

# Logic and Proof - Supervision 1

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1. Prove the following two statements, using any method you like.
  - (a)  $p \rightarrow q \not\equiv q \rightarrow p$
  - (b)  $(p \rightarrow q) \rightarrow ((\neg p \rightarrow q) \rightarrow q) \simeq \top$
2. Reduce each of the following formulas to NNF, CNF, and DNF.
  - (a)  $(q \rightarrow \neg p) \rightarrow p$
  - (b)  $((p \wedge q) \vee r) \wedge \neg(p \vee r)$
3. For each formula  $\varphi$  below, either prove that it is valid, or give an interpretation satisfying  $\neg\varphi$ :
  - (a)  $(p \rightarrow r) \wedge (q \rightarrow r) \rightarrow (p \vee q \rightarrow r)$
  - (b)  $((p \rightarrow q) \rightarrow r) \rightarrow (p \rightarrow (q \rightarrow r))$
  - (c)  $(\neg p \vee (q \rightarrow p)) \rightarrow (\neg p \wedge q)$
4. Prove the following sequents, using the sequent calculus rules given in the notes:
  - (a)  $\neg\neg p \Rightarrow p$
  - (b)  $(p \wedge q) \wedge r \Rightarrow p \wedge (q \wedge r)$
  - (c)  $(p \vee q) \wedge (p \vee r) \Rightarrow p \vee (q \wedge r)$
  - (d)  $\neg(p \vee q) \Rightarrow \neg p \wedge \neg q$

5. Suppose  $\mathcal{L}$  is a first-order language with symbols  $F$  and  $M$ , of arity 2, and  $Ed$ , of arity 0. Let  $F(x, y)$  mean that  $x$  is the father of  $y$ , and let  $M(x, y)$  mean that  $x$  is the mother of  $y$ . The constant  $Ed$  denotes the person Ed. Write for each of the sentences below, a formula of  $\mathcal{L}$  with the same meaning:

- (a) Everybody has a mother.
- (b) Everybody has a mother and a father.
- (c) Whoever has a mother has a father.
- (d) Ed is a grandfather.
- (e) Nobody's grandmother is anybody's father.

6. Consider the formula

$$\varphi = \forall x \exists y \exists z (P(x, y) \wedge P(z, y) \wedge (P(x, z) \rightarrow P(z, x))).$$

Which of the following models satisfy  $\varphi$ ? Justify briefly.

- (a)  $\mathcal{M} = (\mathbb{N}, I)$  where  $I[P] = \{(m, n) \mid m, n \in \mathbb{N} \wedge m < n\}$
- (b)  $\mathcal{M}' = (\mathbb{N}, I)$  where  $I[P] = \{(m, 2 * m) \mid m \in \mathbb{N}\}$
- (c)  $\mathcal{M}'' = (\mathbb{N}, I)$  where  $I[P] = \{(m, n) \mid m, n \in \mathbb{N} \wedge m < n + 1\}$