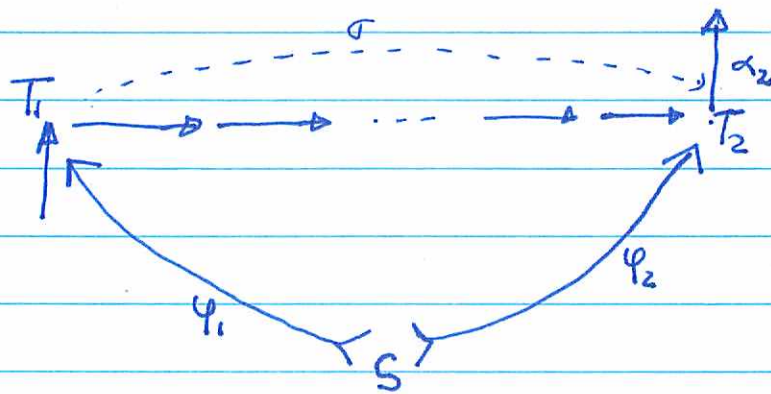


On Strong compression & simulations

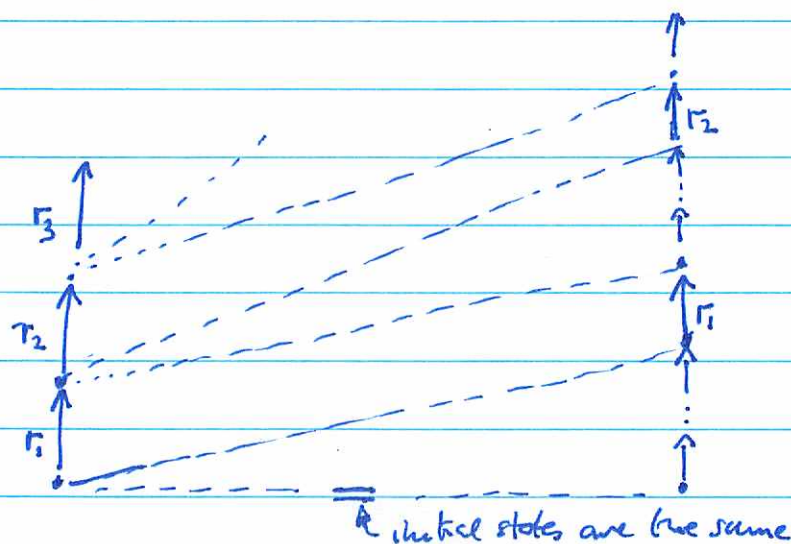
Imagine a strong compression step



associated with a rule $\alpha: S \rightarrow S'$. A strong compression step is possible even in the extreme situation where in the intermediate actions $T_1 \rightarrow \dots \rightarrow T_2$ all agents ~~can~~ have been ~~deleted~~ deleted and then re-created. Given this we can't expect any form of (weak) commutation relation to hold between ϕ_1, ϕ_2 and the intermediate actions. We might say, as I did in Barbados, that ~~what~~ ^{what} was required was a partial bijection σ between agents of T_1 and agents of T_2 that ~~comes~~ makes $\sigma \phi_1 \sim_{Ag} \phi_2$, and then proceed as for weak compression. BUT the fact that the same rule can be applied at T_1 and T_2 , that both embeddings ϕ_1 and ϕ_2 are present, directly enforces such a

partial bijection via S . So that approach contains lots of redundancy.

I think this leads to a view of string compression as a form of simulation:

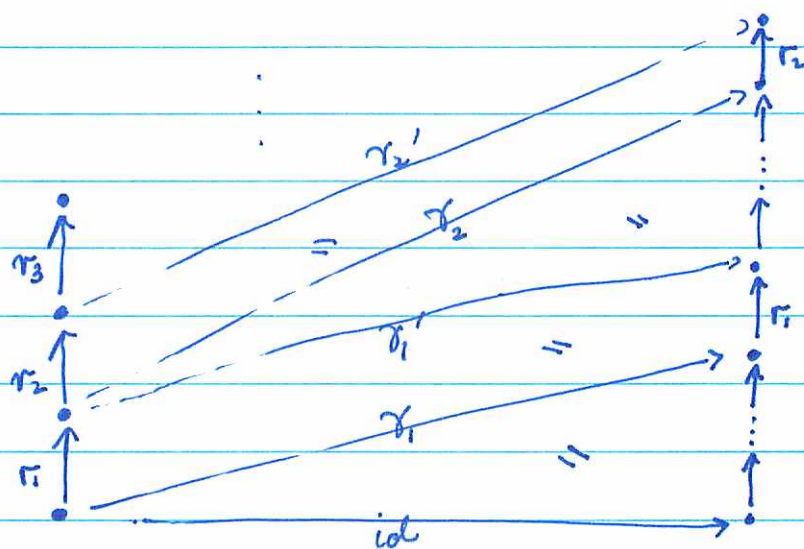


$r_1, r_2, \dots$
names ~~be~~ only
applied.

given simply by a (~~theoretical~~ set-theoretic) relation between states of the compressed trace and states of the original trace to states of the original trace.
[Do you agree?]

The earlier simple and weak compressors are associated with more refined forms of simulation.

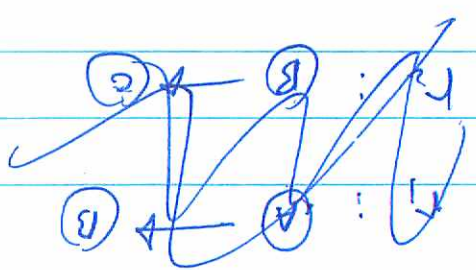
The simulation of simple compressor:



For simple compression ~~is given~~ the function between states is accompanied by action maps γ_i, γ_i' satisfying the commutativity conditions

For weak compression the function is accompanied by ~~trivial~~ partial injections on agents such the commutativity conditions only ~~to~~ need hold up to identity of agents.

Make sense?



Rules:

$r_0 : A \rightarrow B$

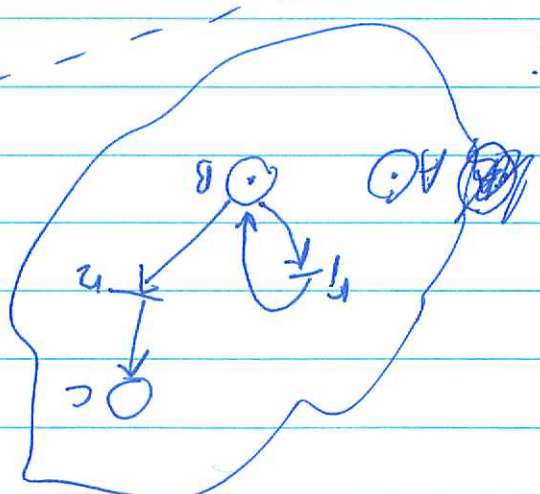
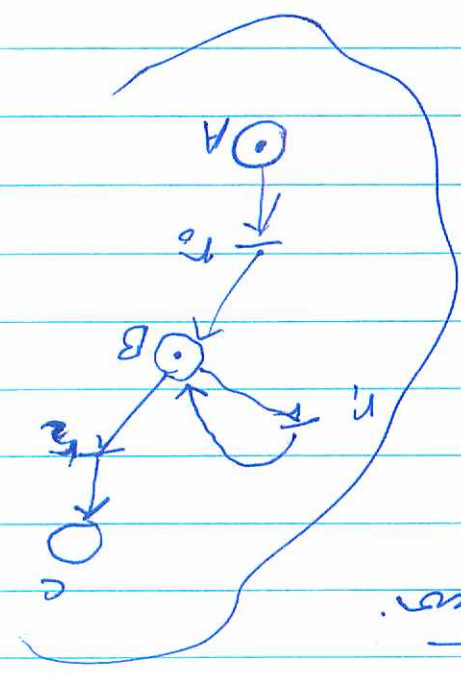
$r_1 : B \rightarrow B$

$r_2 : B \rightarrow C$

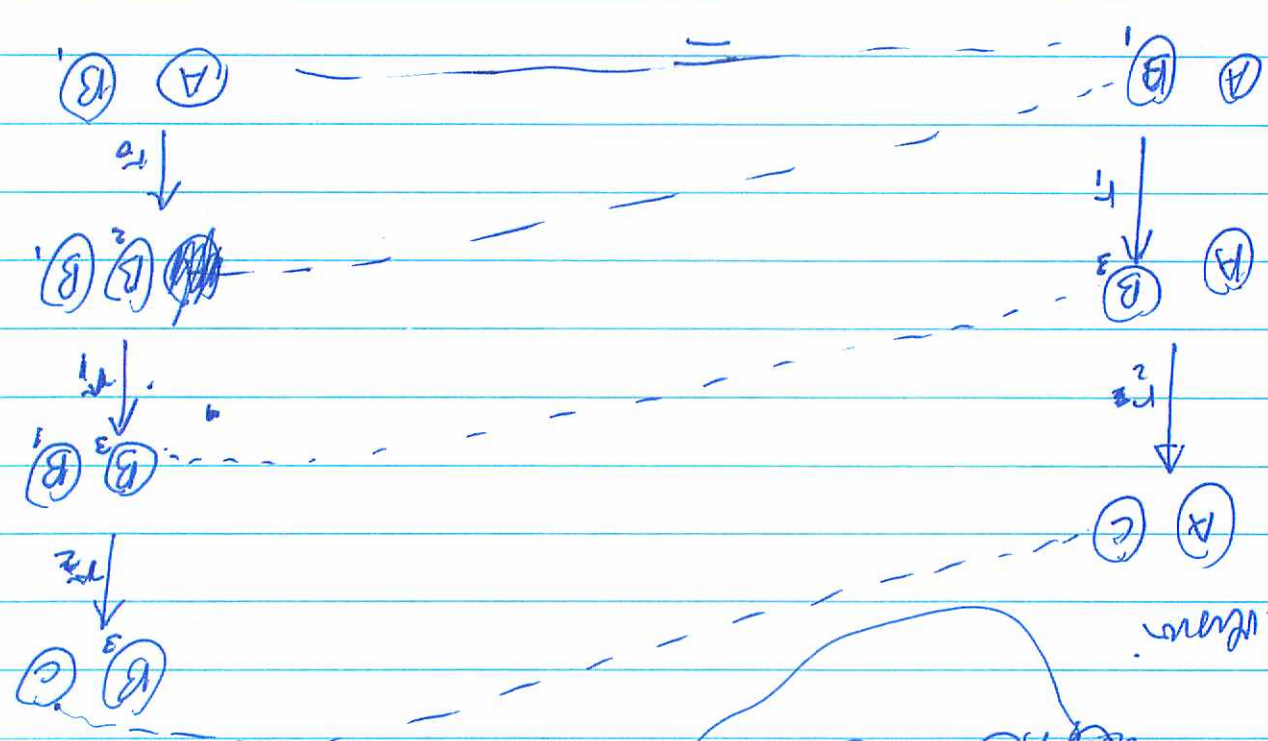
Initial state:

A, B

Example of a
strong compressor



Compressed version:



not a strong compressor

