

Talk presented at IEEE Security and Privacy 2014

Chip and Skim: Cloning EMV cards with the pre-play attack

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EMV – leading system for payments across the world



Chip and Skim. Bond, Choudary, Murdoch, Skorobogatov, Anderson.

EMV – introduced to remove magstripe counterfeiting

- EMV uses CHIP & PIN
- Should protect against card cloning and abuse
- Should decrease fraud



Chip and **PIN**

EMV is not totally secure in practice

- We discovered 2 important flaws in EMV
 - engineering flaw
 - protocol flaw
- In practice these allow same effect as card cloning
 - we can perform a “CHIP & PIN” transaction without the original EMV card

EMV protocol for POS/ATM



EMV protocol – online authorisation



$D = \{\text{Amount, Country, Date, } \mathbf{UN}, \dots\}$

→

$\text{REQ} = \{\text{UN, ATC, IAD, } \dots\}, \text{ AUTH REQ} = \text{MAC}_{\mathbf{K}}(D, \mathbf{ATC}, \text{IAD})$

←

$\text{RESP} = \{\text{OK/BAD}\}, \text{ AUTH RESP} = \text{MAC}_{\mathbf{K}}(\text{RESP}, \text{AUTH REQ}, \dots)$

→

UN = Unpredictable Number

ATC = Application Transaction Counter

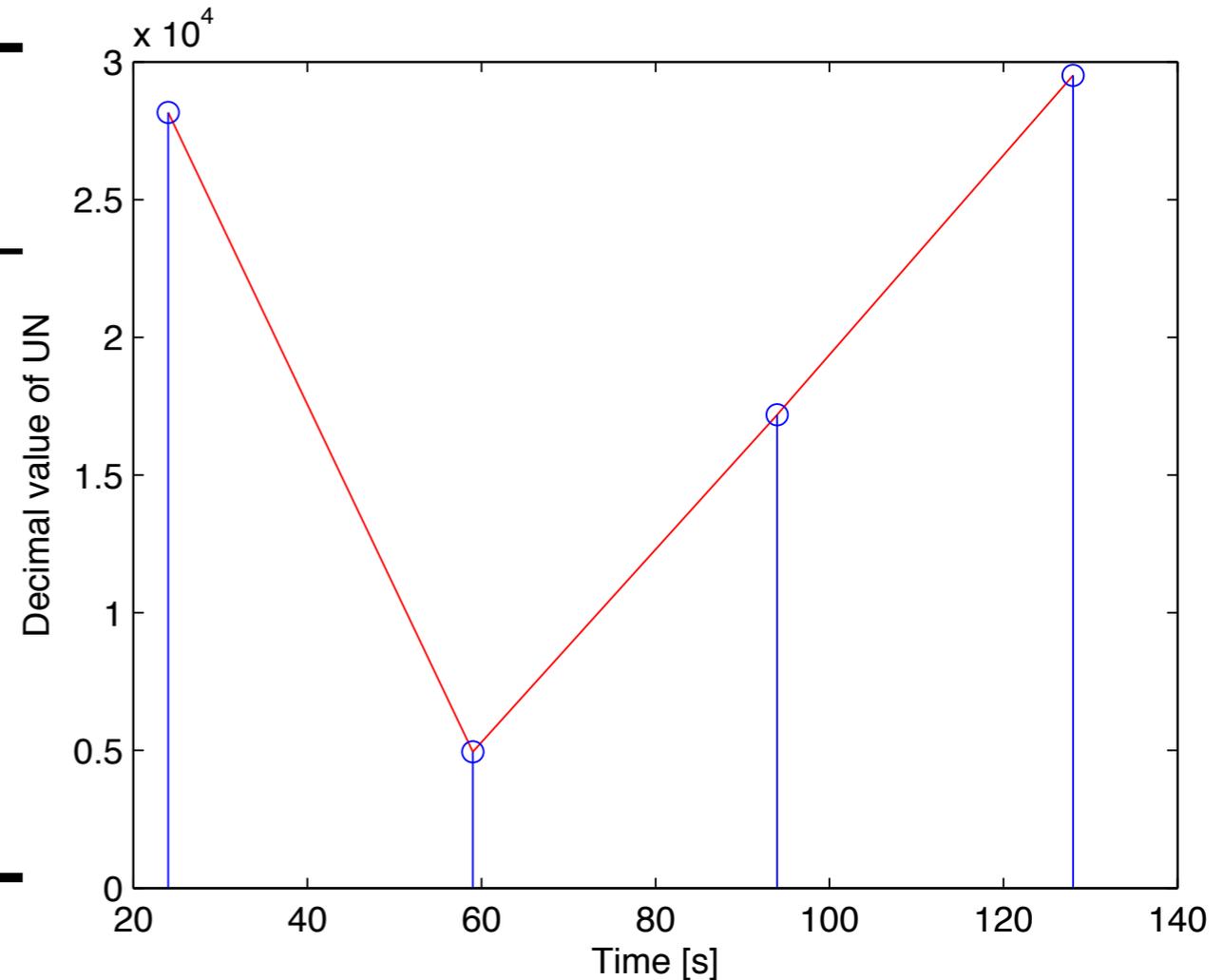
Evidence from real data: UN is a counter!

Time	UN
10:37:24	F1246E04
10:37:59	F1241354
10:38:34	F1244328
10:39:08	F1247348

- 17 bits fixed
- 15 bits seem to follow a linear counter

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No terminal ID



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Pre-play attack: exploit predictable UN

Step 1: Skim PIN & data for set of UNs

$D_1 = \{\text{Amount, Country, Date, } \mathbf{UN}_1, \dots\}$

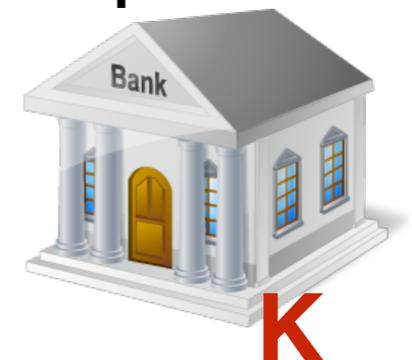
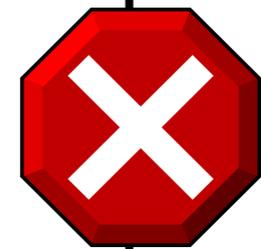
AUTH REQ₁

$D_2 = \{\text{Amount, Country, Date, } \mathbf{UN}_2, \dots\}$

AUTH REQ₂

⋮

ID	UN	AUTH REQ
1	xx	aa
2	yy	bb
...		



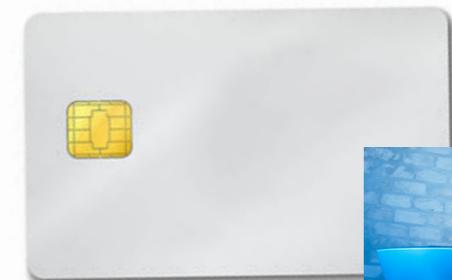
Pre-play attack: exploit predictable UN

Step 2: replay data to get diamond



$D = \{\text{Amount, Country, Date, **UN**, ...}\}$

Replay from table of skimmed data

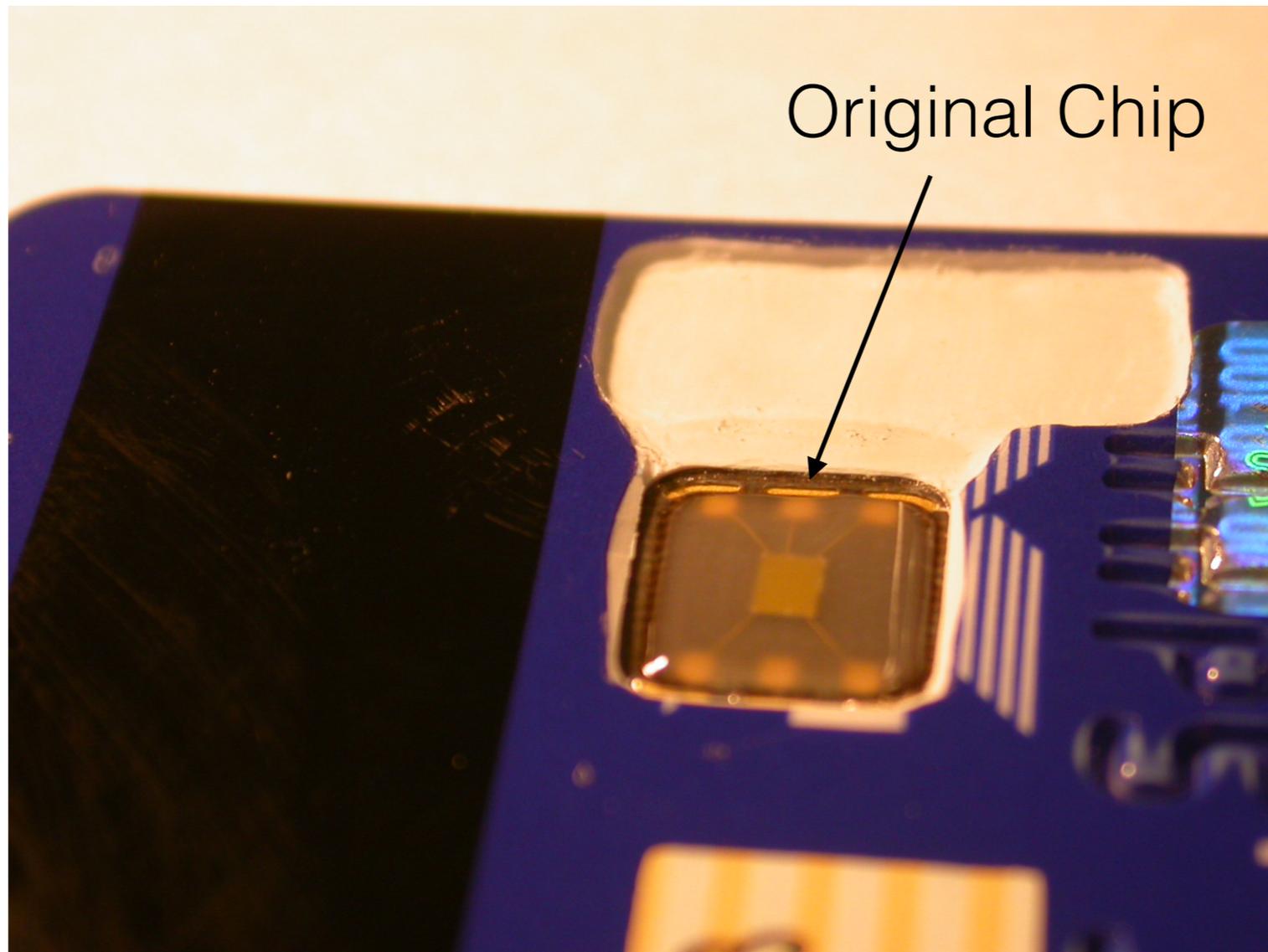


ID	UN	AUTH REQ
1	xx	aa
2	yy	bb
...		

Can we find weak RNGs?

- Previous EMV specs only required 4 consecutive UNs to be different
 - a counter would work better than a secure TRNG
- We decided to find out ...

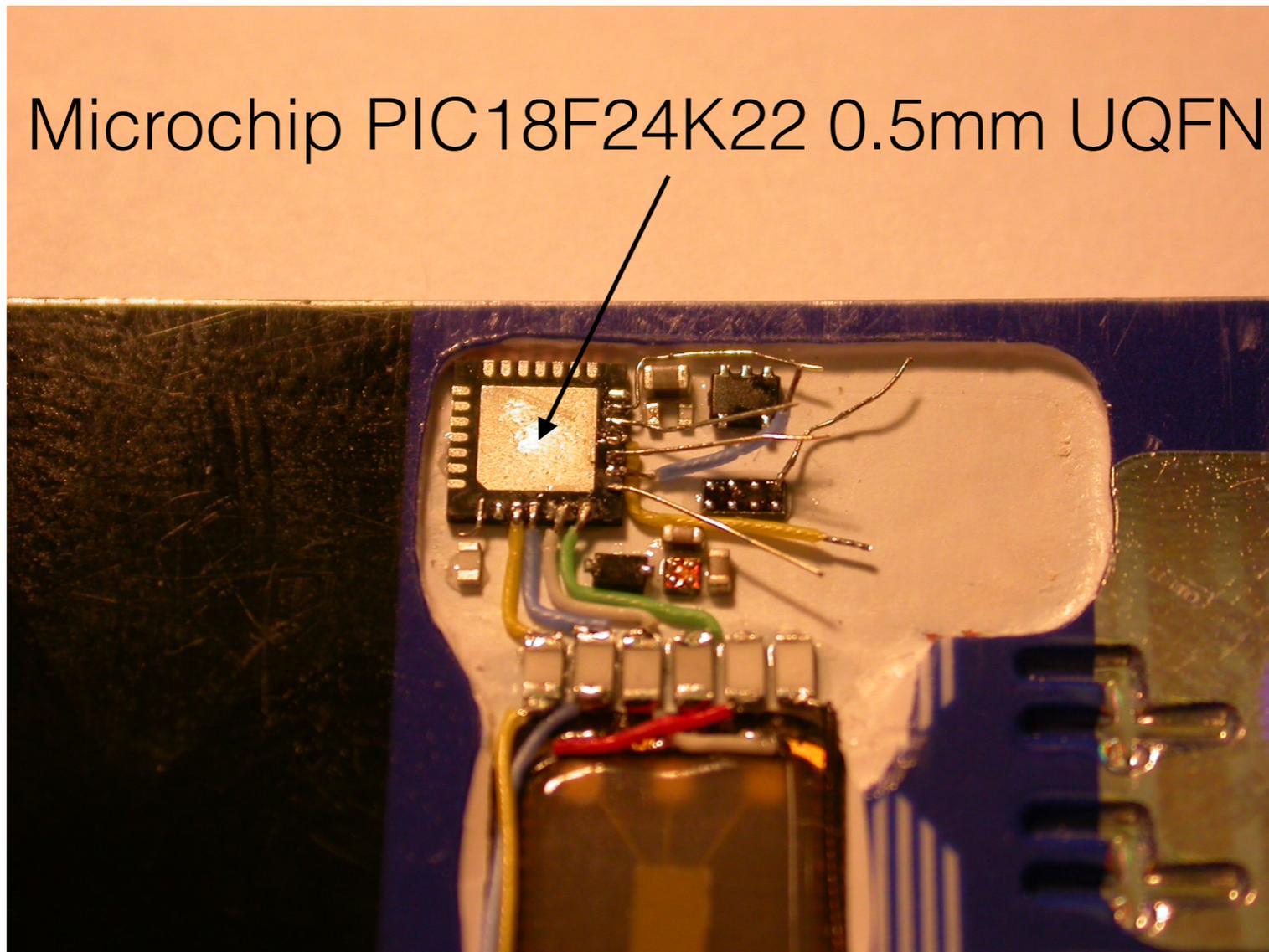
Searching for weak RNG: using ATM logger



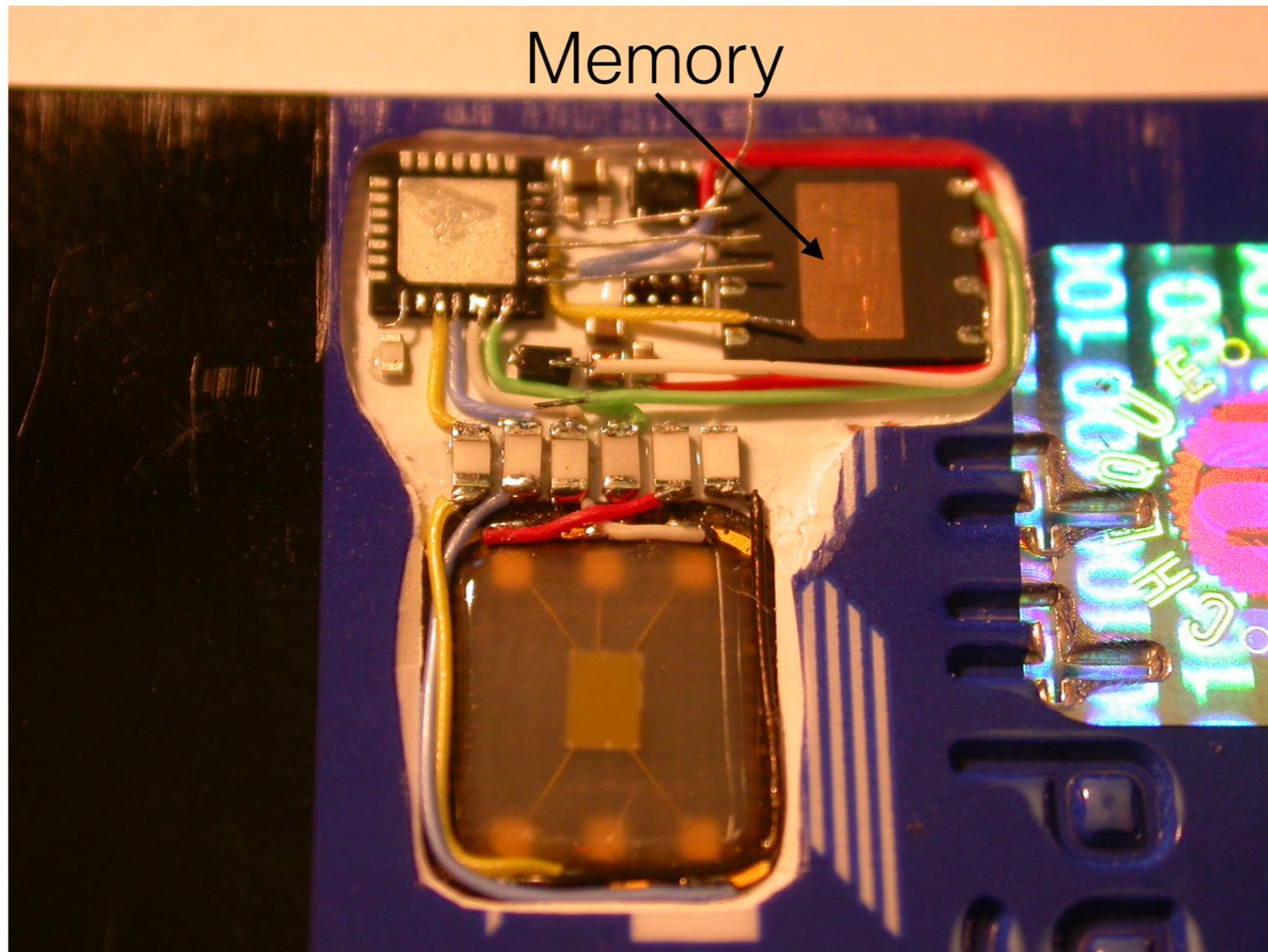
Chip and Skim. Bond, Choudary, Murdoch, Skorobogatov, Anderson.

Searching for weak RNG: using ATM logger

Microchip PIC18F24K22 0.5mm UQFN

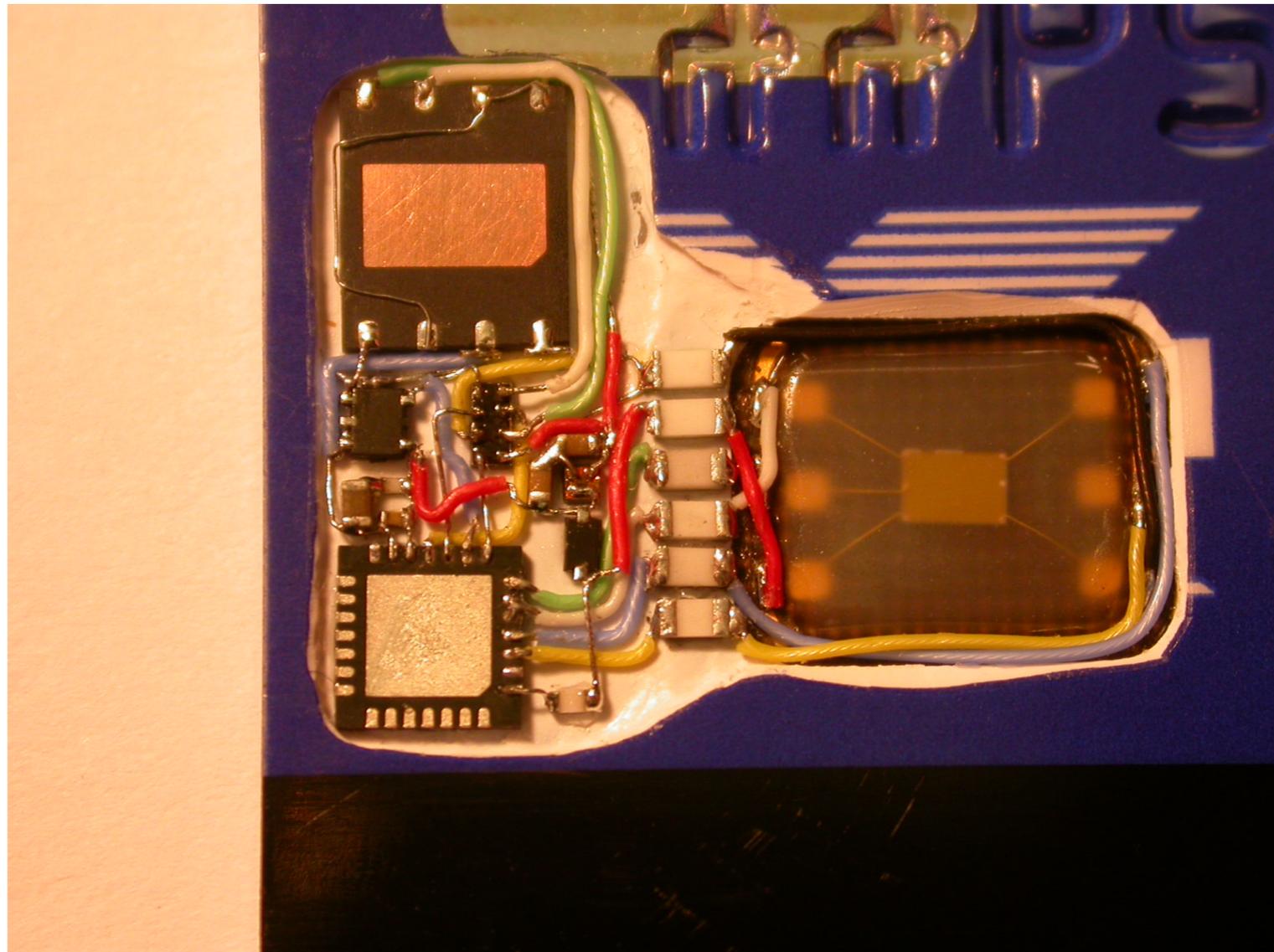


Searching for weak RNG: using ATM logger



Chip and Skim. Bond, Choudary, Murdoch, Skorobogatov, Anderson.

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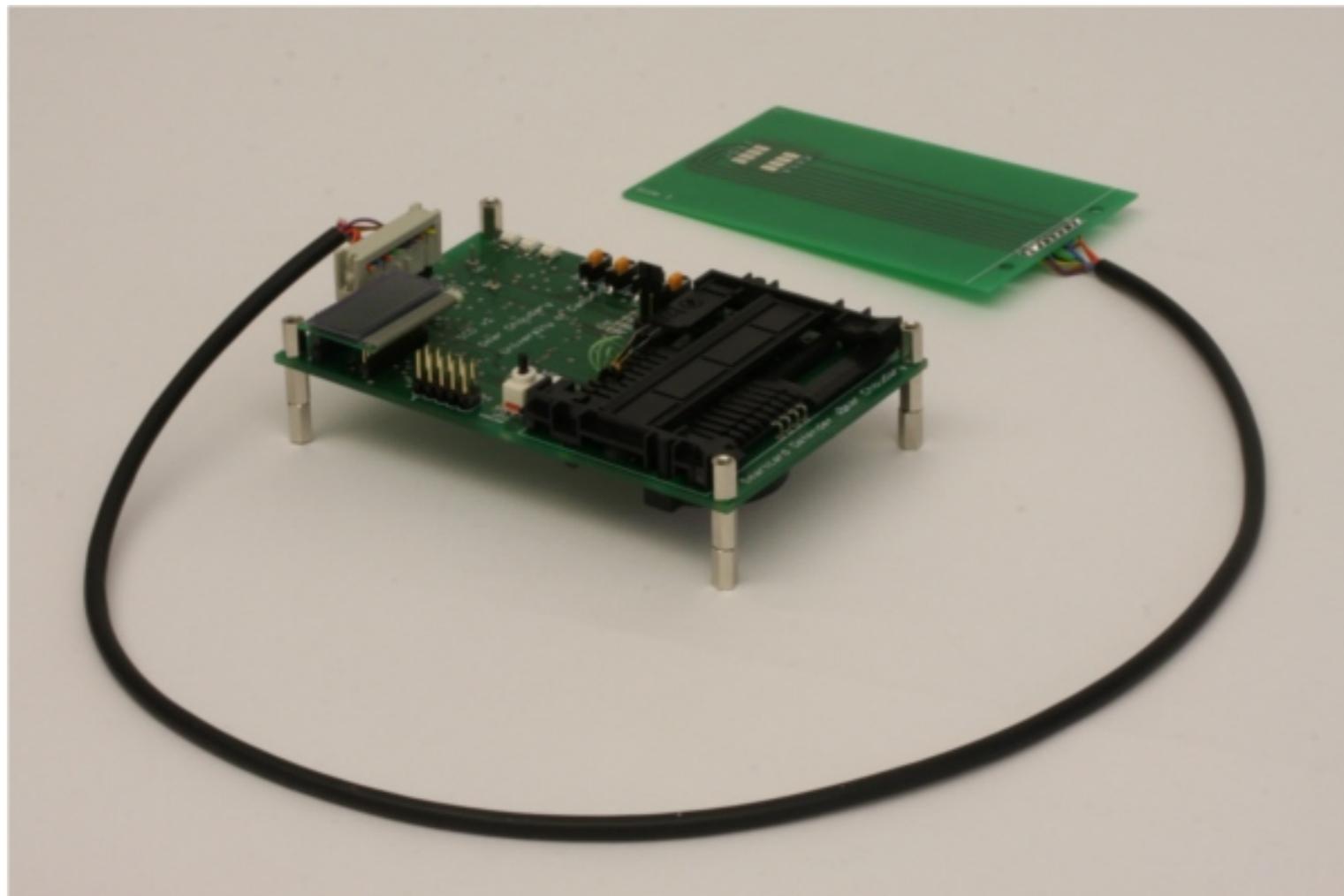


Searching for weak RNG: using ATM logger

- Characteristic C (5 bits fixed):
 - Third nibble is 0
 - First bit is 0
- 11 ATMs had same output
- Possibly due to common lib

Weak RNGs	
ATM1	690d4df2
ATM1	69053549
ATM1	660341c7
ATM1	5e0fc8f2
ATM2	6f0c2d04
ATM2	580fc7d6
ATM2	4906e840
ATM2	46099187

Searching for weak RNG: using SmartCard Detective



Chip and Skim. Bond, Choudary, Murdoch, Skorobogatov, Anderson.

Searching for weak RNG: using SmartCard Detective

- Results from local POS
- First bit still 0, but otherwise could not find clear pattern

Stronger RNGs	
POS1	013A8CE2
POS1	01FB2C16
POS1	2A26982F
POS1	39EB1E19
POS1	293FBA89
POS1	49868033

The deeper problem: We can use our own UN!



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UN generated by Terminal (POS, ATM), not issuer!

The pre-play attack by tampering UN

Step 1: get PIN & data for a **chosen** UN



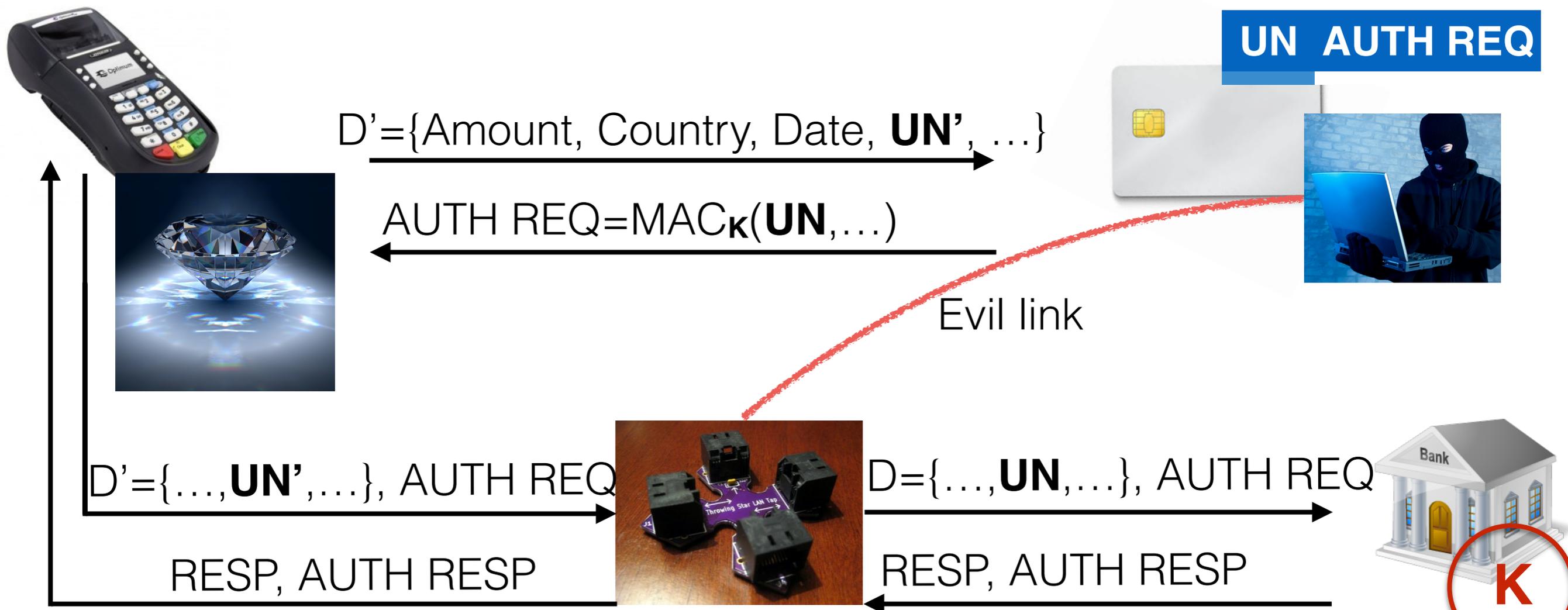
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The pre-play attack by tampering UN

Step 2: replay data & tamper UN to get diamond



Can we actually modify the UN sent by the terminal ?

Likely. It depends on bank, country, regulator, etc.

Issuer

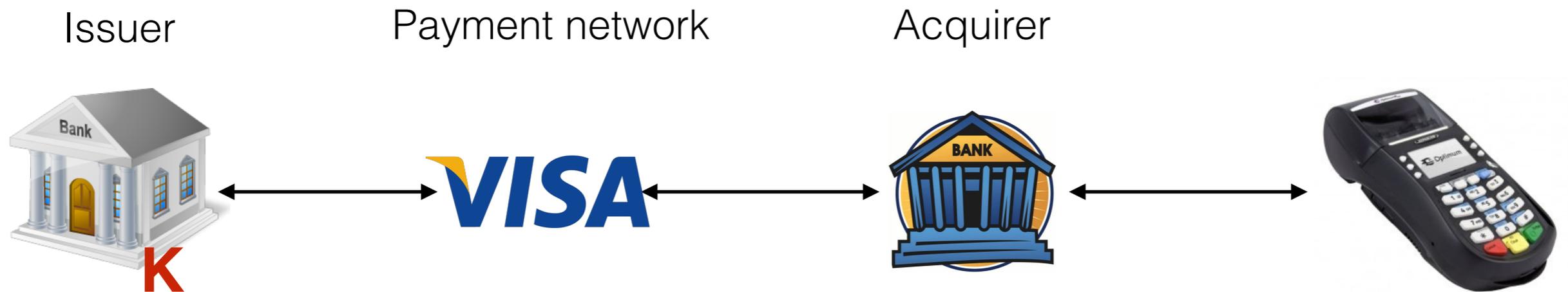


syntax, semantic: ISO 8583, ISO 20022, ...

transport: AS2, AS3, SWIFT, FTP, IFX, ...

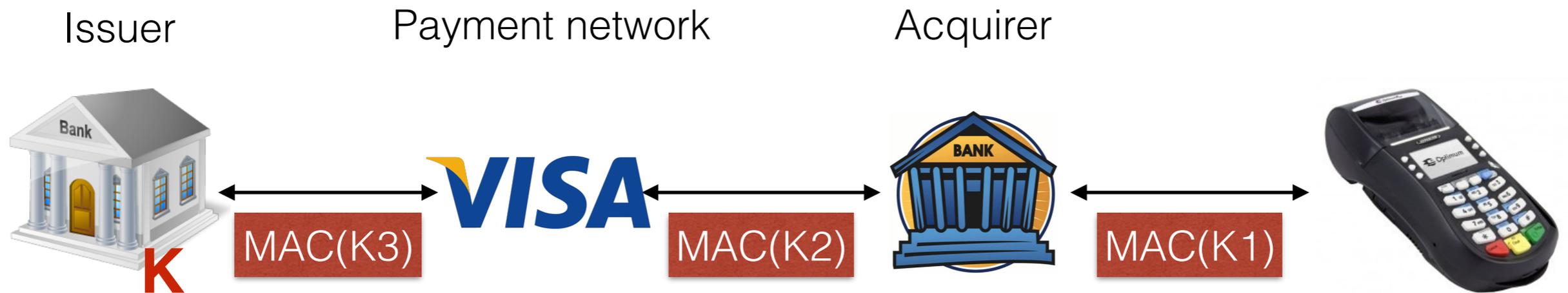
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Can we actually modify the UN sent by the terminal ?

... emergence of new functionality such as authentication methods ...
[VISA "Transactions Acceptance Device Guide" 2013]

Practical example: Maxwell Parsons in UK

- injected data into the bank system (reverse transactions), stealing £2,560,000 in 7 months

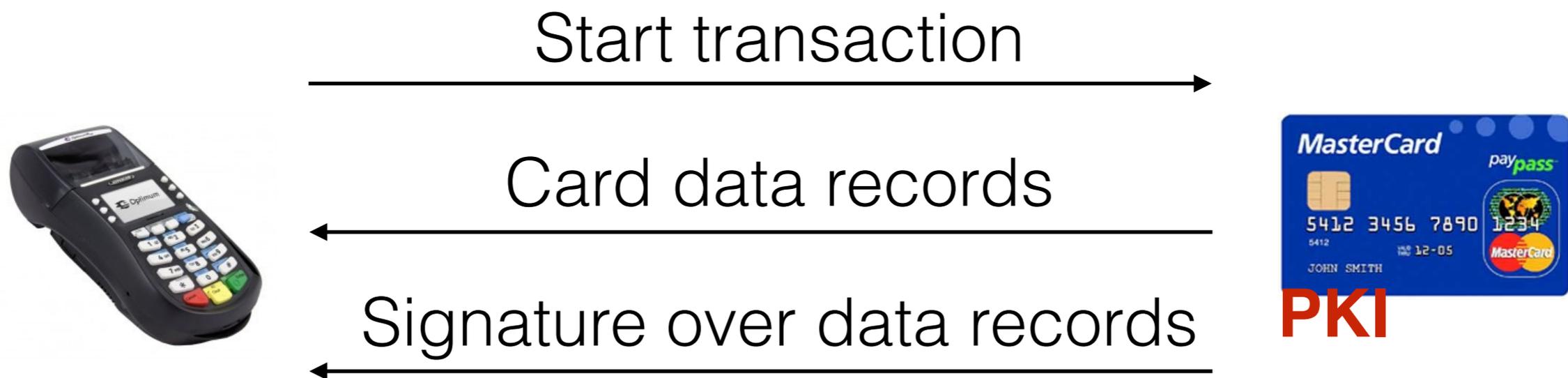
Can we actually modify the UN sent by the terminal ?

- Even if authentication is enabled, there are options:
 - Malware infection of POS/ATM
 - Supply chain attacks (react on covert signal)
 - Collusive or dishonest merchant

It is a protocol problem

- Issuer relies on fresh UN for transaction
- But UN generated by terminal
- Terminal might not have incentive to cooperate

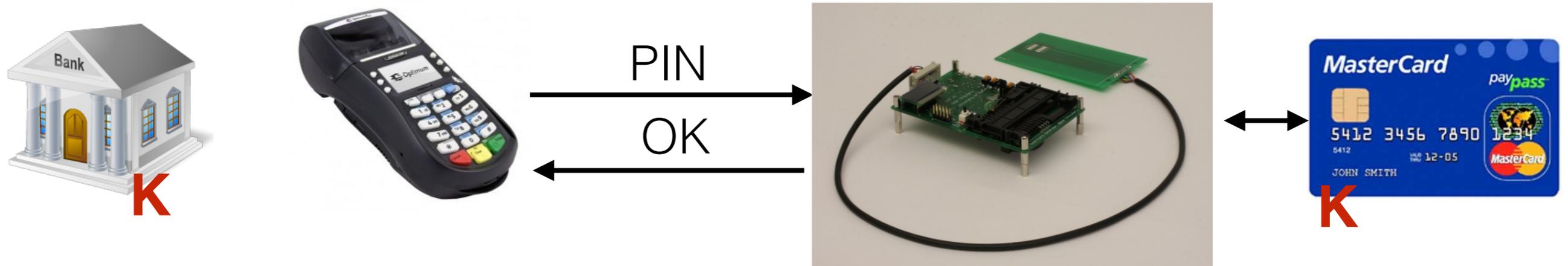
Card authentication via DDA does not help



Same UN for both DDA and ARQC => skim signature as well

PIN verification does not help either

Simply skim PIN during step (1) of attack, or lie [Oakland '10]



Blocking a pre-play attack using the Transaction Certificate (TC)



$D = \{\text{Amount, Country, Date, UN, ...}\}$

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Blocking a pre-play attack using the Transaction Certificate (TC)



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External Authenticate

RESP, AUTH RESP

Final exchange

D'

$\text{TC} = \text{MAC}_K(D', \text{ATC}, \text{IAD})$

Importance of TC not taken into consideration

- Problem 1: TC not routinely kept
 - not needed for clearance, may be discarded
 - only needed to ensure that card does not need to go online (issuer) at next transaction and to provide liability protection to acquirer
- Problem 2: TC may be sent within 24 hours
 - good: send daily TC batches to reduce #messages
 - bad: this leaves system open to pre-play attack

What could EMV do

- Fix RNG everywhere
- Mandatory authentication between all parties
- Request terminal to keep log of UNs for disputes
- Mandatory check or at least storage of TC for every transaction
 - **TC should be the only probative evidence** in case of disputes
- For high-value transactions, check TC before customer leaves the shop!

Conclusions

- We discovered a deep and important flaw in the EMV implementation, indistinguishable from card cloning
- Issuer relies on freshness, but this is generated by another party
 - Changing the protocol is unlikely to happen
 - Practical solution is mandatory use or retention of TC
- Lack of understanding and deliberate overstatement of security may lead to customers being defrauded
- Bank regulators should prohibit EMV liability shift

Questions?

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Industry response

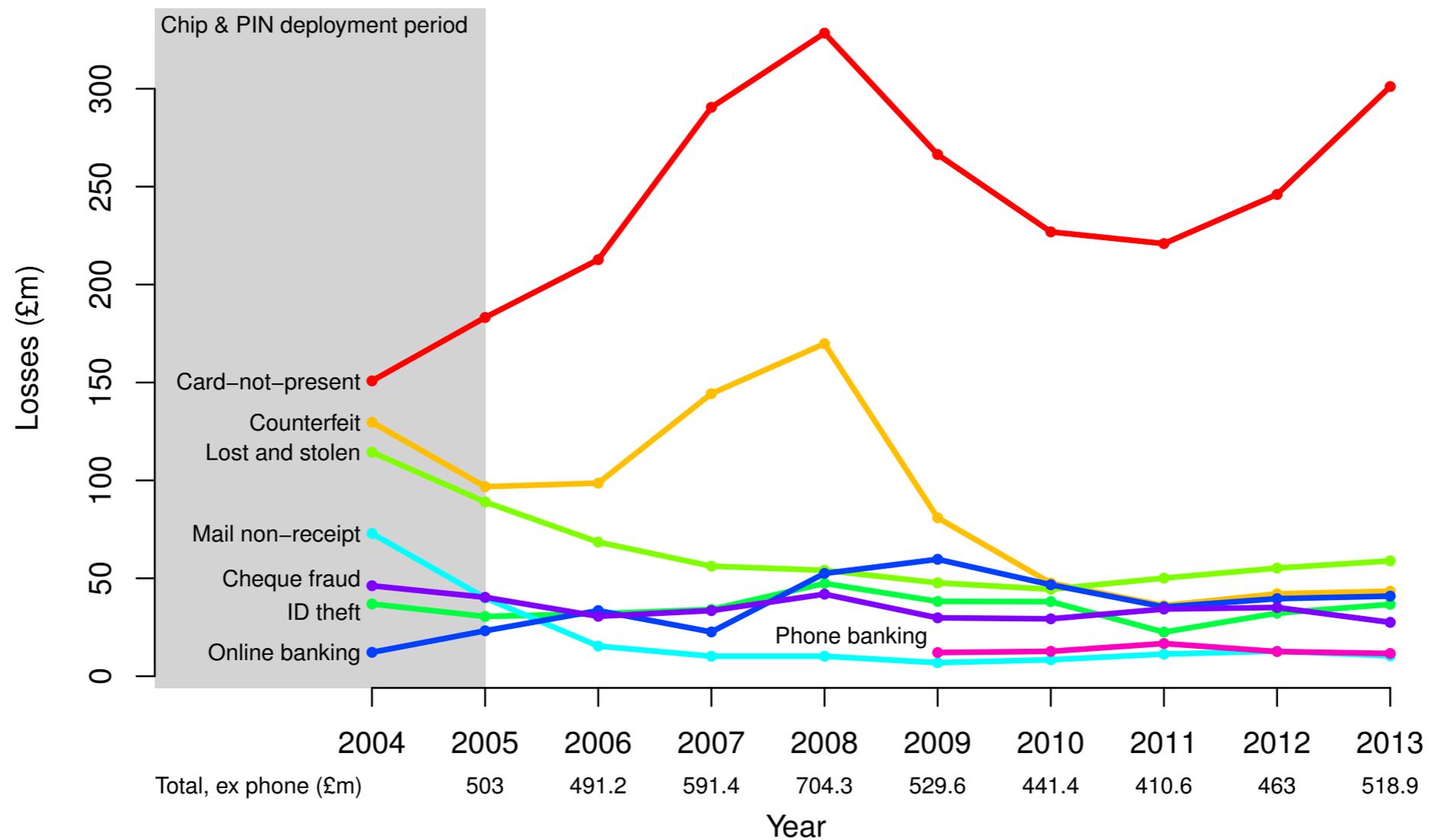
- RNG attack disclosed in early 2012
- Banks and payment switches acknowledge receipt
- April 2012 EMVCo publishes update on RNG
- However, ATMs and terminals still vulnerable to malware
- industry insider mentioned Malta's case may involve ATM malware

ATM reverse engineering



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Bank losses by kind



Fraud levels on UK-issued payments cards

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