

16 Types (AMP)

- (a) Existential types can be encoded in the Polymorphic Lambda Calculus (PLC) by defining  $\exists \alpha (\tau)$  to be  $\forall \beta ((\forall \alpha (\tau \rightarrow \beta)) \rightarrow \beta)$ , where  $\beta \neq \alpha$  and  $\beta$  does not occur free in the PLC type  $\tau$ . Give the following definitions, justifying the typings in each case:
- (i) A closed PLC term `pack` of type  $\forall \alpha (\tau \rightarrow \exists \alpha (\tau))$ . [5 marks]
- (ii) A PLC term `unpack`( $M, x, M', \tau'$ ) satisfying  $\Gamma \vdash \text{unpack}(M, x, M', \tau') : \tau'$  whenever  $\Gamma \vdash M : \exists \alpha (\tau)$  and  $\Gamma, x : \tau \vdash M' : \tau'$  hold, where  $x$  is not in the domain of  $\Gamma$  and  $\alpha$  does not occur free in  $\Gamma$  or  $\tau'$ . [5 marks]
- (b) If  $\Gamma \vdash M : \exists \alpha (\tau)$ ,  $\Gamma, x : \tau \vdash M' : \tau'$  and  $\Gamma \vdash N : \tau[\tau'/\alpha]$  hold, where  $x$  is not in the domain of  $\Gamma$  and  $\alpha$  does not occur free in  $\Gamma$  or  $\tau'$ , to what term does `unpack`(`(pack  $\tau'$   $N$ ),  $x, M', \tau'$ ) beta-reduce? [2 marks]`
- (c) For each PLC type  $\tau$ , let  $\neg \tau$  be the type  $\tau \rightarrow \forall \alpha (\alpha)$ . Give, with justification, closed PLC terms of the following types
- (i)  $\forall \alpha (\neg \tau) \rightarrow \neg \exists \alpha (\tau)$  [4 marks]
- (ii)  $\exists \alpha (\neg \tau) \rightarrow \neg \forall \alpha (\tau)$  [4 marks]