

# Automating elementary number-theoretic proofs using Gröbner bases

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## Divisibility properties over the integers

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Often want to prove tedious lemmas like

$$\forall a n x y. ax \equiv ay \pmod{n} \wedge \text{coprime}(a, n) \Rightarrow x \equiv y \pmod{n}$$

## Expanding divisibility properties

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Eliminate divisibility notions in terms of existentials:

- $s \mid t$  to  $\exists d. t = sd$
- $s \equiv t \pmod{u}$  to  $\exists d. t - s = ud$
- $\text{coprime}(s, t)$  to  $\exists x y. sx + ty = 1$ .

## Applied to the example

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$$\begin{aligned}\forall a n x y. (\exists d. ay - ax = nd) \wedge \\ (\exists u v. au + nv = 1) \\ \Rightarrow (\exists e. y - x = ne)\end{aligned}$$

Pull out the quantifiers in the antecedent:

$$\forall a n x y d u v. ay - ax = nd \wedge au + nv = 1 \Rightarrow \exists e. y - x = ne$$

## Solving a more general problem

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Instead, attempt to prove the property holds *in all rings*.

It turns out that this problem *is* decidable using well-known methods.

## Word problem for rings

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$$\forall \bar{x}. p_1(\bar{x}) = 0 \wedge \cdots \wedge p_n(\bar{x}) = 0 \Rightarrow q(\bar{x}) = 0$$

holds in all rings iff

$$q \in \mathbf{Id}_{\mathbb{Z}} \langle p_1, \dots, p_n \rangle$$

i.e. there exist ‘cofactor’ polynomials with integer coefficients such that

$$p_1 \cdot q_1 + \cdots + p_n \cdot q_n = q$$

## Generalizes to linear existential theorems

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$$\forall \bar{x}. \bigwedge_{i=1}^m e_i(\bar{x}) = 0 \Rightarrow \exists y_1 \cdots y_n. p_1(\bar{x})y_1 + \cdots + p_n(\bar{x})y_n = a(\bar{x})$$

holds in all rings iff (Horn-Herbrand) there are terms in the language of rings s.t.

$$\text{Ring} \vdash \forall \bar{x}. \bigwedge_{i=1}^m e_i(\bar{x}) = 0 \Rightarrow p_1(\bar{x})t_1(\bar{x}) + \cdots + p_n(\bar{x})t_n(\bar{x}) = a(\bar{x})$$

iff (previous theorem)

$$a \in \text{Id}_{\mathbb{Z}} \langle e_1, \dots, e_m, p_1, \dots, p_n \rangle$$



## ... and simultaneous linear existentials

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$$\forall \bar{x}. \bigwedge_{i=1}^m e_i(\bar{x}) = 0 \Rightarrow \exists y_1 \cdots y_n. p_{11}(\bar{x})y_1 + \cdots + p_{1n}(\bar{x})y_n = a_1(\bar{x}) \wedge$$
$$\cdots \wedge$$
$$p_{k1}(\bar{x})y_1 + \cdots + p_{kn}(\bar{x})y_n = a_k(\bar{x})$$

holds in all rings iff

$$(a_1 u_1 + \cdots + a_k u_k)$$
$$\in \text{Id}_{\mathbb{Z}} \langle e_1, \dots, e_m, (p_{11}u_1 + \cdots + p_{k1}u_k), (p_{1n}u_1 + \cdots + p_{kn}u_k) \rangle$$

where the  $u_i$  are fresh variables.

## Solving ideal membership problems

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The most natural approach to solving ideal membership problem is Gröbner bases.

Strictly, should use an integer version. However, can use the rational version speculatively and see if we get integer cofactors.

With an instrumented version of Buchberger's algorithm, can generate cofactors and hence easily generate a rigorous formal proof.

In our example

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This is true because

$$y - x = (ay - ax - nd) \cdot u + (au + nv - 1) \cdot (x - y) + n \cdot (ud + vy - vx)$$

## Extensions

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- Use linear equations  $x + a = b$  to substitute directly
- Add greatest common divisors by characterizing theorem  
 $g \mid a \wedge g \mid b \wedge (\exists u v. au + bv = g)$
- Solve for existential witnesses sequentially to defer nonlinear ones.

## Implementation in HOL Light

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A prototype of the procedure is available in the latest release of HOL Light, 2.20:

```
# INTEGER_RULE
  `!a1 a2 n1 n2:int.
    (a1 == a2) (mod (gcd(n1,n2)))
    ==> ?x. (x == a1) (mod n1) /\ (x == a2) (mod n2)`;;
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4 basis elements and 1 critical pairs
5 basis elements and 0 critical pairs
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Translating certificate to HOL inferences
val it : thm =
|- !a1 a2 n1 n2.
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  ==> (?x. (x == a1) (mod n1) /\ (x == a2) (mod n2))
```

## Various successful examples

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$$d|a \wedge d|b \Rightarrow d|(a - b)$$

$$\text{coprime}(d, a) \wedge \text{coprime}(d, b) \Rightarrow \text{coprime}(d, ab)$$

$$\text{coprime}(d, ab) \Rightarrow \text{coprime}(d, a)$$

$$\text{coprime}(a, b) \wedge x \equiv y \pmod{a} \wedge x \equiv y \pmod{b} \Rightarrow x \equiv y \pmod{ab}$$

$$m|r \wedge n|r \wedge \text{coprime}(m, n) \Rightarrow (mn)|r$$

$$\text{coprime}(xy, x^2 + y^2) \Leftrightarrow \text{coprime}(x, y)$$

$$\text{coprime}(a, b) \Rightarrow \exists x. x \equiv u \pmod{a} \wedge x \equiv v \pmod{b}$$

$$ax \equiv ay \pmod{n} \wedge \text{coprime}(a, n) \Rightarrow x \equiv y \pmod{n}$$

$$\text{gcd}(a, n) | b \Rightarrow \exists x. ax \equiv b \pmod{n}$$



## Failures

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Can't solve problems where special properties of the integers are used

$$2|x^2 + x$$

This fails over some rings, e.g.  $\mathbb{R}[x]$ .

However, such examples very seldom appear in typical routine lemmas.

## Conclusions

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