Randomised Algorithms

Lecture 3: Concentration Inequalities, Application to Quick-Sort, Extensions

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Outline

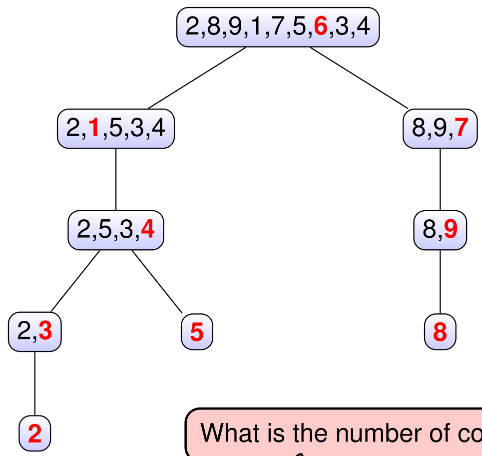
QuickSort

QUICKSORT (Input $A[1], A[2], \ldots, A[n]$)
Pick an element from the array, the so-called pivot
If $|A| = 0$ or $|A| = 1$ **then**
    return A
else
    Create two subarrays A_1 and A_2 (without the pivot) such that:
        A_1 contains the elements that are smaller than the pivot
        A_2 contains the elements that are greater (or equal) than the pivot
    QUICKSORT(A_1)
    QUICKSORT(A_2)
    return A

- Example: Let $A = (2, 8, 9, 1, 7, 5, 6, 3, 4)$ with $A[7] = 6$ as pivot.
  $\Rightarrow A_1 = (2, 1, 5, 3, 4)$ and $A_2 = (8, 9, 7)$
- Worst-Case Complexity (number of comparisons) is $\Theta(n^2)$, while Average-Case Complexity is $O(n \log n)$.

We will now give a proof of this “well-known” result!

QuickSort: How to Count Comparisons



What is the number of comparisons?

Note that the number of comparison by QUICKSORT is equivalent to the sum of the depths of all nodes in the tree (why?). In this case:

$$0 + 1 + 1 + 2 + 2 + 3 + 3 + 3 + 4 = 19.$$

Randomised QuickSort: Analysis (1/4)

How to pick a good pivot? We don't, **just pick one at random.**

This should be your standard answer in this course ☺

Let us analyse QuickSort with random pivots.

1. Assume A consists of n different numbers, w.l.o.g., $\{1, 2, \ldots, n\}$
2. Let H_i be the deepest level where element i appears in the tree. Then the number of comparison is $H = \sum_{i=1}^{n} H_i$
3. We will prove that there exists $C > 0$ such that

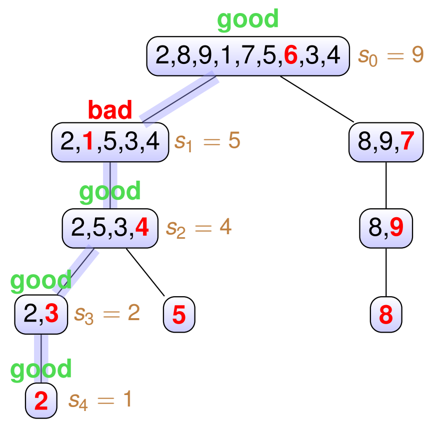
$$\mathbf{P}[H \leq Cn \log n] \geq 1 - n^{-1}.$$

4. Actually, we will prove sth slightly stronger:

$$\mathbf{P}\left[\bigcap_{i=1}^{n} \{H_i \leq C \log n\}\right] \geq 1 - n^{-1}.$$

Randomised QuickSort: Analysis (2/4)

- Let P be a path from the root to the deepest level of some element
  - A node in P is called **good** if the corresponding pivot partitions the array into two subarrays each of size at most $2/3$ of the previous one
  - otherwise, the node is **bad**
- Further let s_t be the size of the array at level t in P.



- Element 2: $(2, 8, 9, 1, 7, 5, \mathbf{6}, 3, 4) \rightarrow (2, \mathbf{1}, 5, 3, 4) \rightarrow (2, 5, 3, \mathbf{4}) \rightarrow (2, \mathbf{3}) \rightarrow (\mathbf{2})$

Randomised QuickSort: Analysis (3/4)

- Consider now any element $i \in \{1, 2, \ldots, n\}$ and construct the path $P = P(i)$ one level by one
- For P to proceed from level k to $k+1$, the condition $s_k > 1$ is necessary

How far could such a path P possibly run until we have $s_k = 1$?

- We start with $s_0 = n$
- **First Case**, **good** node: $s_{k+1} \leq \frac{2}{3} \cdot s_k$.
- **Second Case**, **bad** node: $s_{k+1} \leq s_k$.

This even holds always, i.e., deterministically!

$\Rightarrow$ There are at most $T = \frac{\log n}{\log(3/2)} < 3 \log n$ many **good** nodes on any path P.

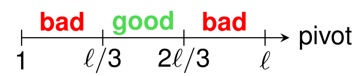
- Assume $|P| \geq C \log n$ for $C := 24$

  $\Rightarrow$ number of **bad** vertices in the first $24 \log n$ levels is more than $21 \log n$.

Let us now upper bound the probability that this "bad event" happens!

Randomised QuickSort: Analysis (4/4)

- Consider the first $24 \log n$ vertices of P to the deepest level of element i.
- For any level $j \in \{0, 1, \ldots, 24 \log n - 1\}$, define an indicator variable X_j:
  - $X_j = 1$ if the node at level j is **bad**,
  - $X_j = 0$ if the node at level j is **good**.



- $\mathbf{P}[X_j = 1 \mid X_0 = x_0, \ldots, X_{j-1} = x_{j-1}] \leq \frac{2}{3}$
- $X := \sum_{j=0}^{24 \log n - 1} X_j$ satisfies relaxed independence assumption (Lecture 2)

Question: Edge Case: What if the path P does not reach level j?

Randomised QuickSort: Analysis (4/4)

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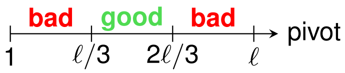

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Question: Edge Case: What if the path P does not reach level j?

Answer: We can then simply define X_j as 0 (deterministically).

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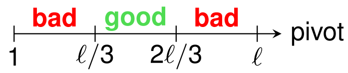

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We can now apply the "nicer" **Chernoff Bound**!

- We have $\mathbf{E}[\, X \,] \leq (2/3) \cdot 24 \log n = 16 \log n$
- Then, by the "nicer" Chernoff Bounds $\mathbf{P}[\, X \geq \mathbf{E}[\, X \,] + t \,] \leq e^{-2t^2/n}$

$$\mathbf{P}[\, X > 21 \log n \,] \leq \mathbf{P}[\, X > \mathbf{E}[\, X \,] + 5 \log n \,]$$

- Hence P has more than $24 \log n$ nodes with probability at most n^{-2}.
- As there are in total n paths, by the union bound, the probability that at least one of them has more than $24 \log n$ nodes is at most n^{-1}.
- This implies $\mathbf{P}\left[\, \bigcap_{i=1}^{n} \{H_i \leq 24 \log n\} \,\right] \geq 1 - n^{-1}$, as needed. □

Randomised QuickSort: Final Remarks

- Well-known: any comparison-based sorting algorithm needs $\Omega(n \log n)$
- A classical result: expected number of comparison of randomised QUICKSORT is $2n \log n + O(n)$ (see, e.g., book by Mitzenmacher & Upfal)

Exercise: *[Ex 2-3.6]* Our upper bound of $O(n \log n)$ whp also immediately implies a $O(n \log n)$ bound on the expected number of comparisons!

- It is possible to deterministically find the best pivot element that divides the array into two subarrays of the same size.
- The latter requires to compute the median of the array in linear time, which is not easy...
- The presented randomised algorithm for QUICKSORT is much easier to implement!

Outline

Hoeffding's Extension

- Besides **sums of independent Bernoulli** random variables, sums of independent and bounded random variables are very frequent in applications.
- Unfortunately the distribution of the X_i may be unknown or hard to compute, thus it will be hard to compute the moment-generating function.
- Hoeffding's Lemma helps us here:

You can always consider $X' = X - \mathbf{E}[X]$

Hoeffding's Extension Lemma

Let X be a random variable with mean 0 such that $a \leq X \leq b$. Then for all $\lambda \in \mathbb{R}$,

$$\mathbf{E}\left[e^{\lambda X}\right] \leq \exp\left(\frac{(b-a)^2\lambda^2}{8}\right)$$

We omit the proof of this lemma!

Hoeffding Bounds

Hoeffding's Inequality

Let $X_1, \ldots, X_n$ be independent random variable with mean μ_i such that $a_i \leq X_i \leq b_i$. Let $X = X_1 + \ldots + X_n$, and let $\mu = \mathbf{E}[X] = \sum_{i=1}^n \mu_i$. Then for any $t > 0$

$$\mathbf{P}[X \geq \mu + t] \leq \exp\left(-\frac{2t^2}{\sum_{i=1}^n (b_i - a_i)^2}\right),$$

and

$$\mathbf{P}[X \leq \mu - t] \leq \exp\left(-\frac{2t^2}{\sum_{i=1}^n (b_i - a_i)^2}\right).$$

Proof Outline (skipped):

- Let $X_i' = X_i - \mu_i$ and $X' = X_1' + \ldots + X_n'$, then $\mathbf{P}[X \geq \mu + t] = \mathbf{P}[X' \geq t]$
- $\mathbf{P}[X' \geq t] \leq e^{-\lambda t} \prod_{i=1}^n \mathbf{E}\left[e^{\lambda X_i'}\right] \leq \exp\left[-\lambda t + \frac{\lambda^2}{8}\sum_{i=1}^n (b_i - a_i)^2\right]$
- Choose $\lambda = \frac{4t}{\sum_{i=1}^n (b_i - a_i)^2}$ to get the result.

This is not magic! you just need to optimise λ!

Method of Bounded Differences

Framework

Suppose, we have independent random variables $X_1, \ldots, X_n$. We want to study the random variable:

$$f(X_1, \ldots, X_n)$$

Some examples:

1. $X = X_1 + \ldots + X_n$ (our setting earlier)
2. In balls into bins, X_i indicates where ball i is allocated, and $f(X_1, \ldots, X_m)$ is the number of empty bins
3. In a randomly generated graph, X_i indicates if the i-th edge is present and $f(X_1, \ldots, X_m)$ represents the number of connected components of G

In all those cases (and more) we can easily prove concentration of $f(X_1, \ldots, X_n)$ around its mean by the so-called **Method of Bounded Differences.**

Method of Bounded Differences

A function f is called Lipschitz with parameters $\mathbf{c} = (c_1, \ldots, c_n)$ if for all $i = 1, 2, \ldots, n$,

$$|f(x_1, x_2, \ldots, x_{i-1}, x_i, x_{i+1}, \ldots, x_n) - f(x_1, x_2, \ldots, x_{i-1}, \widetilde{x_i}, x_{i+1}, \ldots, x_n)| \leq c_i,$$

where x_i and $\widetilde{x_i}$ are in the domain of the i-th coordinate.

McDiarmid's inequality

Let $X_1, \ldots, X_n$ be independent random variables. Let f be Lipschitz with parameters $\mathbf{c} = (c_1, \ldots, c_n)$. Let $X = f(X_1, \ldots, X_n)$. Then for any $t > 0$,

$$\mathbf{P}[X \geq \mu + t] \leq \exp\left(-\frac{2t^2}{\sum_{i=1}^n c_i^2}\right),$$

and

$$\mathbf{P}[X \leq \mu - t] \leq \exp\left(-\frac{2t^2}{\sum_{i=1}^n c_i^2}\right).$$

- Notice the similarity with Hoeffding's inequality! *[Exercise 2/3.14]*
- The proof is omitted here (it requires the concept of martingales).

Outline

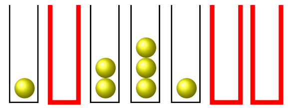
Application 2: Randomised QuickSort

Extensions of Chernoff Bounds

Applications of Method of Bounded Differences

Appendix: More on Moment Generating Functions (non-examinable)

Application 3: Balls into Bins (again...)

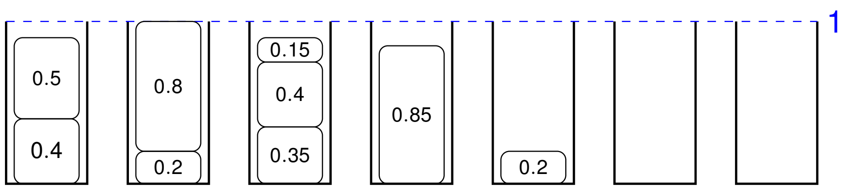


- Consider again m balls assigned uniformly at random into n bins.
- Enumerate the balls from 1 to m. Ball i is assigned to a random bin X_i
- Let Z be the number of empty bins (after assigning the m balls)
- $Z = Z(X_1, \ldots, X_m)$ and Z is Lipschitz with $\mathbf{c} = (1, \ldots, 1)$
  (If we move one ball to another bin, number of empty bins changes by ≤ 1.)
- By McDiarmid's inequality, for any $t \geq 0$,

$$\mathbf{P}[\,|Z - \mathbf{E}[\,Z\,]\,| > t\,] \leq 2 \cdot e^{-2t^2/m}.$$

This is a decent bound, but for some values of m it is far from tight and stronger bounds are possible through a refined analysis.

Application 4: Bin Packing



- We are given n items of sizes in the unit interval $[0, 1]$
- We want to pack those items into the fewest number of unit-capacity bins
- Suppose the item sizes X_i are independent random variables in $[0, 1]$

- Let $B = B(X_1, \ldots, X_n)$ be the optimal number of bins
- The Lipschitz conditions holds with $\boldsymbol{c} = (1, \ldots, 1)$. **Why?**
- Therefore

$$\mathbf{P}[\,|B - \mathbf{E}[\,B\,]\,| \geq t\,] \leq 2 \cdot e^{-2t^2/n}.$$

This is a typical example where proving concentration is much easier than calculating (or estimating) the expectation!

Outline

Application 2: Randomised QuickSort

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Applications of Method of Bounded Differences

Appendix: More on Moment Generating Functions (non-examinable)

Moment Generating Functions (non-examinable)

Moment-Generating Function

The moment-generating function of a random variable X is

$$M_X(t) = \mathbf{E}\left[e^{tX}\right], \qquad \text{where } t \in \mathbb{R}.$$

Using power series of e and differentiating shows that $M_X(t)$ encapsulates all moments of X.

Lemma

1. If X and Y are two r.v.'s with $M_X(t) = M_Y(t)$ for all $t \in (-\delta, +\delta)$ for some $\delta > 0$, then the distributions X and Y are identical.
2. If X and Y are independent random variables, then

$$M_{X+Y}(t) = M_X(t) \cdot M_Y(t).$$

Proof of 2:

$$M_{X+Y}(t) = \mathbf{E}\left[e^{t(X+Y)}\right] = \mathbf{E}\left[e^{tX} \cdot e^{tY}\right] \overset{(!)}{=} \mathbf{E}\left[e^{tX}\right] \cdot \mathbf{E}\left[e^{tY}\right] = M_X(t)M_Y(t) \quad \square$$