPFORBREAK} \\
\\
\frac{}{\text{continue} \cdot [\text{for}(\ ; e_2; \diamond s_3) s] \cdot \kappa_s |_\rho \longrightarrow s_3 \cdot [\text{for}(\ ; \diamond e_2; s_3) s] \cdot \kappa_s |_\rho} \text{STEPFORCONTINUE} \\
\\
\frac{\text{call\_cont } \kappa_s = (\text{Kcall None (Internal } fn) \rho'' \kappa'_s) \\ \text{fn} \cdot (\text{fn\_return}) = \text{Tvoid} \\ \text{ps} = \text{sorted\_pointers\_of\_env } \rho'}{\text{return} \cdot \kappa_s |_{\rho'} \longrightarrow \text{skip} \cdot [\text{free ps; return None}] \cdot \kappa_s |_{\rho'}} \text{STEPRETURNNONE} \\
\\
\frac{}{\text{skip} \cdot [\text{free } p :: ps; \text{return } opt\_v] \cdot \kappa_s |_\rho \xrightarrow{\text{mem (free } p \text{ MObjStack)}} \text{skip} \cdot [\text{free ps; return } opt\_v] \cdot \kappa_s |_\rho} \text{STEPRETURNMEM} \\
\\
\frac{\text{call\_cont } \kappa_s = \kappa'_s \\ \text{get\_fundef } \kappa'_s = \text{Some (Internal } fn) \\ \text{fn} \cdot (\text{fn\_return}) \langle \rangle \text{Tvoid}}{\text{return } e \cdot \kappa_s |_{\rho'} \longrightarrow e \cdot [\text{return } \_ ] \cdot \kappa_s |_{\rho'}} \text{STEPRETURN SOME} \\
\\
\frac{\text{ps} = \text{sorted\_pointers\_of\_env } \rho}{v \cdot [\text{return } \_ ] \cdot \kappa_s |_\rho \longrightarrow \text{skip} \cdot [\text{free ps; return (Some } v)] \cdot \kappa_s |_\rho} \text{STEPRETURN SOME1} \\
\\
\frac{}{\text{switch } (e) ls \cdot \kappa_s |_\rho \longrightarrow e \cdot [\text{switch } (\_) ls] \cdot \kappa_s |_\rho} \text{STEP SWITCH} \\
\\
\frac{s = \text{seq\_of\_labeled\_statement (select\_switch } n \text{ } ls)}{n \cdot [\text{switch } (\_) ls] \cdot \kappa_s |_\rho \longrightarrow s \cdot [\text{switch } \kappa_s] |_\rho} \text{STEP SELECT SWITCH} \\
\\
\frac{}{\text{break} \cdot [\text{switch } \kappa_s] |_\rho \longrightarrow \text{skip} \cdot \kappa_s |_\rho} \text{STEP BREAK SWITCH} \\
\\
\frac{}{\text{continue} \cdot [\text{switch } \kappa_s] |_\rho \longrightarrow \text{continue} \cdot \kappa_s |_\rho} \text{STEP CONTINUE SWITCH} \\
\\
\frac{}{l : s \cdot \kappa_s |_\rho \longrightarrow s \cdot \kappa_s |_\rho} \text{STEP LABEL} \\
\\
\frac{\text{call\_cont } \kappa_s = \kappa'_s \\ \text{get\_fundef } \kappa'_s = (\text{Some (Internal } fn)) \\ \text{find\_label } l \text{ fn} \cdot (\text{fn\_body}) \kappa'_s = \text{Some } (s', \kappa''_s)}{\text{goto } l \cdot \kappa_s |_\rho \longrightarrow s' \cdot \kappa''_s |_\rho} \text{STEP GOTO} \\
\\
\frac{\text{args} = \text{fn} \cdot (\text{fn\_params}) ++ \text{fn} \cdot (\text{fn\_vars}) \\ \text{fd} = \text{Internal } fn}{vs \cdot [\text{opt\_lhs } fd(\_) |_\rho] \cdot \kappa_s \longrightarrow \text{alloc } (vs, \text{args}) \cdot [\text{opt\_lhs } fd(\_) |_\rho] \cdot \kappa_s |_{\rho_{\text{empty}}}} \text{STEP FUNCTION INTERNAL} \\
\\
\frac{n = \text{Int.repr}(\text{sizeof } ty)}{\text{alloc } (vs, id^{ty} :: args) \cdot \kappa_s |_\rho \xrightarrow{\text{mem (alloc } p \ n \ \text{MObjStack)}} \text{alloc } (vs, \text{args}) \cdot \kappa_s |_{\rho \oplus (id \mapsto p)}} \text{STEP ALLOC LOCAL} \\
\\
\frac{\text{args} = \text{fn} \cdot (\text{fn\_params}) \\ \text{fd} = (\text{Internal } fn)}{\text{alloc } (vs, \text{nil}) \cdot [\text{opt\_lhs } fd(\_) |_{\rho'}] \cdot \kappa_s |_{\rho'} \longrightarrow \text{bind } (fn, vs, \text{args}) \cdot [\text{opt\_lhs } fd(\_) |_{\rho'}] \cdot \kappa_s |_{\rho'}} \text{STEP BIND ARGS} \\
\\
\frac{\rho!id = \text{Some } p \\ \text{type\_to\_chunk } ty = (\text{Some } c) \\ \text{cast\_value\_to\_chunk } c \ v_1 = v_2}{\text{bind } (fn, v_1 :: vs, id^{ty} :: args) \cdot \kappa_s |_\rho \xrightarrow{\text{mem (write } p \ c \ v_2)} \text{bind } (fn, vs, \text{args}) \cdot \kappa_s |_\rho} \text{STEP BIND ARGS} \\
\\
\frac{s = \text{fn} \cdot (\text{fn\_body})}{\text{bind } (fn, \text{nil}, \text{nil}) \cdot \kappa_s |_\rho \longrightarrow s \cdot \kappa_s |_\rho} \text{STEP TRANSFER FUN}
\end{array}$$

$\text{true} \ (* \ \text{event\_match} \ (\text{external\_function} \ \text{id} \ \text{targs} \ \text{ty}) \ \text{vs} \ \text{t} \ \text{vres} \ \rightarrow \ *)$   
 $\text{fd} = \text{External} \ \text{id} \ \text{ty}^* \ \text{ty}$   
 $\text{vs} = \text{map} \ \text{val\_of\_eval} \ \text{evs}$

---

STEPEXTERNALCALL

$$\text{vs} \cdot [\text{opt\_lhs} \ \text{fd}(\_) \mid_{\rho}] \cdot \kappa_s \xrightarrow{\text{ext}(\text{call} \ \text{id} \ \text{evs})} \text{opt\_lhs} \ \text{ext}(\_ \text{^TYP}) \cdot \kappa_s \mid_{\rho}$$

$\text{Val.has\_type} \ v \ \text{typ}$   
 $\text{fd} = \text{External} \ \text{id} \ \text{ty}^* \ \text{ty}$   
 $\text{typ} = \text{match} \ (\text{opttyp\_of\_type} \ \text{ty}) \ \text{with} \ | \ \text{Some} \ x \ \Rightarrow \ x \ | \ \text{None} \ \Rightarrow \ \text{Ast.Tint} \ \text{end}$   
 $v = \text{val\_of\_eval} \ \text{evl}$

---

STEPEXTERNALRET

$$\text{opt\_lhs} \ \text{ext}(\_ \text{^TYP}) \cdot \kappa_s \mid_{\rho} \xrightarrow{\text{ext}(\text{return} \ \text{typ} \ \text{evl})} \text{opt\_lhs} \ v \cdot \kappa_s \mid_{\rho}$$

$\rho! \ \text{id} = \text{Some} \ p$   
 $\text{type\_to\_chunk} \ \text{ty} = \text{Some} \ c$   
 $\text{cast\_value\_to\_chunk} \ c \ v_1 = v_2$

---

STEPEXTERNALSTORESOMELOCAL

$$(\text{id}:\text{ty}) = v_1 \cdot \kappa_s \mid_{\rho} \xrightarrow{\text{mem}(\text{write} \ p \ c \ v_2)} \text{skip} \cdot \kappa_s \mid_{\rho}$$

$\rho! \ \text{id} = \text{None}$   
 $\text{Genv.find\_symbol} \ \text{ge} \ \text{id} = \text{Some} \ p$   
 $\text{type\_to\_chunk} \ \text{ty} = \text{Some} \ c$   
 $\text{cast\_value\_to\_chunk} \ c \ v_1 = v_2$

---

STEPEXTERNALSTORESOMEGLOBAL

$$(\text{id}:\text{ty}) = v_1 \cdot \kappa_s \mid_{\rho} \xrightarrow{\text{mem}(\text{write} \ p \ c \ v_2)} \text{skip} \cdot \kappa_s \mid_{\rho}$$

STEPEXTERNALSTORENONE

$$v \cdot \kappa_s \mid_{\rho} \longrightarrow \text{skip} \cdot \kappa_s \mid_{\rho}$$

STEPSKIP

$$\text{skip} \cdot [\_ ; s_2] \cdot \kappa_s \mid_{\rho} \longrightarrow s_2 \cdot \kappa_s \mid_{\rho}$$

STEPWHILELOOP

$$\text{skip} \cdot [\text{while} \ (e) \ \text{do} \ s] \cdot \kappa_s \mid_{\rho} \longrightarrow \text{while} \ (e) \ \text{do} \ s \cdot \kappa_s \mid_{\rho}$$

STEPDOWHILELOOP

$$\text{skip} \cdot [\text{do} \ s \ \text{while} \ (e)] \cdot \kappa_s \mid_{\rho} \longrightarrow e \cdot [\text{do} \ s \ \text{while} \ (-e)] \cdot \kappa_s \mid_{\rho}$$

STEPSKIPSWITCH

$$\text{skip} \cdot [\text{switch} \ \kappa_s] \mid_{\rho} \longrightarrow \text{skip} \cdot \kappa_s \mid_{\rho}$$

STEPRETURNNONEFINISH

$$\text{call\_cont} \ \kappa_s = [\text{fd}(\_) \mid_{\rho'}] \cdot \kappa'_s$$


---


$$\text{skip} \cdot [\text{free} \ \text{nil}; \ \text{return} \ \text{opt\_v}] \cdot \kappa_s \mid_{\rho''} \longrightarrow \text{skip} \cdot \kappa'_s \mid_{\rho'}$$

$\text{type\_to\_chunk} \ \text{ty} = (\text{Some} \ c)$   
 $\rho! \ \text{id} = \text{Some} \ p$   
 $\text{call\_cont} \ \kappa_s = [(\text{id}:\text{ty}) = \text{fd}(\_) \mid_{\rho'}] \cdot \kappa'_s$   
 $\text{cast\_value\_to\_chunk} \ c \ v_1 = v_2$

---

STEPRETURN SOME FINISH LOCAL

$$\text{skip} \cdot [\text{free} \ \text{nil}; \ \text{return} \ (\text{Some} \ v_1)] \cdot \kappa_s \mid_{\rho''} \xrightarrow{\text{mem}(\text{write} \ p \ c \ v_2)} \text{skip} \cdot \kappa'_s \mid_{\rho'}$$

$\text{type\_to\_chunk} \ \text{ty} = (\text{Some} \ c)$   
 $\rho! \ \text{id} = \text{None}$   
 $\text{Genv.find\_symbol} \ \text{ge} \ \text{id} = \text{Some} \ p$   
 $\text{call\_cont} \ \kappa_s = [(\text{id}:\text{ty}) = \text{fd}(\_) \mid_{\rho'}] \cdot \kappa'_s$   
 $\text{cast\_value\_to\_chunk} \ c \ v_1 = v_2$

---

STEPRETURN SOME FINISH GLOBAL

$$\text{skip} \cdot [\text{free} \ \text{nil}; \ \text{return} \ (\text{Some} \ v_1)] \cdot \kappa_s \mid_{\rho''} \xrightarrow{\text{mem}(\text{write} \ p \ c \ v_2)} \text{skip} \cdot \kappa'_s \mid_{\rho'}$$

STEPSTOP

$$\text{skip} \cdot \text{stop} \mid_{\rho} \xrightarrow{\text{exit}} \text{skip} \cdot \text{stop} \mid_{\rho}$$

Definition rules:            94 good      0 bad  
Definition rule clauses: 178 good      0 bad