

7 Hoare Logic and Model Checking (jp622)

Consider a programming language that consists of commands  $C$  composed from assignments  $X := E$  (where  $X$  is a program variable, and  $E$  is an arithmetic expression), heap allocation  $X := \text{alloc}(E_1, \dots, E_n)$ , heap assignment  $[E_1] := E_2$ , heap dereference  $X := [E]$ , disposal of heap locations  $\text{dispose}(E)$ , the no-op  $\text{skip}$ , sequencing  $C_1; C_2$ , conditionals  $\text{if } B \text{ then } C_1 \text{ else } C_2$  (where  $B$  is a boolean expression), and loops  $\text{while } B \text{ do } C$ .  $\text{null}$  is 0

- (a) Explain informally what it means for a separation logic partial correctness triple  $\{P\} C \{Q\}$  to be valid. [3 marks]
- (b) Explain informally what it means in terms of the executions of  $C$  for the separation logic partial correctness triple  $\{\top\} C \{\perp\}$  to be valid. [2 marks]
- (c) Recall the list representation predicate  $\text{list}$ :

$$\text{list}(t, []) = (t = \text{null}) \quad \text{list}(t, h :: \alpha) = \exists y. ((t \mapsto h) * ((t + 1) \mapsto y) * \text{list}(y, \alpha))$$

We write  $[]$  for the empty mathematical list;  $h :: \alpha$  for the mathematical list the head of which is  $h$ , and the tail of which is  $\alpha$ ;  $\alpha ++ \beta$  for the concatenation of mathematical lists  $\alpha$  and  $\beta$ ;  $\alpha[i]$  for the  $i$ -th element of the list  $\alpha$ , starting at 0; and  $[k, \dots, n]$  for the ascending list of integers from  $k$  to  $n$ , including  $k$  and  $n$ . Give a proof outline, including a loop invariant, for the following triple:

$$\begin{aligned} & \{N = n \wedge N \geq 0\} \\ & X := \text{null}; \text{while } N > 0 \text{ do } (X := \text{alloc}(N, X); N := N - 1) \\ & \{\text{list}(X, [1, \dots, n])\} \end{aligned} \quad [4 \text{ marks}]$$

- (d) Also recall the partial list representation predicate  $\text{plist}$ :

$$\begin{aligned} \text{plist}(t, [], u) &= (t = u) \\ \text{plist}(t, h :: \alpha, u) &= \exists y. ((t \mapsto h) * ((t + 1) \mapsto y) * \text{plist}(y, \alpha, u)) \end{aligned}$$

Give a loop invariant for the following list sum triple:

$$\begin{aligned} & \{\text{list}(X, \alpha)\} \\ & Y := X; N := 0; \text{while } Y \neq \text{null} \text{ do } (M := [Y]; N := N + M; Y := [Y + 1]) \\ & \left\{ \text{list}(X, \alpha) \wedge N = \sum_{i=0}^{\text{length}(\alpha)-1} \alpha[i] \right\} \end{aligned} \quad [4 \text{ marks}]$$

- (e) Give a loop invariant for the following list concatenation triple:

$$\begin{aligned} & \{\text{list}(X, \alpha) * \text{list}(Y, \beta)\} \\ & \text{if } X = \text{null} \text{ then } Z := Y \text{ else} \\ & \left( \begin{array}{l} Z := X; U := Z; V := [Z + 1]; \\ \text{while } V \neq \text{null} \text{ do } (U := V; V := [V + 1]); \\ [U + 1] := Y \end{array} \right) \\ & \{\text{list}(Z, \alpha ++ \beta)\} \end{aligned} \quad [5 \text{ marks}]$$

- (f) Describe precisely a stack and a heap that satisfy  $\text{list}(X, [1, \dots, 3])$ . [2 marks]