## COMPUTER SCIENCE TRIPOS Part II – 2022 – Paper 9

## 5 Denotational Semantics (mpf23)

You may use standard results provided that you state them clearly.

(a) For a domain D, let fix be the function mapping a continuous function  $f \in (D \to D)$  to its least pre-fixed point  $fix(f) \in D$ .

Prove that  $fix: (D \to D) \to D$  is continuous. [4 marks]

(b) For a PCF type  $\tau$ , let  $\Omega_{\tau} = \mathbf{fix}(\mathbf{fn} \ x : \tau, x)$  and consider the following closed PCF terms of type  $(\tau \to \tau) \to (nat \to \tau)$ .

$$\begin{split} \mathbf{M}_{\tau} &= \mathbf{fn} \ f : \tau \to \tau. \ \mathbf{fn} \ n : nat. \ \mathbf{fix}(f) \\ \mathbf{N}_{\tau} &= \mathbf{fn} \ f : \tau \to \tau. \\ \mathbf{fix} \Big( \ \mathbf{fn} \ h : nat \to \tau. \ \mathbf{fn} \ n : nat. \\ f \Big( \mathbf{if} \ \mathbf{zero}(n) \ \mathbf{then} \ \Omega_{\tau} \ \mathbf{else} \ h(\mathbf{pred}(n)) \Big) \ \Big) \end{split}$$

Give an explicit description of the denotations  $\llbracket M_{\tau} \rrbracket$  and  $\llbracket N_{\tau} \rrbracket$  in the domain  $(\llbracket \tau \rrbracket \to \llbracket \tau \rrbracket) \to (\mathbb{N}_{\perp} \to \llbracket \tau \rrbracket).$  [4 marks]

(c) Recall that the contextual preorder  $\vdash M \leq_{\text{ctx}} N : \tau$  holds whenever M and N are closed PCF terms of type  $\tau$  and for all PCF contexts  $\mathcal{C}$  for which  $\mathcal{C}[M]$  and  $\mathcal{C}[N]$  are closed PCF terms of type  $\gamma \in \{nat, bool\}$  and for all values V of type  $\gamma$ , if  $\mathcal{C}[M] \Downarrow_{\gamma} V$  then  $\mathcal{C}[N] \Downarrow_{\gamma} V$ .

Say whether the following statements concerning the PCF terms in Part (b) are true or false and, respectively, either prove or disprove them:

(i) For all PCF types  $\tau$ ,  $\vdash M_{\tau} \leq_{ctx} N_{\tau} : (\tau \to \tau) \to (nat \to \tau).$ 

(Hint: Consider the case  $\tau = nat \rightarrow nat$ .) [6 marks]

(*ii*) For all PCF types  $\tau$ ,  $\vdash N_{\tau} \leq_{ctx} M_{\tau} : (\tau \to \tau) \to (nat \to \tau)$ .

(Hint: Recall that every PCF type is of the form  $\tau_1 \to (\cdots (\tau_{\ell} \to \gamma) \cdots)$ where  $\ell \in \mathbb{N}, \tau_i \ (1 \le i \le \ell)$  are types, and  $\gamma \in \{nat, bool\}$ .) [6 marks]