

BIG

UNIONS and INTERSECTIONS

Sets and logic

$\mathcal{P}(U)$	$\{ \text{false}, \text{true} \}$
\emptyset	false
U	true
\cup	\vee
\cap	\wedge
$(\cdot)^c$	$\neg(\cdot)$
\cup \cap	\exists \forall

Example: Big union

- $\mathcal{T}_0 =_{\text{def}} \{ T \subseteq [5] \mid \begin{array}{l} \text{the sum of the elements} \\ \text{of } T \text{ is less than or equal } 2 \end{array} \}$
 $= \{ \emptyset, \{0\}, \{1\}, \{0,1\}, \{0,2\} \}$

- $\bigcup \mathcal{T}_0$ is the union of the sets in \mathcal{T}_0

$$n \in \bigcup \mathcal{T}_0 \Leftrightarrow \exists T \in \mathcal{T}_0. n \in T$$

$$\bigcup \mathcal{T}_0 = \{0, 1, 2\}$$

Big unions

Definition 90 Let \mathcal{U} be a set. For a collection of sets $\mathcal{F} \in \mathcal{P}(\mathcal{P}(\mathcal{U}))$, we let the big union (relative to \mathcal{U}) be defined as

$$\bigcup \mathcal{F} = \{ x \in \mathcal{U} \mid \exists A \in \mathcal{F}. x \in A \} \in \mathcal{P}(\mathcal{U}) \quad .$$

Examples:

- $\bigcup (\mathcal{P}(U)) = \{x \in U \mid \exists S \in \mathcal{P}(U). x \in S\}$
 $= \{x \in U \mid \underline{\text{true}}\} = U$

- $\bigcup \emptyset = \{x \in U \mid \exists S \in \emptyset, x \in S\}$
 $= \{x \in U \mid \underline{\text{false}}\} = \emptyset$

ASSOCIATIVITY
(idea / intuition)

$$F \subseteq \mathcal{P}(\mathcal{P}(U))$$

$$F = \{ \dots, \underset{\parallel}{A}, \underset{\parallel}{B}, \dots \}$$

$$\{ \dots, B, B', \dots \}$$

$$\{ \dots, A, A', \dots \}$$

- $\cup F = \dots \cup A \cup B \cup \dots$

$$= \{ \dots, \dots A, A', \dots, \dots B, B', \dots, \dots \}$$

$$\cup(\cup F) = (\dots \cup A \cup A' \cup \dots \cup B \cup B' \cup \dots)$$

- $\cup \{ \dots, \cup A, \cup B, \dots \}$

$$= \dots \cup (\dots \cup A \cup A' \cup \dots) \cup (\dots \cup B \cup B' \cup \dots) \cup \dots$$

Proposition 91 For all $\mathcal{F} \in \mathcal{P}(\mathcal{P}(\mathcal{P}(\mathcal{U})))$,

$$\bigcup (\bigcup \mathcal{F}) = \bigcup \left\{ \bigcup \mathcal{A} \in \mathcal{P}(\mathcal{U}) \mid \mathcal{A} \in \mathcal{F} \right\} \in \mathcal{P}(\mathcal{U}) .$$

PROOF:

NB(1): pattern-matching notation for
 $\{ X \in \mathcal{P}(\mathcal{U}) \mid \exists \mathcal{A} \in \mathcal{F}. X = \bigcup \mathcal{A} \}$

NB(2): (Type-checking) as $\mathcal{F} \in \mathcal{P}(\mathcal{P}(\mathcal{P}(\mathcal{U})))$
we have $\bigcup \mathcal{F} \in \mathcal{P}(\mathcal{P}(\mathcal{U}))$ and then

$$\bigcup (\bigcup \mathcal{F}) \in \mathcal{P}(\mathcal{U})$$

$$U(U\mathcal{F})^{\text{RTP}} = U \{ U\mathcal{A} \in \mathcal{P}(U) \mid \mathcal{A} \in \mathcal{F} \}$$

$\forall x \in U.$

$$x \in U(U\mathcal{F}) \Leftrightarrow x \in U \{ U\mathcal{A} \in \mathcal{P}(U) \mid \mathcal{A} \in \mathcal{F} \}.$$

- $$\begin{aligned}
 x \in U(U\mathcal{F}) &\Leftrightarrow \exists S \in U\mathcal{F}. x \in S \\
 &\Leftrightarrow \exists S. S \in U\mathcal{F} \wedge x \in S \\
 &\Leftrightarrow \exists S. \exists \mathcal{A} \in \mathcal{F}. S \in \mathcal{A} \wedge x \in S
 \end{aligned}$$

- $$x \in U \{ U\mathcal{A} \in \mathcal{P}(U) \mid \mathcal{A} \in \mathcal{F} \}$$

$$\Leftrightarrow \exists \mathcal{A} \in \mathcal{F}. x \in U\mathcal{A}.$$

$$\Leftrightarrow \exists \mathcal{A} \in \mathcal{F}. \exists S \in \mathcal{A}. x \in S.$$



PROOF: For $x \in U$, we show:

$$x \in \bigcup (UF) \Leftrightarrow x \in \bigcup \{X \in \mathcal{P}(U) \mid \exists A \in \mathcal{F}. X = \bigcup A\}$$

On the one hand,

$$x \in \bigcup (UF) \Leftrightarrow \exists S \in UF. x \in S$$

$$\Leftrightarrow \exists A \in \mathcal{F}. \exists S \in A. x \in S$$

On the other hand,

$$x \in \bigcup \{X \in \mathcal{P}(U) \mid \exists A \in \mathcal{F}. X = \bigcup A\}$$

$$\Leftrightarrow \exists X \in \mathcal{P}(U). \exists A \in \mathcal{F}. X = \bigcup A \wedge x \in X$$

$$\Leftrightarrow \exists A \in \mathcal{F}. x \in \bigcup A$$

$$\Leftrightarrow \exists A \in \mathcal{F}. \exists S \in A. x \in S.$$



Example: Big intersection

- $\mathcal{S} =_{\text{def}} \left\{ S \subseteq [5] \mid \begin{array}{l} \text{the sum of the elements} \\ \text{of } S \text{ equals } 6 \end{array} \right\}$

$$= \left\{ \{2, 4\}, \{0, 2, 4\}, \{1, 2, 3\} \right\}$$

- $\bigcap \mathcal{S}$ is the intersection of the sets in \mathcal{S}

$$n \in \bigcap \mathcal{S} \iff \forall S \in \mathcal{S}. n \in S$$

$$\bigcap \mathcal{S} = \{2\}$$

Big intersections

Definition 92 Let \mathcal{U} be a set. For a collection of sets $\mathcal{F} \subseteq \mathcal{P}(\mathcal{U})$, we let the big intersection (relative to \mathcal{U}) be defined as

$$\bigcap \mathcal{F} = \{x \in \mathcal{U} \mid \forall A \in \mathcal{F}. x \in A\} \quad .$$

Examples:

- $\bigcap (\mathcal{P}(U)) = \{x \in U \mid \forall S \in \mathcal{P}(U). x \in S\}$
 $= \{x \in U \mid \underline{\text{false}}\} = \emptyset$

- $\bigcap \emptyset = \{x \in U \mid \forall S \in \emptyset. x \in S\}$
 $= \{x \in U \mid \text{true}\} = U.$

Theorem 93 Let

$$\mathcal{F} = \left\{ S \subseteq \mathbb{R} \mid (0 \in S) \wedge (\forall x \in \mathbb{R}. x \in S \implies (x+1) \in S) \right\}.$$

Then, (i) $\mathbb{N} \in \mathcal{F}$ and (ii) $\mathbb{N} \subseteq \bigcap \mathcal{F}$. Hence, $\bigcap \mathcal{F} = \mathbb{N}$.

PROOF:

$$\underline{\text{RTP}}: \bigcap \mathcal{F} = \mathbb{N} \iff (\bigcap \mathcal{F} \subseteq \mathbb{N} \wedge \mathbb{N} \subseteq \bigcap \mathcal{F}).$$

$$(i) \underline{\text{RTP}}: \bigcap \mathcal{F} \subseteq \mathbb{N}.$$

We show $\mathbb{N} \in \mathcal{F}$, which is the case.

$$(ii) \underline{\text{RTP}}: \mathbb{N} \subseteq \bigcap \mathcal{F}$$

$$\iff \forall n \in \mathbb{N}. n \in \bigcap \mathcal{F} \iff \forall n \in \mathbb{N}. \forall S \in \mathcal{F}. n \in S.$$

We show

$$\forall n \in \mathbb{N}. P(n)$$

where $P(n) =_{\text{def}} \forall S \in \mathcal{F}. n \in S.$

By induction:

BASE CASE $n=0$: $\forall S \in \mathcal{F}. 0 \in S$. holds by definition of \mathcal{F} .

INDUCTIVE STEP. Let $n \in \mathbb{N}$.

$$(IH) \quad \forall S \in \mathcal{F}. n \in S$$

$$\underline{RTP} \quad \forall S \in \mathcal{F}. (n+1) \in S.$$

Let $S \in \mathcal{F} \Rightarrow \textcircled{1} (\forall x \in \mathbb{R}. x \in S \Rightarrow (x+1) \in S)$

Then by $(IH) \textcircled{2} n \in S$

and so, by $\textcircled{1}$ and $\textcircled{2}$, $(n+1) \in S$.



Proposition: Let U be a set and let $\mathcal{F} \subseteq \mathcal{P}(U)$ be a collection of subsets of U .

(1) For all $S \in \mathcal{P}(U)$,

$$\text{iff } S = \bigcup \mathcal{F}$$

$$[\forall A \in \mathcal{F}. A \subseteq S]$$

$$\text{and } [\forall X \in \mathcal{P}(U). (\forall A \in \mathcal{F}. A \subseteq X) \Rightarrow S \subseteq X]$$

(2) For all $T \in \mathcal{P}(U)$,

$$\text{iff } T = \bigcap \mathcal{F}$$

$$[\forall A \in \mathcal{F}. T \subseteq A]$$

$$\text{and } [\forall Y \in \mathcal{P}(U). (\forall A \in \mathcal{F}. Y \subseteq A) \Rightarrow Y \subseteq T]$$

Union axiom

Every collection of sets has a union.

$$\bigcup \mathcal{F}$$

$$x \in \bigcup \mathcal{F} \iff \exists X \in \mathcal{F}. x \in X$$

For non-empty \mathcal{F} we also have

$$\bigcap \mathcal{F}$$

defined by

$$\forall x. x \in \bigcap \mathcal{F} \iff (\forall X \in \mathcal{F}. x \in X) \quad .$$

$$\{1\} \times A = \{(1, a) \mid a \in A\}$$

$$\{2\} \times B = \{(2, b) \mid b \in B\}$$

$$(\{1\} \times A) \cap (\{2\} \times B) = \emptyset$$

Disjoint unions

Definition 94 The disjoint union $A \uplus B$ of two sets A and B is the set

$$A \uplus B = (\{1\} \times A) \cup (\{2\} \times B) .$$

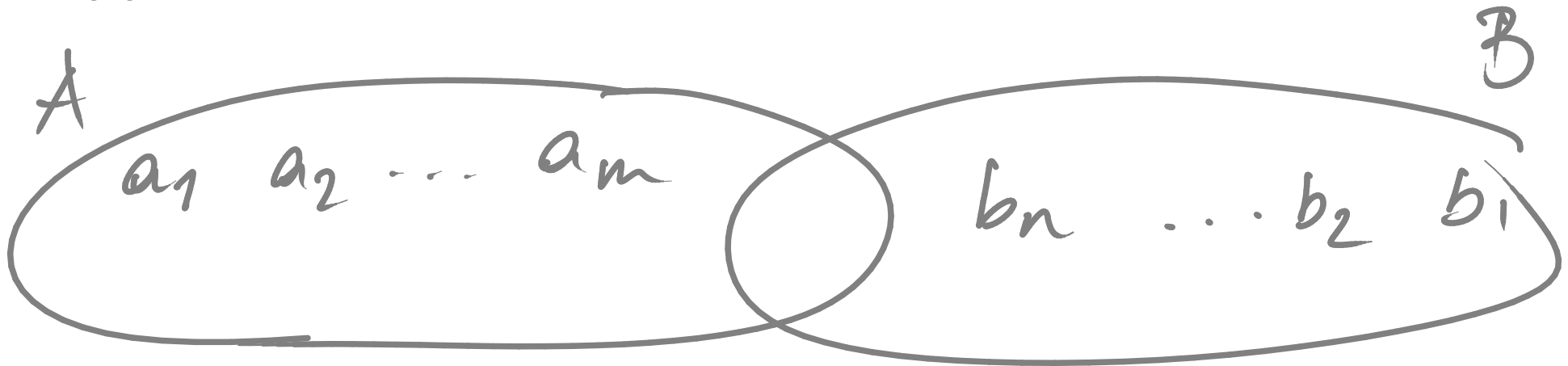
Thus,

$$\forall x. x \in (A \uplus B) \iff (\exists a \in A. x = (1, a)) \vee (\exists b \in B. x = (2, b)) .$$

Proposition 96 For all finite sets A and B ,

$$A \cap B = \emptyset \implies \#(A \cup B) = \#A + \#B .$$

PROOF IDEA:



Corollary 97 For all finite sets A and B ,

$$\#(A \uplus B) = \#A + \#B .$$

Relations

Definition 99 A (binary) relation R from a set A to a set B

$$R : A \rightarrowtail B \quad \text{or} \quad R \in \text{Rel}(A, B) \quad ,$$

is

$$R \subseteq A \times B \quad \text{or} \quad R \in \mathcal{P}(A \times B) \quad .$$

Notation 100 One typically writes $a R b$ for $(a, b) \in R$.

Informal examples:

- ▶ Computation.
- ▶ Typing.
- ▶ Program equivalence.
- ▶ Networks.
- ▶ Databases.

PROGRAM SEMANTICS

$$\underline{sq} : \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}$$

given by all pairs

$$(x, y) \text{ such that } x = y^2$$

In particular, $1 \underline{sq} 1$

$$1 \underline{sq} -1$$

TYPING

$P : \alpha$

E.g.

$(\underline{fn} x \rightarrow x, \underline{bool} \rightarrow bool)$

$(\underline{fn} x \rightarrow x, \underline{nat} \rightarrow nat)$

are in the Typing relation

$(\underline{fn} x \rightarrow x, \underline{bool} \rightarrow \underline{nat})$

is not.

NETWORKS

N — nodes

C — connections

$$C: N \rightarrow N$$

DATABASES

A relation R on sets A_1, A_2, \dots, A_n is defined as a subset

$$R \subseteq A_1 \times A_2 \times \dots \times A_n$$

E.g. $R \subseteq \text{Movies} \times \text{Directors} \times \text{Years} \times \text{Person}$
consisting of all quadruples (m, d, y, p)
such that movie m was directed by director
 d in year y with person p as a member.

Examples:

- Empty relation.

$$\emptyset : A \multimap B$$

$$(a \emptyset b \iff \text{false})$$

- Full relation.

$$(A \times B) : A \multimap B$$

$$(a (A \times B) b \iff \text{true})$$

- Identity (or equality) relation.

$$\text{id}_A = \{ (a, a) \mid a \in A \} : A \multimap A$$

$$(a \text{id}_A a' \iff a = a')$$

- Integer square root.

$$R_2 = \{ (m, n) \mid m = n^2 \} : \mathbb{N} \multimap \mathbb{Z}$$

$$(m R_2 n \iff m = n^2)$$