

## 1998 Paper 8 Question 15

### Communicating Automata and Pi Calculus

A process is *deterministic* if, for any state  $P$  and any action  $a$ , there is at most one transition  $P \xrightarrow{a} P'$ . If  $s = a_1 \cdots a_n$  is a sequence of actions, then  $P \xrightarrow{s} P'$  means  $P \xrightarrow{a_1} \cdots \xrightarrow{a_n} P'$ .

By drawing a transition graph, define a deterministic sequential process  $Q$  which uses the four actions  $\text{in}_1, \text{out}_1, \text{in}_2, \text{out}_2$  such that, for every action sequence  $s$ ,  $Q \xrightarrow{s} Q'$  for some  $Q'$  if and only if  $s$  satisfies the four conditions

$$\begin{array}{ll} (a) & 0 \leq \#_s \text{in}_1 - \#_s \text{out}_1 \leq 1 \\ (b) & 0 \leq \#_s \text{in}_2 - \#_s \text{out}_2 \leq 1 \\ (c) & 0 \leq \#_s \text{in}_1 - \#_s \text{in}_2 \leq 1 \\ (d) & 0 \leq \#_s \text{out}_1 - \#_s \text{out}_2 \leq 1 \end{array}$$

where  $\#_s a$  is defined to be the number of times  $a$  occurs in  $s$ . For example the sequence  $s = \text{in}_1 \text{in}_2 \text{out}_1$  satisfies the conditions, but  $s = \text{in}_1 \text{in}_2 \text{out}_2$  violates (d). Also write down the defining equations of  $Q$ .

[Hint:  $Q$  need have no more than eight states.] [8 marks]

Define *weak bisimulation* and *weak equivalence* ( $\approx$ ), and explain their advantages over the strong versions. [4 marks]

With the help of one or more extra actions  $\vec{a}$ , define two communicating processes  $R_1$  and  $R_2$  with no more than four states each, such that

$$Q \approx \text{new } \vec{a}(R_1 \mid R_2)$$

and verify this weak equivalence. [8 marks]