

## 2008 Paper 9 Question 16

### Optimising Compilers

Consider the language

$$e ::= x \mid \lambda x.e \mid e_1 e_2 \mid e_1; e_2 \mid \xi?x.e \mid \xi!e_1.e_2 \mid \text{if } e_1 \text{ then } e_2 \text{ else } e_3.$$

in which  $\xi$  represents a communication channel (from a fixed set), and the forms  $e_1; e_2$ ,  $\xi?x.e$  and  $\xi!e_1.e_2$  respectively represent sequencing, reading from a channel (binding  $x$ ) and writing to a channel.

- (a) Construct an *effect system* for the above language where effects,  $F$ , are represented as *sets* of actions of the form  $\xi?$  or  $\xi!$  representing side-effects of input from or output to  $\xi$ . Explain the two principal occurrences of effects in the judgement form of your system. [8 marks]
- (b) Assess the *safety* of your analysis making clear any respects in which execution behaviour may fail to match your analysis. [2 marks]
- (c) Let us say a general program analysis framework is *any-path*-like (as opposed to *all-path*-like) if the analyses of  $\text{if } e_1 \text{ then } e_2 \text{ else } e_3$  and  $e_1; e_2; e_3$  coincide. Is your effect system any-path-like? Justify your answer. [2 marks]
- (d) Augment the above language with constructs

$$e ::= \text{letchan } \xi \text{ in } e \mid \text{parsum}(e_1, e_2)$$

which allow a *local channel* to be created, and also inter-thread communication ( $e_1$  and  $e_2$  are evaluated in parallel and their sum returned when both have completed). Extend your effect system to the augmented language, noting that reads and writes to local channels are *not* to be reflected in the overall effect of a **letchan**. [6 marks]

- (e) Suggest an alternative data structure for  $F$  that might enable effects of the form “after getting two inputs from channel  $\xi_1$  or getting one input from channel  $\xi_2$  then an output is written to channel  $\zeta$ ” to be represented. [A modified effect system is not required.] [2 marks]