

Ch.3. Constructions on Sets

cf. constructions on types
in prog. languages.

Sets as extensions of properties

$$\{x \mid P(x)\}$$

↑
a property

always a set?

↳ Russell's paradox (really a contradiction)

A safe way to build sets:

$$\{x \in S \mid P(x)\}$$

a set a property

But this needs sufficiently big sets S

By fiat: \mathbb{N} is a set

power set: $\mathcal{P}(X) = \{A \mid A \subseteq X\}$ is a set

↑



$$R = \{x \mid x \notin x\}$$

$$R \in R ?$$

$$R \in R \Rightarrow R \notin R$$

$$R \notin R \Rightarrow R \in R$$

Further constructions on sets

Let A, B be sets

Their product:

$$A \times B = \{ (a, b) \mid a \in A \text{ and } b \in B \}$$

↑
ordered pair:

$$(a, b) = (a', b') \text{ iff } a = a' \text{ and } b = b'$$

Can be realised as a set.

$$(a, b) = \{ \{a\}, \{a, b\} \}$$

Their disjoint union (or disjoint sum.)

$$A \uplus B = (\overset{\cup}{\{1\}} \times A) \cup (\overset{\cup}{\{2\}} \times B)$$

$(1, a)$ $(2, b)$

let I be a set s.t.

A_i is a set for all $i \in I$.

The big union.

$$\bigcup_{i \in I} A_i \quad \left[\text{or } \bigcup \{A_i \mid i \in I\} \right]$$
$$= \{x \mid \exists i \in I. x \in A_i\} \quad \text{is a set.}$$

The big intersection.

$$\bigcap_{i \in I} A_i \quad \left[\text{or } \bigcap \{A_i \mid i \in I\} \right]$$
$$= \{x \mid \forall i \in I. x \in A_i\}.$$

[Needs I non empty, or universe \mathcal{U} so empty intersections are \mathcal{U} .]

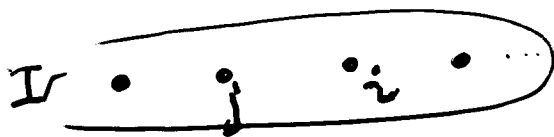
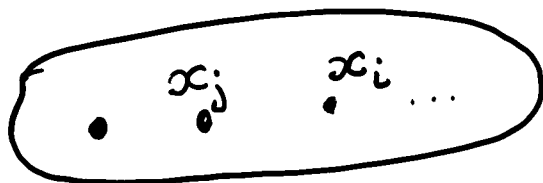
Indexed sets:

let I be a set s.t.

x_i is an element for all $i \in I$

Then,

$\{x_i \mid i \in I\}$ is a set.



Cantor's diagonal argument revisited

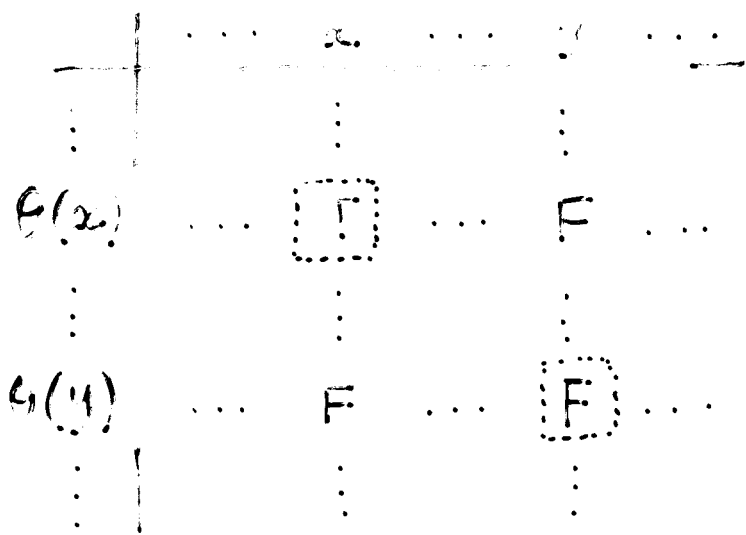
Let X be a set. There is no bijection $\theta: X \rightarrow \mathcal{P}(X)$.

Proof By contradiction. Suppose

$$\theta: X \rightarrow \mathcal{P}(X)$$

is a bijection.

Define $Y = \{x \in X \mid x \notin \theta(x)\}$.



As θ is a bijection, there is $y \in X$ such that $\theta(y) = Y$. But ...

Some Consequences

\leadsto set of relations between sets X and Y :

$$\mathcal{P}(X \times Y).$$

\leadsto set of partial functions from set X to set Y :

$$(X \rightarrow Y) = \{f \in \mathcal{P}(X \times Y) \mid f \text{ is a partial fn.}\}.$$

\leadsto set of total functions from set X to set Y :

$$(X \rightarrow Y) = \{f \in \mathcal{P}(X \times Y) \mid f \text{ is a function}\}.$$

Higher Order Logic:

If $(x \in X \Rightarrow e \in Y)$, then

$$\lambda x \in X. e \in (X \rightarrow Y)$$

$$= \{(x, e) \mid x \in X\}$$