

10 Discrete Mathematics (js2878)

- (a) Let  $\Sigma = \mathbb{N}$  be the alphabet whose symbols are given by natural numbers  $0, 1, 2, \dots$ , and let  $L_n$  be the language over  $\Sigma$  consisting of finite strings of digits that sum to  $n$ .

(i) Draw a diagram of a DFA recognising  $L_3$ . [3 marks]

(ii) For arbitrary  $n \in \mathbb{N}$ , define a DFA  $M_n = (Q_n, \Sigma, \Delta_n, s_n, F_n)$  over  $\Sigma = \mathbb{N}$  that recognises the language  $L_n$ . [4 marks]

- (b) Let  $M = (Q, \Sigma, \Delta, s, F)$  be an NFA. We shall refer to  $M$  as **complete** when for any  $q \in Q$  and  $a \in \Sigma$ , there exists some  $q' \in Q$  such that  $q \xrightarrow{a} q'$ .

We shall refer to  $M$  as **partially deterministic** when for any states  $q, q', q'' \in Q$  and input symbol  $a \in \Sigma$ , if both  $q \xrightarrow{a} q'$  and  $q \xrightarrow{a} q''$  then  $q' = q''$  holds.

(i) Prove that an NFA is *deterministic* if and only if it is both *complete* and *partially deterministic*. [2 marks]

(ii) Suppose that  $M$  is a *partially deterministic* NFA; construct a DFA  $M' = (Q', \Sigma, \Delta', s', F')$  over the same alphabet such that any finite string  $u \in \Sigma^*$  is accepted by  $M$  if and only if it is accepted by  $M'$  and, moreover, such that the cardinality of  $Q'$  is bounded by  $\#Q' < 2^{\#Q}$  assuming  $\#Q > 1$ . Argue that  $M'$  meets these requirements. [8 marks]

- (c) Let  $M = (Q, \Sigma, \Delta, s, F)$  and  $M' = (Q', \Sigma, \Delta', s', F')$  be two DFAs over the same alphabet, writing  $\delta: Q \times \Sigma \rightarrow Q$  and  $\delta': Q' \times \Sigma \rightarrow Q'$  for the next-state functions corresponding to the total functional relations  $\Delta$  and  $\Delta'$  respectively.

A **homomorphism of DFAs** from  $M$  to  $M'$ , written  $f: M \rightarrow M'$ , is defined to be a function  $f: Q \rightarrow Q'$  satisfying the following conditions:

- $f$  preserves the starting state, *i.e.*  $fs = s'$ ;
- $f$  sends accepting states to accepting states, *i.e.* for  $q \in F$  we have  $fq \in F'$ ;
- $f$  preserves transitions, *i.e.*  $f(\delta(q, a)) = \delta'(f(q), a)$  for all  $q \in Q$  and  $a \in \Sigma$ .

- (i) Let  $M_1 = (Q_1, \Sigma, \Delta_1, s_1, F_1)$  and  $M_2 = (Q_2, \Sigma, \Delta_2, s_2, F_2)$  be two DFAs over  $\Sigma$ . Define a new DFA  $M_1 \times M_2$  over  $\Sigma$  with states in the cartesian product  $Q_1 \times Q_2$ , such that the projections  $\pi_1: Q_1 \times Q_2 \rightarrow Q_1$  and  $\pi_2: Q_1 \times Q_2 \rightarrow Q_2$  form homomorphisms  $\pi_1: M_1 \times M_2 \rightarrow M_1$  and  $\pi_2: M_1 \times M_2 \rightarrow M_2$  of DFAs, with proof. [3 marks]