Formalising a Number Theory Textbook: Lessons Learnt

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Formalising maths – up to 2015

A machine-checked proof of the odd order theorem, using Coq (Gonthier et al.)

 ... of Gödel's incompleteness theorems, using Isabelle (Paulson)

 ... of the Kepler conjecture, using HOL Light and Isabelle (Hales et al.)

The Lean phenomenon (2017–)

- Sophisticated definitions: schemes, perfectoid spaces
- Big libraries of advanced mathematics (mathlib)
- Liquid tensor experiment: verifying the brand-new work of a Fields medallist

Experiments to confirm "that a proof assistant can handle complexity ..., which is rather different from formalising a long proof about simple objects." — Kevin Buzzard

ALEXANDRIA

(ERC Project GA 742178)

Aim: to support working mathematicians

... by developing tools and libraries

What sorts of mathematics—and proofs—can we formalise?

Using Isabelle/HOL, by the way

Some of our topics

quantum computation • projective geometry • countingroots • Budan–Fourier theorem • algebraically closed fields

ordinal partition theory

Grothendieck schemes

Szemerédi's regularity lemma

Roth: arithmetic progressions

Balog–Szemerédi–Gowers theorem, Khovanskii's theorem ...

ω-categories

Aiming for variety; trying to test the limits



Graduate Texts in Mathematics

- Elliptic and modular functions
- The Dedekind eta function
- Approximation theorems (Kronecker's and others)
- The Riemann zeta function

Lots of advanced material

Tom M. Apostol Modular Functions and Dirichlet Series in Number Theory

Second Edition



7.2 Dirichlet's approximation theorem

Theorem 7.1. Given any real θ and any positive integer N, there exist integers h and k with $0 < k \le N$ such that

$$(1) |k\theta - h| < \frac{1}{N}.$$

PROOF. Let $\{x\} = x - [x]$ denote the fractional part of x. Consider the N + 1 real numbers

$$0, \{\theta\}, \{2\theta\}, \ldots, \{N\theta\}.$$

All these numbers lie in the half open unit interval $0 \le \{m\theta\} < 1$. Now divide the unit interval into N equal half-open subintervals of length 1/N. Then some subinterval must contain at least two of these fractional parts, say $\{a\theta\}$ and $\{b\theta\}$, where $0 \le a < b \le N$. Hence we can write

$$|\{b\theta\} - \{a\theta\}| < \frac{1}{N}.$$

But

$$\{b\theta\} - \{a\theta\} = b\theta - [b\theta] - a\theta + [a\theta] = (b - a)\theta - ([b\theta] - [a\theta]).$$

Therefore if we let

$$k = b - a$$
 and $h = [b\theta] - [a\theta]$

inequality (2) becomes

$$|k\theta - h| < \frac{1}{N}$$
, with $0 < k \le N$.

This proves the theorem.

- ★ Unfortunately, the **simultaneous** version turned out to be necessary: approximate $\theta_1, ..., \theta_n$ as $|k\theta_i h_i| < \frac{1}{N}$ where $k \le N^n$.
- Proof obtained from Hardy and Wright, An Introduction to the Theory of Numbers.
- Still, an elementary proof by the pigeon hole principle, easily formalised (almost on a single slide!)

```
theorem Dirichlet approx simult:
  fixes \vartheta :: "nat \Rightarrow real" and N n :: nat
  assumes "N > 0"
  obtains q p where "0 < q" "q \le int (N^n)" and "\land i < n \implies \circ of int q * \vartheta i - of int(p i) < 1/N"
proof -
  have lessN: "nat |x|^* real N| < N" if "0 \le x" "x < 1" for x
  proof -
    have "|x * real N| < N"
      using that by (simp add: assms floor_less_iff)
    with assms show ?thesis by linarith
  qed
  define interv where "interv \equiv \lambda k. {real k/N..< Suc k/N}"
  define fracs where "fracs \equiv \lambda k. map (\lambda i. frac (real k * \vartheta i)) [0..<n]"
  define X where "X \equiv fracs ` {..N^n}"
  define Y where "Y \equiv set (List.n lists n (map interv [0..<N]))"
  have interv iff: "interv k = interv k' \leftrightarrow k = k'" for k k'
    using assms by (auto simp: interv def Ico eq Ico divide strict right mono)
  have in interv: "x \in interv (nat |x|^* real N|)" if "x \ge 0" for x
    using that assms by (simp add: interv def divide simps) linarith
  have False
    if non: "\forall a b. b \leq N^n \longrightarrow a < b \longrightarrow \neg (\forall i < n. | frac (real b * \vartheta i) - frac (real a * \vartheta i) | < 1/N)"
  proof - [35 lines]
  qed
  then obtain a b where "a<b" "b \leq N^n" and *: "\Lambdai. i<n \implies {frac (real b * \vartheta i) - frac (real a * \vartheta i) { < 1/N"
    by blast
  let ?k = "b-a"
  let ?h = "\lambdai. |b * \vartheta i| - |a * \vartheta i|"
  show ?thesis
  proof
    fix i
    assume "i<n"
    have "frac (b * \vartheta i) - frac (a * \vartheta i) = ?k * \vartheta i - ?h i"
      using <a < b> by (simp add: frac def left diff distrib' of nat diff)
    then show "of int ?k * \vartheta i - ?h i < 1/N"
       by (metis "*" <i < n> of int of nat eq)
  qed (use <a < b> <b < N^n> in auto)
qed
```

The chapter continues ...

- refinements to Dirichlet's approximation theorem
- Liouville's approximation theorem (done elsewhere)
- Kronecker's approximation theorem
- ... and the simultaneous version of Kronecker
- Advanced examples, e.g. to periodic functions

Theorem 7.13. If f has three periods $\omega_1, \omega_2, \omega_3$ which are linearly independent over the integers, then f has arbitrarily small nonzero periods.

PROOF. Suppose first that ω_2/ω_1 is real. If ω_2/ω_1 is rational then ω_1 and ω_2 are linearly dependent over the integers, hence $\omega_1, \omega_2, \omega_3$ are also dependent, contradicting the hypothesis. If ω_2/ω_1 is irrational, then f has arbitrarily small nonzero periods by Theorem 7.12.

Now suppose ω_2/ω_1 is not real. Geometrically, this means that ω_1 and ω_2 are not collinear with the origin. Hence ω_3 can be expressed as a linear combination of ω_1 and ω_2 with real coefficients, say

 $\omega_3 = \alpha \omega_1 + \beta \omega_2$, where α and β are real.

Now we consider three cases:

(a) Both α and β rational.

(b) One of α , β rational, the other irrational.

(c) Both α and β irrational.

Case (a) implies $\omega_1, \omega_2, \omega_3$ are dependent over the integers, contradicting the hypothesis.

For case (b), assume α is rational, say $\alpha = a/b$, and β is irrational. Then we have

 $\omega_3 = \frac{a}{b}\omega_1 + \beta\omega_2$, so $b\omega_3 - a\omega_1 = \beta(b\omega_2)$.

This gives us two periods $b\omega_3 - a\omega_1$ and $b\omega_2$ with irrational ratio, hence f has arbitrarily small periods. The same argument works, of course, if β is rational and α is irrational.

Now consider case (c), both α and β irrational. Here we consider two subcases.

Oops!

- Unfortunately, Apostol set things up for Kronecker's theorem when actually he needed Dirichlet's
- Despite including a redundant case analysis, he *didn't* establish the preconditions for Kronecker's theorem

... and he hadn't bothered to present Dirichlet's in its simultaneous form

```
theorem
  fixes f:: "complex \Rightarrow complex" and \omega 1 \ \omega 2 \ \omega 3:: complex
  assumes \omega: "is periodic \omega 1 f" "is periodic \omega 2 f" "is periodic \omega 3 f"
     and indp: "module.independent (\lambda r. (*) (complex of int r)) {\omega 1, \omega 2, \omega 3}"
     and dist: "distinct [\omega 1, \omega 2, \omega 3]"
     and "\varepsilon > 0"
  obtains \omega where "is periodic \omega f" "0 < cmod \omega" "cmod \omega < \varepsilon"
proof -
  interpret C: Modules.module "(\lambdar. (*) (complex of int r))"
     by (simp add: Modules.module.intro distrib left mult.commute)
  have nz: "\omega 1 \neq 0" "\omega 2 \neq 0" "\omega 3 \neq 0"
     using indp C.dependent zero by force+
  show thesis
  proof (cases "\omega 2/\omega 1 \in \mathbb{R}") [16 lines]
  next
     case False
     then obtain \alpha \beta where \alpha\beta: "\omega3 = of real \alpha * \omega1 + of real \beta * \omega2"
        using complex is Real iff gen lattice.\omega 1 \omega 2 decompose gen lattice.intro by blast
     show ?thesis
     proof (cases "\alpha \in \mathbb{O}")
        case True
        then obtain m1 n1 where mn1: "\alpha = of int m1 / of int n1" and "n1 > 0"
          by (meson Rats cases')
        show ?thesis
        proof (cases "\beta \in \mathbb{Q}")
          case True
          then obtain m2 n2 where mn2: "\beta = of int m2 / of int n2" and "n2 > 0"
             by (meson Rats cases')
          have "of int(m1*n2)*\omega1 + of int(m2*n1)*\omega2 + of int(-n1*n2)*\omega3 = 0"
             using \alpha\beta (n1 > 0) (n2 > 0) by (simp add: mn1 mn2 add frac eq)
          then have "C.dependent \{\omega 1, \omega 2, \omega 3\}" [5 lines]
          with indp show ?thesis
             by blast
        next
          case False
          define \omega where "\omega \equiv n1 * \omega 3 - m1 * \omega 1"
          have "\omega = \beta * (n1 * \omega 2)"
             using \langle n1 \rangle \rangle by (simp add: \omega def \alpha\beta mn1 algebra_simps)
          moreover have "is periodic \omega f" "is periodic (n1 * \omega2) f"
             by (simp all add: \omega \,\,\omega def is periodic diff is periodic times int)
          ultimately show ?thesis
             using that \langle \beta \notin Q \rangle nz \langle 0 < n1 \rangle \langle \varepsilon > 0 \rangle small periods real irrational [of "n1*\omega2" f \omega \varepsilon]
             by auto
        qed
```

- * This first part covers when ω_2/ω_1 is real, and if not obtains real α and β where $\omega_3 = \alpha \omega_1 + \beta \omega_2$.
- * Then it considers whether α (or β) is rational.
- In the final case, both are irrational and there is a big calculation using Dirichlet's approximation theorem

```
case False
      show ?thesis
      proof (cases "\beta \in \mathbb{Q}") [11 lines]
      next
         case False
         show ?thesis
         proof -
            define \vartheta where "\vartheta \equiv case nat \alpha (\lambda . \beta)"
            define \delta where "\delta \equiv \varepsilon / (1 + cmod \omega1 + cmod \omega2)"
            have "\delta > 0"
               by (smt (verit, best) \delta def \langle \varepsilon > 0 \rangle divide pos pos norm not less zero)
            obtain N where N: "1 / real N < \delta" and "N>0"
               by (meson \langle 0 < \delta \rangle nat approx posE zero less Suc)
            then obtain k q where kh: "\wedgei. i < 2 \implies {of int k * \vartheta i - of int (q i)} < \delta" and "\theta < k"
               by (metis Dirichlet approx simult[of N 2 \vartheta] less trans)
            define h1 where "h1 \equiv q 0" define h2 where "h2 \equiv q 1"
            have "cmod (\mathbf{k} * \alpha * \omega \mathbf{1} - \mathbf{h}\mathbf{1} * \omega \mathbf{1}) = \text{cmod} (\mathbf{k} * \alpha - \mathbf{h}\mathbf{1}) * \text{cmod} \omega \mathbf{1}"
               by (metis left diff distrib norm mult of real diff of real of int eq)
            also have "... = abs (k * \alpha - h1) * cmod \omega1"
               by (metis norm of real)
            also have "... < \delta * cmod \omega1"
               using kh [of 0] by (simp add: \vartheta def nz h1 def)
            finally have 1: "norm (k * \alpha * \omega1 - h1 * \omega1) < \delta * cmod \omega1".
            have "cmod (k * \beta * \omega<sup>2</sup> - h<sup>2</sup> * \omega<sup>2</sup>) = cmod (k * \beta - h<sup>2</sup>) * cmod \omega<sup>2</sup>"
               by (metis left diff distrib norm mult of real diff of real of int eq)
            also have "... = abs (\mathbf{k} * \beta - \mathbf{h2}) * \text{ cmod } \omega \mathbf{2}"
               by (metis norm of real)
            also have "... < \delta * cmod \omega2"
               using kh [of 1] by (simp add: \vartheta def nz h2 def)
            finally have 2: "cmod (k * \beta * \omega2 - h2 * \omega2) < \delta * cmod \omega2".
            define \omega where "\omega \equiv k * \omega 3 - h1 * \omega 1 - h2 * \omega 2"
            have "\omega = (k * \alpha * \omega 1 - h1 * \omega 1) + (k * \beta * \omega 2 - h2 * \omega 2)"
               by (simp add: \omega def \alpha\beta algebra simps)
            then have "cmod \omega < \text{cmod}(\mathbf{k} * \alpha * \omega \mathbf{1} - \mathbf{h}\mathbf{1} * \omega \mathbf{1}) + \text{cmod}(\mathbf{k} * \beta * \omega \mathbf{2} - \mathbf{h}\mathbf{2} * \omega \mathbf{2})"
               using norm triangle ineq by blast
            also have "... < \delta * cmod \omega1 + \delta * cmod \omega2"
               using "1" "2" by linarith
            also have "... < \varepsilon"
               using \langle \varepsilon \rangle > 0 \rangle nz
               by (simp add: \delta def divide simps) (auto simp add: distrib left pos add strict)
            finally have "cmod \omega < \varepsilon".
            have "is periodic \omega f"
               by (simp add: \omega \omega def is periodic diff is periodic times int)
            moreover have "\omega \neq 0" [9 lines]
            ultimately show ?thesis
               by (simp add: \langle \text{cmod } \omega \langle \varepsilon \rangle that)
         qed
      qed
   qed
qed
```

qed

Apostol's application to the Riemann zeta function

- * Obtaining the inf and sup of $|\zeta(\sigma + it)|$ where σ is held constant
- Apostol's proof contains a circularity that I broke with the help of an elaborate argument from *MathOverflow*
- You also need to understand that *s* is synonymous with $\sigma + it$ (*except when it isn't*)

Definition. For fixed σ , we define

$$m(\sigma) = \inf_{i=1}^{\infty} |\zeta(\sigma + it)|$$
 and $M(\sigma) = \sup_{i=1}^{\infty} |\zeta(\sigma + it)|$,

where the infimum and supremum are taken over all real t.

Theorem 7.11. For each fixed $\sigma > 1$ we have

$$M(\sigma) = \zeta(\sigma)$$
 and $m(\sigma) = \frac{\zeta(2\sigma)}{\zeta(\sigma)}$.

PROOF. For $\sigma > 1$ we have $|\zeta(\sigma + it)| \leq \zeta(\sigma)$ so $M(\sigma) = \zeta(\sigma)$, the supremum being attained on the real axis. To obtain the result for $m(\sigma)$ we estimate the reciprocal $|1/\zeta(s)|$. For $\sigma > 1$ we have

(17)
$$\left|\frac{1}{\zeta(s)}\right| = \prod_{p} |1 - p^{-s}| \le \prod_{p} (1 + p^{-\sigma}) = \frac{\zeta(\sigma)}{\zeta(2\sigma)}.$$

Hence $|\zeta(s)| \ge \zeta(2\sigma)/\zeta(\sigma)$ so $m(\sigma) \ge \zeta(2\sigma)/\zeta(\sigma)$.

Now we wish to prove the reverse inequality $m(\sigma) \leq \zeta(2\sigma)/\zeta(\sigma)$. The idea is to show that the inequality

$$|1 - p^{-s}| \le 1 + p^{-s}$$

used in (17) is very nearly an equality for certain values of t. Now

$$1 - p^{-s} = 1 - p^{-\sigma - it} = 1 - p^{-\sigma} e^{-it \log p} = 1 + p^{-\sigma} e^{i(-t \log p - \pi)},$$

so we need to show that $-t \log p - \pi$ is nearly an even multiple of 2π for certain values of t. For this we invoke Kronecker's theorem. Of course, there are infinitely many terms in the Euler product for $1/\zeta(s)$ and we cannot expect to make $-t \log p - \pi$ nearly an even multiple of 2π for all primes p. But we will be able to do this for enough primes to obtain the desired inequality.

```
theorem Inf 7 11:
  fixes \sigma:: real
  assumes "\sigma > 1"
  shows "(INF t. cmod (zeta(Complex \sigma t))) = Re (zeta (2 * \sigma)) / Re (zeta \sigma)" (is "Inf ?F = ?rhs")
proof (intro antisym)
  show rhs le INFF: "?rhs < Inf ?F"</pre>
    by (metis Complex eq UNIV I empty iff norm zeta same Im ge [OF assms] cINF greatest)
  interpret Modules.module "(\lambdar. (*) (real of int r))"
    by (simp add: Modules.module.intro distrib left mult.commute)
  define pr where "pr \equiv enumerate {p::nat. prime p}"
    — <enumeration of the primes, starting from 0 (not 1 as in the text) >
  have [simp]: "strict mono pr"
    by (simp add: pr def primes infinite strict mono enumerate)
  have prime iff: "prime n \leftrightarrow (\exists k. n = pr k)" for n
    using enumerate Ex enumerate in set pr def primes infinite by blast
  then have pnp: "prime (pr k)" for k
    using prime iff by blast
  then have pr gt1: "real (pr k) > 1" for k
    by (metis of nat 1 of nat less iff prime gt 1 nat)
  have pr gt: "n < pr n" for n [9 lines]</pre>
  define \vartheta where "\vartheta \equiv \lambda k. - ln (pr k) / (2 * pi)"
  have "inj pr"
    by (simp add: strict mono imp inj on)
  then have inj\vartheta: "inj \vartheta"
    by (auto simp: inj_on_def θ_def pnp prime_gt_0_nat)
  have [simp]: "pr 0 = 2"
    by (simp add: pr def enumerate.simps Least_equality prime_ge_2_nat)
  have prod if prime eq: "(\prod p \le pr n. if prime p then w p else 1) = (\prod k \le n. w (pr k))" (is "?L=?R") [12 lines]
  have prod if prime eq real: "(\prod p \le pr n. if prime p then w p else 1) = (\prod k \le n. w (pr k))" [9 lines]
  have zeta nz: "zeta (Complex \sigma t) \neq 0" for t
    using assms complex.sel(1) zeta Re gt 1_nonzero by presburger
  then obtain zeta pos: "Re (zeta \sigma) > 0" "Re (zeta (of real \sigma * 2)) > 0"
    by (smt (verit) Complex eq Re complex of real assms complex of real def mult 2 right
             norm zeta same Im le of real add zero less norm iff zeta Re gt 1 nonzero)
  then have rhs pos: "?rhs > 0"
    by (auto simp: field simps)
  with rhs le INFF have INFF pos: "Inf ?F > 0"
    by linarith
```

- Here we see only the boilerplate from the first part of this 400 line proof.
- The main part is the calculation of the required *t*,
 including the removal of the circular dependence
- Apostol is good at conveying general ideas, but terrible with details
- ... let's see some examples from Chapter 1

Three trivial proofs

Theorem 1.5. If an elliptic function f has no zeros in some period parallelogram, then f is constant.

PROOF. Apply Theorem 1.4 to the reciprocal 1/f.

Over 60 lines of dense calculations

Theorem 1.6. The contour integral of an elliptic function taken along the boundary of any cell is zero.

PROOF. The integrals along parallel edges cancel because of periodicity.

Theorem 1.7. The sum of the residues of an elliptic function at its poles in any period parallelogram is zero.

PROOF. Apply Cauchy's residue theorem to a cell and use Theorem 1.6.

Nearly 200 lines

Remarks and conclusions

- ✤ We did chapters 1–3 and 7.
- The material is straightforward to formalise, once you understand the conventions
- ... but errors and gaps waste lots of time.
- Expertise in the field you're formalising is necessary!