2005 Paper 5 Question 6

Compiler Construction

(*a*) Explain how a parse tree representing an expression can (*i*) be converted into stack-oriented intermediate code and then (*ii*) be translated into simple machine code for a register-oriented architecture (e.g. ARM or IA32) on an instruction-by-instruction basis. Also indicate how this code might be improved to remove push–pop pairs introduced by (*ii*). Your answer need only consider expression forms encountered in the expression:

```
h(a, g(b), c) * 3 + d
```

[12 marks]

(*b*) In Java, expressions are evaluated strictly left-to-right. Consider compiling the function `f` in the following Java class definition:

```
class A
{
    static int a,b;
    void f() { ... <<C>> ... }
    int g(int x) { ... a++; ... }
};
```

Indicate what *both* the intermediate code *and* (improved as above) target code might be for `<<C>>` for the cases where `<<C>>` is:

(*i*) `b = g(7) + a;`

(*ii*) `b = a + g(7);`

(*iii*) `b = (-g(7)) + a;`

(*iv*) `b = a - g(7);`

Comment on any inherent differences in efficiency at both the intermediate code and target code levels.

[8 marks]