## COMPUTER SCIENCE TRIPOS Part IB – 2021 – Paper 4

## 3 Compiler Construction (tgg22)

(a) Suppose we are writing a compiler for an ML-like language and we want to employ the equation

as a left-to-right rewrite rule for optimisation. The symbol @ represents list append.

Discuss the merits of this idea. Is it always correct? If so, state clearly what assumptions you are making about @ and map. [5 marks]

(b) A compiler's front-end will often expand some syntactic constructs into lower-level representations. Consider the following fragment for the abstract syntax of a SLANG-like language.

```
type var = string
type exp =
                       (* concrete syntax *)
  (* abstract syntax *)
                            (* x *)
  | Var of var
  | Project of int * exp
                           (* proj i e *)
  | Tuple of exp list
                            (* (e1, e2, ..., en) *)
  | Let of var * exp * exp
                           (* let x = e1 in e2 *)
  | Apply of exp * exp
                            (* e1 e2 *)
  | Function of var * arg_pattern * exp
                                       (* fun f p = e *)
and arg_pattern =
                                  (* x *)
  | APvar of var
                                 (* (p1, p2, ... pn) *)
  | APtuple of arg_pattern list
```

This language has general projections for n-tuples so

proj 
$$i$$
  $(e_1, e_2, \cdots, e_k)$ 

will evaluate to  $v_i$ , the value of  $e_i$ . If i is not in the range between 1 and k there will be a run-time error.

In this language we can write functions with simple (possibly nested) patterns for function arguments:

fun f 
$$(a, b, (c, (d, e)) = b a$$

[continued ...]

Now suppose we want our front-end to eliminate such patterns. That is, we want to write a function of type

```
eliminate_tuple_patterns : exp -> exp
```

so that the resulting expression contains functions with patterns only of the form  $\texttt{APvar}\ \mathtt{x}$  for some (new) variable  $\mathtt{x}$ . For example, the code for  $\mathtt{f}$  above should be translated to a semantically equivalent expression of the form

fun f 
$$x = \dots$$

that contains only simple variable arguments (that is, only APvar patterns in the abstract syntax).

Your task is to write this function in OCaml. You can assume that you have a function for generating fresh variable strings.

[15 marks]