# The Dataflow Model:

#### A Practical Approach to Balancing Correctness, Latency, and Cost in Massive-Scale, Unbounded, Out-of-Order Data Processing

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\*Not the Eric Schmidt you think...

### Problem

- Unbounded, unordered datasets
  - $\circ$  Web logs
  - Mobile usage statistics
  - Sensor networks
- Users have complex requirements:
  - Event-time ordering
  - Windowing by features of the data
  - Low latency
- One can never fully optimize along all dimensions of correctness, latency, and cost.
- How do we reconcile these conflicting requirements?

### **Previous Work: Need for Data Processing**

- Mapreduce, Hadoop, Pig, Hive, Spark enabled **scale**
- SQL Systems enabled
  - Query systems
  - $\circ$  Windowing
  - Data Streams
  - Time Domains
  - Semantic Models
- Spark streaming, Millwheel, Storm enabled low-latency processing

### But something is missing

Performance: Many good solutions but none have everything we want

- **High Latency** batch systems
- Not Fault Tolerant at Scale Aurora, TelegraphCQ, Niagara, Esper
- Fail on Correctness Pulsar, Storm, Samza (No Exactly once semantics)
- Lack Expressiveness MillWheel and Spark Streaming (Need for high-level models)
- **Too Complex** Lambda Architecture Systems (Need to maintain batch and stream)

Paradigm:

- Focus on input data as something which at some point will become complete
- Nearly all distinguish batch and streaming

### Key Aim of Paper: Shift In Approach

"Fully embrace the assumption that we never know if or when we have seen all of our data, only that data will arrive, old data may be retracted, and the only way to make the problem tractable is via principled abstractions that allow the practitioner the choice of appropriate tradeoffs between correctness, latency and cost."

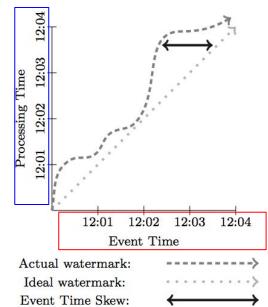
> "Execution engine [should not] dictate system semantics; properly designed and built batch, micro-batch, and streaming systems can all provide equal levels of correctness"

### **Contribution: The Dataflow Model**

- A Unified Model allowing:
  - Event-time ordered results windowed by features of the data themselves
  - Unbounded, unordered data source
  - Correctness, Latency, and Cost tunable
- Decomposes pipeline implementation across four related dimensions, providing clarity, composability and flexibility
  - What results are being computed
  - Where in event time they are being computed
  - Where in processing time they are materialized
  - How earlier results relate to later refinements
- Separates logic of data processing from the underlying physical implementation
  - $\circ$  choice of batch, micro-batch, or streaming engine  $\rightarrow$  correctness, latency, and cost.

### What time is it?

- **Event time** time at which **event actually occurred**, never changes (e.g. when someone searched for "dog")
- Processing time time at which event is observed at a given point during processing
  - changes as moves event moves through pipeline
- No global clock



### Primitives: What results are being computed

#### **Two Core Transforms**

- ParDo generic parallel processing
  - Translates well to unbounded data

(fix, 1), (fit, 2)

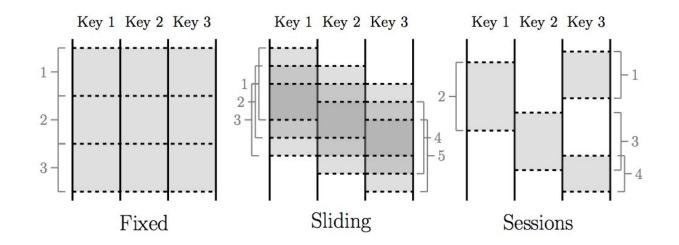
 $\begin{array}{c}
ParDo(\\ExpandPrefixes)\\(f,1),(fi,1),(fix,1),(f,2),(fi,2),(fit,2)\end{array}$ 

- **GroupByKey** grouping (key, value) pairs
  - Not so easy with unbounded data

(f, 1), (fi, 1), (fix, 1), (f, 2), (fi, 2), (fit, 2)  $\downarrow GroupByKey$ (f, [1, 2]), (fi, [1, 2]), (fix, [1]), (fit, [2])

#### Windowing Model: Where in event time results are computed

- Window: Time-based slices of dataset for processing as a group
- Aligned applied across all data
- Unaligned applied across given subset (e.g. per key)



#### Windowing Model: Where in event time results are computed

- Two operations
  - Set<Window> AssignWindows(T datum)

 $(k, v_1, 12:00, [0, \infty)), (k, v_2, 12:01, [0, \infty))$ 

 $\begin{array}{c} & AssignWindows(\\ & Sliding(2m, 1m)) \\ (k, v_1, 12:00, [11:59, 12:01)), \\ (k, v_1, 12:00, [12:00, 12:02)), \\ (k, v_2, 12:01, [12:00, 12:02)), \\ (k, v_2, 12:01, [12:01, 12:03)) \end{array}$ 

- Set<Window> MergeWindows(Set<Window> windows)
  - Typically redefine GroupByKey to GroupByKeyAndWindow

• Instead of (key,value) pairs, system is now handling (key, value, event time, window)

### Windowing Model: GroupByKeyAndWindow

 $(k_1, v_1, 13:02, [0, \infty)),$  $(k_2, v_2, 13:14, [0, \infty)),$  $(k_1, v_3, 13:57, [0, \infty)),$  $(k_1, v_4, 13:20, [0, \infty))$ AssignWindows( Sessions(30m))  $(k_1, v_1, 13:02, [13:02, 13:32)),$  $(k_2, v_2, 13:14, [13:14, 13:44)),$  $(k_1, v_3, 13:57, [13:57, 14:27)),$  $(k_1, v_4, 13:20, [13:20, 13:50))$ DropTimestamps  $(k_1, v_1, [13:02, 13:32)),$  $(k_2, v_2, [13:14, 13:44)),$  $(k_1, v_3, [13:57, 14:27)),$  $(k_1, v_4, [13:20, 13:50))$ GroupByKey  $(k_1, [(v_1, [13:02, 13:32)),$  $(v_3, [13:57, 14:27)),$  $(v_4, [13:20, 13:50))]),$  $(k_2, [(v_2, [13:14, 13:44))])$ 

 $(k_1, [(v_1, [13:02, 13:32)),$  $(v_3, [13:57, 14:27)),$  $(v_4, [13:20, 13:50))]),$  $(k_2, [(v_2, [13:14, 13:44))])$ MergeWindows( Sessions(30m))  $(k_1, [(v_1, [13:02, 13:50)),$  $(v_3, [13:57, 14:27)),$  $(v_4, [13:02, 13:50))]),$  $(k_2, [(v_2, [13:14, 13:44))])$ GroupAlsoByWindow  $(k_1, [([\mathbf{v}_1, \mathbf{v}_4], [13:02, 13:50)),$  $([v_3], [13:57, 14:27))]),$  $(k_2, [([\mathbf{v}_2], [13:14, 13:44))])$ *ExpandToElements*  $(k_1, [v_1, v_4], 13:50, [13:02, 13:50)),$  $(k_1, [v_3], 14:27, [13:57, 14:27)),$  $(k_2, [v_2], 13:44, [13:14, 13:44))$ 

### Windowing Model: In Practice

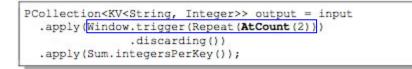
• E.g. Window data into 30 minute sessions

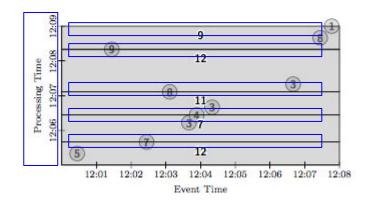
PCollection<KV<String, Integer>> input = IO.read(...); PCollection<KV<String, Integer>> output = input .apply(Sum.integersPerKey());

PCollection <kv<string,< th=""><th></th><th></th></kv<string,<>		
PCollection <kv<string,< td=""><td>Integer&gt;&gt;</td><th>output = input</th></kv<string,<>	Integer>>	output = input
.apply(Window.into(Se Duration.standard		a la constante 🖷 constante esta esta esta esta esta esta esta es
.apply(Sum.integersPe	erKey());	Contraction Contraction

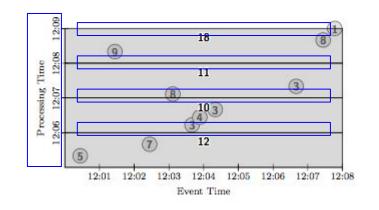
#### Triggering Model: When in processing time results are materialized

- Mechanism for stimulating the production of GroupByKeyAndWindow results in response to internal or external signals
- Allows you to control latency





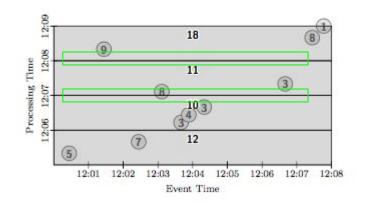
Collection <kv<string,< th=""><th>Integer&gt;&gt; (</th><th>output =</th><th>input</th><th></th></kv<string,<>	Integer>> (	output =	input	
.apply (Window.trigge	r (Repeat (Ati	Period(1,	MINUTE)	))
.discar	ding())			
.apply(Sum.integersPo	erKey());			

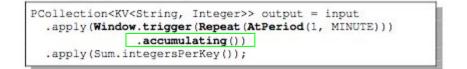


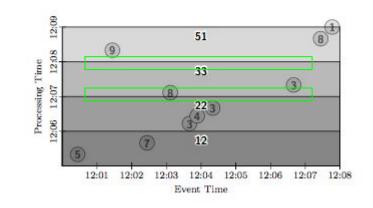
#### Incremental Model: How earlier results relate to later refinements

- Discarding
- Accumulating
- Accumulating and Retracting









## Putting it all together

What results are being computed

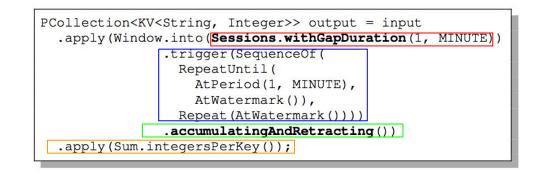
Where in event time they are being computed

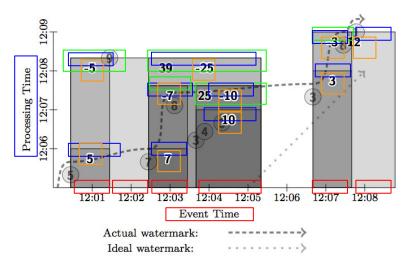
When in processing time they are materialized

How earlier results relate to later refinements

"Session windowing with 1 minute timeout, enabling retractions"

- Sessions joined as more data received
- Results retracted as more data received





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- Scalable implementations on FlumeJava and Millwheel

### How does it stack up?

- Low latency
  - via windowing and triggering
- Scalable and Fault Tolerant
  - Millwheel, FlumeJava
- Correctness
  - Incremental model with accumulations and retractions
- Greater Expressiveness
  - Windowing by features, complex triggering
- Reduced Complexity
  - Abstracted, Unified framework

### **But No Magic Bullet**

- That which was impractical in existing systems remains so
  - Framework for parallel computation independent of underlying execution engine
  - Balance latency, correctness for a problem
- Aimed at ease of use, pragmatic, real world massive scale data processing
- Hard to reason about the underlying performance.
- What is the **Complexity** of these operations?
- What is the **Overhead**?
- Abstractions mean less control
  - Where is my computation happening?
  - But that's the point of Dataflow Model...
  - Do I need to know?
- Paper doesn't explore how this model is to be **implemented** 
  - But open source is available

Thank You. Questions?