Maru: SGX-Spark Deep Dive

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Trust Issues: Provider Perspective

Cloud provider does not trust users

Use virtual machines to isolate users from each other and the host

VMs only provide one way protection
Trust Issues: User Perspective

Users trust their applications

Users must implicitly trust cloud provider

Existing applications implicitly assume trusted operating system
Trusted Execution Support with Intel SGX

- Users create HW-enforced trusted environment (enclave)
- Supports unprivileged user code
- Protects against strong attacker model
- Remote attestation
- Available on commodity CPUs
Trusted Execution Environments

Trusted execution environment (TEE) in process

- Own code & data
- Controlled entry points
- Provides confidentiality & integrity
- Supports multiple threads
- Full access to application memory
Overview of Intel SGX
Intel Software Guard Extensions (SGX)

Extension of Instruction Set Architecture (ISA) in recent Intel CPUs

Protects confidentiality and integrity of code & data in untrusted environments
  – Platform owner considered malicious
  – Only CPU chip and isolated region trusted
SGX Enclaves

SGX introduces notion of enclave
- Isolated memory region for code & data
- New CPU instructions to manipulate enclaves and new enclave execution mode

Enclave memory encrypted and integrity-protected by hardware
- Memory encryption engine (MEE)
- No plaintext secrets in main memory

Enclave memory can be accessed only by enclave code
- Protection from privileged code (OS, hypervisor)

Application has ability to defend secrets
- Attack surface reduced to just enclaves and CPU
- Compromised software cannot steal application secrets
SGX Enclave API & Operation

Enclave interface functions: **E Calls** to provide input data to enclave

Calls outside enclave: **O Calls** to return results from enclave

Constitute enclave boundary interface
**Enclave Page Cache (EPC)**

Physical memory region protected by MEE
- EPC holds enclave contents

Shared resource between all enclaves running on platform
- Currently only 128 MB
- ~96 MB available to user, rest for metadata

Content encrypted while in DRAM, decrypted when brought to CPU
- Plaintext in CPU caches
SGX allows multiple threads to enter same enclave simultaneously
- One thread control structure (TCS) per thread
- Part of enclave, reflected in measurement

TCS limits number of enclave threads
- Upon thread entry TCS is blocked and cannot be used by another thread

Each TCS contains address of entry point
- Prevents jumps into random locations inside of enclave
SGX Paging

SGX provides mechanism to evict EPC page to unprotected memory
- EPC limited in size

Paging performed by OS
- Validated by HW to prevent attacks on address translations
- Metadata (MAC, version) kept within EPC

Accessing evicted page results in page fault
- Page is brought back into EPC by OS
- Hardware verifies integrity of page
- Another page might be evicted if EPC is full
SGX SDK Code Sample

```c
char request_buf[BUFFER_SIZE];
char response_buf[BUFFER_SIZE];

int main()
{
    ...
    while(1)
    {
        receive(request_buf);
        ret = EENTER(request_buf, response_buf);
        if (ret < 0)
            fprintf(stderr, "Corrupted message\n");
        else
            send(response_buf);
    }
    ...
}
```

Server:
- Receives encrypted requests
- Processes them in enclave
- Sends encrypted responses

---

```c
char input_buf[BUFFER_SIZE];
char output_buf[BUFFER_SIZE];

int process_request(char *in, char *out)
{
    copy_msg(in, input_buf);
    if(verify_MAC(input_buf))
    {
        decrypt_msg(input_buf);
        process_msg(input_buf, output_buf);
        encrypt_msg(output_buf);
        copy_msg(output_buf, out);
        EEXIT(0);
    } else
        EEXIT(-1);
}
```
Enclave populated using special instruction (EADD)
• Contents initially in untrusted memory
• Copied into EPC in 4KB pages
Both data & code copied before execution commences in enclave
SGX Enclave Construction

Enclave contents distributed in plaintext
- Must not contain any (plaintext) confidential data

Secrets provisioned after enclave constructed and integrity verified

Problem: what if someone tampers with enclave?
- Contents initially in untrusted memory

```c
int process_request(char *in, char *out)
{
    copy_msg(in, input_buf);
    if(verify_MAC(input_buf))
    {
        decrypt_msg(input_buf);
        process_msg(input_buf, output_buf);
        encrypt_msg(output_buf);
        copy_msg(output_buf, out);
        EEXIT(0);
    }
    else
    
    EEXIT(-1);
}
```

Write unencrypted response to outside memory
CPU calculates enclave measurement hash during enclave construction
- Each new page extends hash with page content and attributes (read/write/execute)
- Hash computed with SHA-256

Measurement can be used to attest enclave to local or remote entity

CPU calculates enclave measurement hash during enclave construction
Different measurement if enclave modified

Different measurement if enclave modified
SGX Enclave Attestation

Is my code running on remote machine intact?
Is code really running inside an SGX enclave?

**Local** attestation
- Prove enclave’s identity (= measurement) to another enclave on same CPU

**Remote** attestation
- Prove enclave’s identity to remote party

Once attested, enclave can be trusted with secrets
Local Attestation

Prove identity of A to local enclave B

1. Target enclave B measurement required for key generation
2. Report contains information about target enclave B, including its measurement
3. CPU fills in report and creates MAC using report key, which depends on random CPU fuses and target enclave B measurement
4. Report sent back to target enclave B
5. Verify report by CPU to check that generated on same platform, i.e. MAC created with same report key (available only on same CPU)
6. Check MAC received with report and do not trust A upon mismatch
Remote Attestation

Transform local report to remotely verifiable “quote”

Based on provisioning enclave (PE) and quoting enclave (QE)
- Architectural enclaves provided by Intel
- Execute locally on user platform

Each SGX-enabled CPU has unique key fused during manufacturing
- Intel maintains database of keys
Remote Attestation

PE communicates with Intel attestation service
- Proves it has key installed by Intel
- Receives asymmetric attestation key

QE performs local attestation for enclave
- QE verifies report and signs it using attestation key
- Creates quote that can be verified outside platform

Quote and signature sent to remote attester, which communicates with Intel attestation service to verify quote validity
SGX Limitations & Research Challenges

Amount of memory enclave can use needs to be known in advance
- Dynamic memory support in SGX v2

Security guarantees not perfect
- Vulnerabilities within enclave can still be exploited
- Side-channel attacks possible

Performance overhead
- Enclave entry/exit costly
- Paging very expensive

Application partitioning? Legacy code? …
SGX-Spark
Secure Big Data Processing

Processing of large amounts of sensitive information
Outsourcing of data storage and processing
Cloud provider can access processed data
- Not acceptable for number of industries

```scala
def main(args: Array[String]) {
  new SparkContext(new SparkConf()
  .textFile(args(0))
  .flatMap(line => {line.split(" ")})
  .map(word => {(word, 1)})
  .reduceByKey{case (x, y) => x + y}
  .saveAsTextFile(args(1))
}
```

Distributed File System

Peter Pietzuch - Imperial College London
Secure machine learning (ML) killer application for Maru
- Resource-intensive thus good use case for cloud usage
- Raw training data comes with security implications

Complex implementations of ML algorithms cannot be adapted for SGX
- Consider Spark MLlib with 100s of algorithms

Challenges
- Extremely data-intensive domain
- Must support existing frameworks (Spark, TensorFlow, MXNet, CNTK, …)
- ML requires accelerators support (GPUs, TPUs, …)
- Prevention of side-channel attacks
State of the Art

Protect confidentiality and integrity of tasks and input/output data

**Opaque** [Zheng, NSDI 2017]
- Hide access patterns of distributed data analytics (Spark SQL)
- Introduces new oblivious relational operators
- Does not support arbitrary/existing Scala Spark jobs

**VC3** [Schuster, S&P 2015]
- Protects MapReduce Hadoop jobs
- Confidentiality/integrity of code/data; correctness/completeness of results
- No support for existing jobs → Re-implement for VC3
SGX Support for Spark

SGX-Spark
– Protect data processing from infrastructure provider
– Protect confidentiality & integrity of existing jobs
– No modifications for end users
– Acceptable performance overhead

Idea:
Execute only sensitive parts of Spark inside enclave
– Code that accesses/processes sensitive data

Code outside of enclave only accesses encrypted data
– Partition Spark
– Run two collaborating JVMs, inside enclave and outside of enclave
Supporting Managed Runtimes in SGX Enclaves

Many applications need runtime support
- JVM
- .NET
- JavaScript/V8/Node.js

Requires complex system support
- Dynamic library loading
- Filesystem support
- Signal handling
- Complete networking stack
SGX-LKL Architecture

Runs **unmodified Linux applications** in SGX enclaves
Applications and dependencies provided via **disk image**
**Full Linux kernel functionality** available

**Custom memory allocator**

**User-level threading**
- In-enclave synchronisation primitives
Challenges & Current State

1. Partitioning Spark

2. Data movement between JVMs

3. Memory efficiency
### 1. Partitioning Spark

**Goal:** Move minimal amount of Spark code to enclave

<table>
<thead>
<tr>
<th><strong>Outside</strong></th>
<th><strong>Enclave</strong></th>
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<td><strong>HadoopRDD</strong></td>
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<td>Execute user-provided function (f)</td>
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<td>(eg <code>flatMap(line =&gt; {line.split(&quot; &quot;)})</code>)</td>
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<tr>
<td>(i) Serialise user-provided function f</td>
<td></td>
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<tr>
<td>(ii) Send f and it to enclave JVM</td>
<td></td>
</tr>
<tr>
<td>(iv) Receive result iterator <code>it_result</code></td>
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1. Partitioning Spark
1. Partitioning Spark

- HadoopRDD A
  - <infile>
  - MapPartitions RDD C
    - k'
    - MapPartitions RDD B
      - j'
      - i'
      - HadoopRDD A
        - <outfile>

- Outside Enclave
  - Tasks A,B,C,D
  - Iterators i,j,k
  - Iterate via shm
1. Partitioning Spark

- HadoopRDD A
- Iterators via shm
- RDD B
- RDD C
- ResultTask
- ItConsumer k
- MapPartitions RDD B
- MapPartitions RDD C
- ItProvider i
- ItProvider k
- SgxTask B
- SgxTask C
- ItConsumer i

Communication

Tasks: A,B,C,D

Iterators: i,j,k

Outside

Enclave
2. Data Movement between JVMs

Goal: **Shared memory**

Use use host OS shared memory between two JVMs
- Outside access by enclave JVM

Manage shared memory between outside and enclave
Implement high-level read/write primitives
1. Partitioning Spark

- RDD
  - MapPartitions RDD
  - HadoopRDD
- Task
  - ResultTask
- SgxTask
  - SgxFirstTask
  - SgxOtherTask
  - SgxFct2
- ItConsumer
- ItProvider
  - Encryption
  - Encryption
- Shm Communication Manager
  - RingBuffer
  - Serialisation
- Shm Communication Manager
  - RingBuffer
  - Serialisation
- Host OS
  - shmem
2. Data Movement between JVMs

- **HadoopRDD A**
- **ItProvider i**
- **MapPartitions RDD B**
- **MapPartitions RDD C**
- **ResultTask**
- **ItConsumer k**

Iterate via shm

Enclave

Outside

Tasks: A,B,C,D

Iterators: i,j,k

Shm Communicator

Shm-enc-to-out

Shm-out-to-enc
3. Memory Efficiency

Only ~80 MB available
Native Spark Worker: ~320 MB 😞

Our Spark in-enclave code: 94 MB
  - With Java flags (-XX:InitialCodeCacheSize=2m
    -XX:ReservedCodeCacheSize=2m -Xms2m -Xmx3m
    -XX:CompressedClassSpaceSize=2m
    -XX:MaxMetaspaceSize=8m
    -XX:+UseCompressedClassPointers) → 50 MB

SGX-LKL:  
8 MB Kernel + 18 MB other → 26 MB
  - Working on memory efficiency
  - Eg thread stack size, kernel size, deactivating features
1. Security model for shielded data science jobs
   - How to harden shielded jobs? How to deal with vulnerabilities, bugs?
   - What about external dependencies/libraries?

2. Integration of language runtimes with secure enclaves
   - How to integrate SGX support for the JVM?
   - What is the right programming model for SGX enclaves?

3. Unikernel support for secure enclaves
   - How to support existing legacy binaries?
   - How to build type-safe minimal secure enclaves for data science jobs?

4. Prototype platform implementation and evaluation
   - Integration with Apache Flink or other dataflow frameworks

5. Dataflow attacks and mitigations strategies
   - What attacks are possible by observing encrypted dataflows?
   - Can we apply techniques for unobservable communication?
Primary System Goals:

1. Ensure **integrity & confidentiality** for tasks, input data & output results

2. Support **arbitrary** workloads and tasks

3. **Low performance overhead** (throughput and latency)
SGX-Spark: **Summary of V1**

Version 1 (V1):

1. **Prototype** of Apache Spark using TEEs (Intel SGX)
   - Shows the **feasibility** of the approach

2. Each **worker partitioned** into trusted and untrusted JVM:
   - Minimize the trusted computing base (TCB)
   - Provides **confidentiality** for input data & result

3. Use “**pull**” model for cross JVM communication:
   - Trusted JVM “asks” for data, objects, context
   - Blocking message requests → **high overhead**
SGX-Spark: **Goals of V2**

**Version 2 (V2):**

1. Provide **integrity** as well as confidentiality
   - Protect driver with TEE to enforce **job integrity**

2. Improve performance with a “**push**” model:
   - Push data into trusted JVM to avoid blocking requests (e.g. use pre-fetching and leverage EPC memory)
   - Batch messages (reduce communication costs)
   - Use lock free and highly parallelized comm. channels

3. Support generic RDDs / avoid assumptions about data layout (e.g. avoid DataFrames and DataSets)
SGX-Spark: Design of V2

High-level Overview:
SGX-Spark: Worker Execution

Worker Execution Diagram:

1. **ResultTask**
   - HadoopInput
   - MapPartitions
   - FilterPartitions
   - ResultConsumer

2. **Control Messages**
   - ClosureProvider
   - Socket
   - BatchedItProvider

3. **SGXTask**
   - Function
   - BatchedItConsumer
   - BatchedItProvider

4. **Unsecure JVM**
   - Lock free MPMC Queue (JNI)

5. **Enclave JVM**
**Spark Worker Execution:**

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