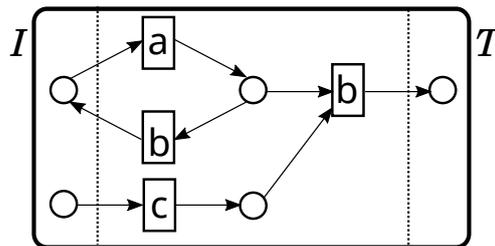


14 Topics in Concurrency (JMH)

- (a) Define the *token game* for basic Petri nets. [3 marks]
- (b) When is a basic Petri net *safe* from an initial marking? [2 marks]
- (c) An LB-net is a basic Petri net  $(B, E, pre, post)$  accompanied by
  - a labelling function  $\lambda : E \rightarrow Act$  from its events to a set of actions  $Act$
  - subsets of conditions  $I \subseteq B$  and  $T \subseteq B$ . The initial conditions  $I$  are marked when the process starts and the terminal conditions  $T$  are marked when the process has terminated.

LB-nets are drawn with labels inside events and boxes surrounding the initial and terminal conditions.

- (i) Draw the labelled transition system of the following LB-net. The initial state should correspond to the initial conditions  $I$  being marked and labels on the transitions should correspond to actions, not events. [4 marks]



- (ii) Ignoring the particular sets that states represent, is there an LB-net with an *injective* labelling function  $\lambda$  that gives rise to the same labelled transition system? Justify your answer briefly. [2 marks]
- (iii) A simple process language has the following syntax.

$$p ::= \alpha \mid p + p' \mid p \parallel p' \mid p; p'$$

where  $\alpha \in Act$ . As in CCS,  $+$  represents the nondeterministic sum of processes and  $\parallel$  represents the parallel composition. The process  $p; p'$  represents the sequential composition of  $p$  and  $p'$ .

Draw diagrams to describe the inductive definition of an LB-net semantics for this fragment. [7 marks]

- (iv) An iteration operator  $p^*$  is proposed with LB-net semantics such that its sets of initial and terminal conditions are equal:  $I = T$ . Discuss briefly how this affects the semantics you gave in part (c)(iii). [2 marks]