• Kleene's fixed point theorem: every continuous function  $f: D \to D$  on a domain D has a least (pre-)fixed point given by  $\bigsqcup_{n \in \mathbb{N}} f^n(\bot_D)$ 

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• We need to construct domains and continuous functions!

- + Flat domains: "base" cases ( $\mathbb{N}_{\perp}$ , but also  $\mathbb{B}_{\perp}$ )
- Products of domains are domains, everything is componentwise

Let  $D_1$  and  $D_2$  be cpos. The projections

$$\pi_{1}: \begin{array}{ccc} D_{1} \times D_{2} & \rightarrow & D_{1} \\ (d_{1}, d_{2}) & \mapsto & d_{1} \end{array} \qquad \begin{array}{ccc} \pi_{2}: & D_{1} \times D_{2} & \rightarrow & D_{2} \\ (d_{1}, d_{2}) & \mapsto & d_{1} \end{array}$$
are continuous functions.
$$\Pi_{1}\left(\left(\begin{array}{c} \left(\begin{array}{c} d_{1} & d_{2} \end{array}\right)\right) - \Pi_{1}\left(\begin{array}{c} d_{1} & d_{1} & d_{2} \end{array}\right) \\ \prod_{n} d_{n} & \prod_{n} d_{n} \end{array}\right) - \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) - \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \\ d_{n} & d_{n} \end{array}\right) + \prod_{n} \left(\begin{array}{c} d_{n} & d_{n} \end{array}\right) + \prod_{n}$$

Let  $D_1$  and  $D_2$  be cpos. The projections

are continuous functions.

If  $f_1: D \to D_1$  and  $f_2: D \to D_2$  are continuous functions from a cpo D, then the pairing function

$$\begin{array}{cccc} \langle f_1, f_2 \rangle : & D & \to & D_1 \times D_2 \\ & d & \mapsto & (f_1(d), f_2(d)) \end{array}$$

is continuous.

#### The **conditional** function

$$\begin{array}{rcl} \text{if} : & \mathbb{B}_{\perp} \times (D \times D) & \to & D \\ & & (x,d) & \mapsto & \begin{cases} \pi_1(d) & \text{if } x = \text{true} \\ \pi_2(d) & \text{if } x = \text{false} \\ \perp_D & \text{if } x = \perp \end{cases}$$

is continuous.

 $(b_{M} d_{1} d_{2}, n) c (B_{1} \times (D_{1} \times D_{2}))_{m \in \mathbb{N}}$  $U_{n}(b_{n}, d_{n}, d_{e_{n}}) = (L_{n}b_{n}, (I_{n}, d_{n}, L_{n}, d_{e_{n}})$ if  $(U_n(D_n, t_{1,n}, d_{2,n})) = U_n if (b_n, d_{1,n}, d_{2,n})$ is enough to look at  $b_n$  constant = b b = Lthen if (time Ud, n (Iden)= Ud, n = Uif (time, 1, n, 1, n) · b= true anilan , W = Jalse

Given a set I, suppose that for each  $i \in I$  we are given a set  $X_i$ . The (cartesian) product of the  $X_i$  is

 $\prod_{i\in I} X_i$ 

Two ways to see it:

• tuples:  $(\ldots, x_i, \ldots)_{i \in I}$  such that  $x_i \in X_i$ ;

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Special case:  $\prod_{i\in\mathbb{B}}D_i$  corresponds to  $D_{ ext{true}} imes D_{ ext{false}}.$ Projections (for any  $i\in I$ ):

$$\pi_i: \left(\prod_{i\in I} X_i\right) \to X_i$$

Given a set I, suppose that for each  $i \in I$  we are given a cpo  $(D_i, \sqsubseteq_i)$ . The **product** of this whole family of cpos has

• underlying set equal to  $\prod_{i \in I} D_i$ ;

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- underlying set equal to  $\prod_{i \in I} D_i$ ;
- componentwise order

$$p \sqsubseteq p' \stackrel{\text{def}}{\Leftrightarrow} \forall i \in I. \ p_i \sqsubseteq_i p'_i.$$

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*I*-indexed products of cpos (domains) are cpos (domains), and projections are continuous.

# CONSTRUCTIONS ON DOMAINS

Given two cpos  $(D, \sqsubseteq_D)$  and  $(E, \sqsubseteq_E)$ , the function cpo  $(D \to E, \sqsubseteq)$  has underlying set

 $\{f: D \to E \mid \text{ is a continuous function}\}$ 

equipped with the pointwise order:

$$f \sqsubseteq f' \stackrel{\text{def}}{\Leftrightarrow} \forall d \in D. \ f(d) \sqsubseteq_E f'(d).$$

## CPO/DOMAIN OF CONTINUOUS FUNCTIONS

Given two cpos  $(D, \sqsubseteq_D)$  and  $(E, \sqsubseteq_E)$ , the function cpo  $(D \to E, \sqsubseteq)$  has underlying set  $\{f : D \to E \mid \text{ is a continuous function}\}$ 

equipped with the pointwise order:

$$f \sqsubseteq f' \stackrel{\text{def}}{\Leftrightarrow} \forall d \in D. \ f(d) \sqsubseteq_E f'(d).$$

$$\frac{f \sqsubseteq_{D \to E} g \qquad x \sqsubseteq_D y}{f(x) \sqsubseteq_E g(y)}$$

Given two cpos  $(D, \sqsubseteq_D)$  and  $(E, \sqsubseteq_E)$ , the function cpo  $(D \to E, \sqsubseteq)$  has underlying set

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equipped with the pointwise order:

$$f \sqsubseteq f' \stackrel{\text{def}}{\Leftrightarrow} \forall d \in D. f(d) \sqsubseteq_E f'(d).$$

Argumentwise least elements and lubs:

$$\perp_{D \to E}(d) = \perp_E \qquad \qquad \left(\bigsqcup_{n \ge 0} f_n\right)(d) = \bigsqcup_{n \ge 0} f_n(d)$$

 $\in D \rightarrow E$ Take  $f_0 = f_1 = f_2 = \dots$  $f_{a}: d \mapsto \mathcal{U}(f_{a}(d))$ John ded (H)E- f(d) E... So U (full) exists it Dan uper hound for fu (d) So for I for it for is hourd then In(d) Ef(d) for all d Arostime J] Fr  $D_{\Theta} = \bigcup_{n \in \mathcal{I}} (d) = \bigcup_{n \in \mathcal{I}} (d)$  $D_{0}$   $f_{0}(d)$  Ef(d)Do fried

Remains. La iscontinuous do Ed, Edm -- chain in D  $\int_{M} (Udm) = U(f_m(Udm)) = U(Uf_m(dm))$ 

**Evaluation**, currying  $(f : (D' \times D) \rightarrow E)$  and composition

eval: 
$$(D \to E) \times D \to E$$
  
 $(f, d) \mapsto f(d)$ 

$$\operatorname{cur}(f): D' \to (D \to E)$$
$$d' \mapsto \lambda d \in D. f(d', d)$$

$$\circ: \ \begin{pmatrix} (E \to F) \times (D \to E) \end{pmatrix} \longrightarrow \ (D \to F) \\ (f,g) \mapsto \lambda d \in D. \ g(f(d)) \end{pmatrix}$$

are all well-defined and continuous.

 $f: (D \times D) \rightarrow E$  then  $cur(1): D \rightarrow (D' \rightarrow E)$ is intimuous  $\operatorname{Cur}: (D \times D') \rightarrow E) \rightarrow (D \rightarrow (D' \rightarrow E))$  is continuous?

eval: 
$$D^{\lambda}(D \rightarrow E) \rightarrow E$$
  
Take  $d_{0} \in d_{1} \subseteq \dots d_{n} \in D$   $\int_{0} \subseteq f_{4} \subseteq \dots \in D \rightarrow E$   
 $eval((\underbrace{U}(\underbrace{d_{m}}, d_{n}))) = eval((\underbrace{U}f_{m}, \underbrace{U}d_{n})) = eval((\underbrace{U}f_{m}, \underbrace{U}d_{n}))$   
 $= (\underbrace{U}f_{m})(\underbrace{U}d_{n})$   
 $= \underbrace{U}(\underbrace{f_{m}}(\underbrace{U}d_{n}))$   
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### fix: $(D \rightarrow D) \rightarrow D$

is continuous.

 $(abe f_0 \subseteq f_1 \subseteq \dots = f_n \in D \Rightarrow D$ First slow fix monstone: JEE (D->1) g(fix(e)) I e (fi(e)) be JEE I fix(e) so fix(e) is a prefixed point of f  $j(x(1) \subseteq f(e)$ We wantpix(Infm) = In fix(fm) it is enough to show fix(Unfn) = Unfix(fm) fix(Unfn) = Unfix(fm) evough: Unfix(fm) is a pre-fixed point for Unfm

 $(U, J_{\mathcal{M}})$  $(U, J_{\mathcal{M}})$  $(U, J_{\mathcal{M}})$ = $(U, U, J_{\mathcal{M}})$  $(for (f_{\mathcal{M}}))$  $(for (f_{\mathcal{M}}))$  $= \bigcup_{k} f_{k}(f_{k}(f_{k}))$ = U fix(fh) U fixlight is a fixed with of Ufm

# **CONSTRUCTIONS ON DOMAINS**

BACK TO THE INTRODUCTION

$$\llbracket while X > 0 \text{ do } (Y \coloneqq X * Y; X \coloneqq X - 1) \rrbracket$$

is a fixed point of the following  $F: D \rightarrow D$ , where D is (State  $\rightarrow$  State):

$$F(w)([X \mapsto x, Y \mapsto y]) = \begin{cases} [X \mapsto x, Y \mapsto y] & \text{if } x \le 0\\ w([X \mapsto x - 1, Y \mapsto x \cdot y]) & \text{if } x > 0. \end{cases}$$

$$\llbracket \texttt{while } X > 0 \texttt{ do } (Y \coloneqq X * Y; X \coloneqq X - 1) \rrbracket$$

is a fixed point of the following  $F: D \rightarrow D$ , where D is  $(State_{\perp} \rightarrow State_{\perp})$ :

$$F(w)([X \mapsto x, Y \mapsto y]) = \begin{cases} [X \mapsto x, Y \mapsto y] & \text{if } x \le 0\\ w([X \mapsto x - 1, Y \mapsto x \cdot y]) & \text{if } x > 0. \end{cases}$$
$$F(\bot) = \bot$$

 $State_{\perp} \rightarrow State_{\perp}$  is a domain!

Kleene's fixed point theorem:

$$w_{\infty} = \bigsqcup_{i \in \mathbb{N}} F^n(\bot)$$

is the least fixed point of F, and in particular a fixed point.

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$$w_{\infty} = \bigsqcup_{i \in \mathbb{N}} F^n(\bot)$$

is the least fixed point of F, and in particular a fixed point.

We can compute explicitly

$$w_{\infty}[X \mapsto x, Y \mapsto y] = \begin{cases} [X \mapsto x, Y \mapsto y] & \text{if } x < 0\\ [X \mapsto 0, Y \mapsto (x!) \cdot y] & \text{if } x \ge 0 \end{cases}$$

And **check** this agrees with the operational semantics.