

## 1998 Paper 7 Question 10

### Types

Each natural number  $n \in \mathbb{N}$  can be encoded in the polymorphic lambda calculus (PLC) by the beta-normal form  $Num_n \stackrel{\text{def}}{=} \Lambda \alpha (\lambda x : \alpha (\lambda f : \alpha \rightarrow \alpha (f^n x)))$ , where the expression  $f^n x$  is an abbreviation for the PLC expression inductively defined by

$$f^0 x \stackrel{\text{def}}{=} x \quad \text{and} \quad f^{n+1} x \stackrel{\text{def}}{=} f (f^n x).$$

For which PLC type  $nat$  does  $\vdash Num_n : nat$  hold? [3 marks]

Say that a function  $\phi : \mathbb{N} \rightarrow \mathbb{N}$  is *PLC-representable* if there is a PLC expression  $F$  such that the following typing and beta-conversions hold:

$$\vdash F : nat \rightarrow nat \quad \text{and} \quad F Num_n =_{\beta} Num_{\phi(n)} \quad (\text{all } n \in \mathbb{N}).$$

Show that the successor function,  $s(n) \stackrel{\text{def}}{=} n + 1$ , is PLC-representable. [5 marks]

Given a PLC-representable function  $\phi$  and a number  $a \in \mathbb{N}$ , show that the function  $\psi$  inductively defined by

$$\psi(0) \stackrel{\text{def}}{=} a \quad \text{and} \quad \psi(n+1) \stackrel{\text{def}}{=} \phi(\psi(n))$$

is PLC-representable. [8 marks]

Is every function  $\mathbb{N} \rightarrow \mathbb{N}$  PLC-representable? Justify your answer. [4 marks]