

Foundations of Computer Science

Lecture #4: More on Lists

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Warm-Up

Question 1a: What is the cost of evaluating `xs @ ys`?

$O(\text{List.length } xs)$

Question 1b: What is the cost of evaluating `x :: xs`?

$O(1)$

Question 2: What is the type of this function?

```
let rec flatten = function
  | []          -> []
  | l :: ls    -> l @ flatten ls
```

Out: `val flatten : 'a list list -> 'a list = <fun>`

Question 3a: What does this return?

```
In [1]: let a = [2];;
```

```
Out[1]: val a : int list = [2]
```

```
In [2]: let b = [3; 4; 5];;
```

```
Out[2]: val b : int list = [3; 4; 5]
```

```
In [3]: a::b;;
```

```
Error: This expression has type int list  
but an expression was expected of type int list list  
Type int is not compatible with int list
```

Question 3b: How to concatenate a and b?

```
In [4]: a @ b;;
```

```
Out[4]: - : int list = [2; 3; 4; 5]
```

Question 3c: Redefine b so that a::b works.

```
In [3]: let b = [b];
```

```
Out[3]: val b : int list list = [[3; 4; 5]]
```

```
In [4]: a::b;;
```

```
Out[4]: - : int list list = [[2]; [3, 4, 5]]
```

A Note on Notation

```
In : let rec append1 = function
      | ([], ys)      -> ys
      | (x::xs, ys)  -> x :: append1 (xs, ys)
```

```
Out: val append : 'a list * 'a list -> 'a list = <fun>
```

```
In : let rec append2 pair =
      match pair with
      | ([], ys)      -> ys
      | (x::xs, ys)  -> x :: append2 (xs, ys)
```

```
Out: val append2 : 'a list * 'a list -> 'a list = <fun>
```

A Note on Notation

```
In : let rec append3 xs ys =  
      match (xs, ys) with  
      | ([], ys)      -> ys  
      | (x::xs, ys) -> x :: append3 xs ys
```

```
Out: val append3 : 'a list -> 'a list -> 'a list = <fun>
```

```
In : let rec append4 xs ys =  
      match xs with  
      | []      -> ys  
      | x::xs -> x :: append4 xs ys
```

```
Out: val append : 'a list -> 'a list -> 'a list = <fun>
```

List Utilities: take and drop

$$xs = \underbrace{[x_0, \dots, x_{i-1}]}_{take(xs, i)}, \underbrace{[x_i, \dots, x_{n-1}]}_{drop(xs, i)}$$

List Utilities: take and drop

```
let rec take = function
| ([], _) -> []
| (x::xs, i) ->
    if i > 0 then
        x :: take (xs, i - 1)
    else
        []

let rec drop = function
| ([], _) -> []
| (x::xs, i) ->
    if i > 0 then
        drop (xs, i - 1)
    else
        x::xs
```

wildcard pattern

List Utilities: take and drop

```
Out: val take : 'a list * int -> 'a list = <fun>
```

```
Out: val drop : 'a list * int -> 'a list = <fun>
```

```
In : let a = [1; 2; 3; 4; 5; 6];;
```

```
In : take (a, 3);;
```

```
Out: - : int list = [1; 2; 3]
```

```
In: drop (a, 3);;
```

```
Out: - : int list = [4; 5; 6]
```


Linear Search

find x in list $[x_1, \dots, x_n]$ by comparing with each element

obviously $O(n)$ TIME

simple & general

ordered searching needs only $O(\log n)$

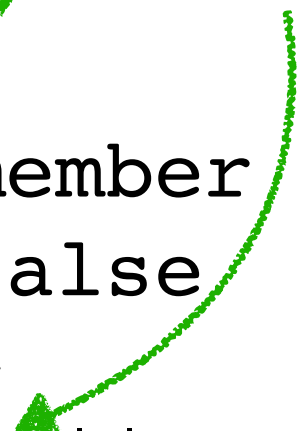
indexed lookup needs only $O(1)$

more about search in Lecture 10...

Equality Tests

if (x=y) then true, else ...

```
let rec member x = function
| [] -> false
| y::l ->
    x = y || member x l
```



Equality testing is *OK* for integers but **NOT** for functions.

Equality Tests (cont.)

```
let rec inter xs ys =  
  match xs, ys with  
  | [], ys      -> []  
  | x::xs, ys ->  
    if member x ys then  
      x :: inter xs ys  
    else  
      inter xs ys
```

Building a List of Pairs

$$\left. \begin{array}{l} [x_1, \dots, x_n] \\ [y_1, \dots, y_n] \end{array} \right\} \mapsto [(x_1, y_1), \dots, (x_n, y_n)]$$

```
let rec zip xs ys =  
  match xs, ys with  
  | (x::xs, y::ys) -> (x, y) :: zip xs ys  
  | _ -> []
```

Building a List of Pairs

```
let rec zip xs ys =  
  match xs, ys with  
  | (x::xs, y::ys) -> (x, y) :: zip xs ys  
  | _ -> []
```

wildcard



The *wildcard* pattern (`_`) matches *anything*.

THE PATTERNS ARE TESTED IN ORDER.

For example, `_` will match: `([], (y::ys))`

```
In : zip [1;2;3;4] ['a';'b';'c'];;
```

```
Out: - : (int * char) list = [(1, 'a'); (2, 'b'); (3, 'c')]
```

Building a List of Pairs

Two functions: **zip** and **unzip**

```
zip : 'a list -> 'b list -> ('a * 'b) list
unzip : ('a * 'b) list -> ('a list * 'b list)
```

Some Syntax

Expressions

let D in E

- Embeds declaration D within expression E
- Useful within a function
- Can perform intermediate computations with function arguments

Building a Pair of Results

Version 1: With a local declaration.

```
let rec unzip = function
  | [] -> ([], [])
  | (x, y)::pairs ->
  declaration let xs, ys = unzip pairs in
  expression (x::xs, y::ys)
```

The `let` construct binds `xs` and `ys` to the results of the recursive call.

Example:

```
In : unzip [(1, 'a'); (2, 'b')];;
```

```
Out: - : int list * char list = ([1; 2], ['a'; 'b'])
```


Building a Pair of Results

Version 2: Replacing local declaration by a function.

```
let conspair ((x, y), (xs, ys)) = (x::xs, y::ys)
```

```
val conspair :  
  ('a * 'b) * ('a list * 'b list) ->  
  'a list * 'b list = <fun>
```

```
let rec unzip = function  
| [] -> ([], [])  
| xy :: pairs -> conspair (xy, unzip pairs)
```

↑
1 pair

↑
list
(of pairs)

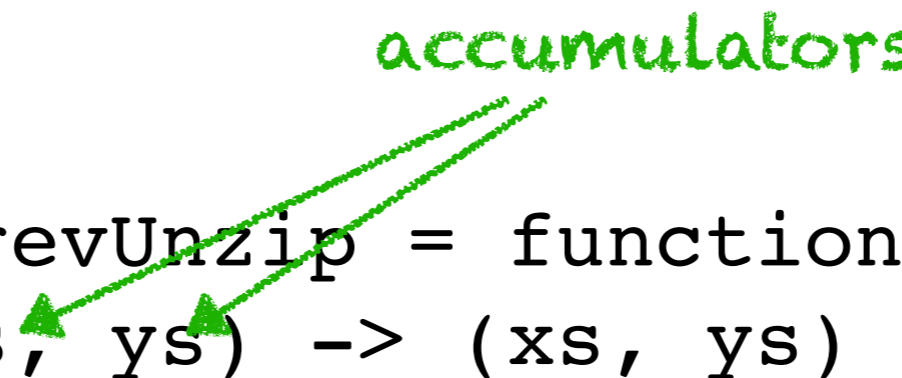
↑
pair of lists

Building a Pair of Results

Version 3: Iterative.

```
let rec revUnzip = function
| ([], xs, ys) -> (xs, ys)
| ((x, y)::pairs, xs, ys) ->
    revUnzip (pairs, x::xs, y::ys)
```

accumulators



Question: How to call revUnzip?

```
revUnzip (pairs, [], []);
```

Question: What's the result of the following?

```
let pairs = [("a", 1); ("b", 2)];;
revUnzip (pairs, [], []);;
```

Out: - : string list * int list = (["b"; "a"], [2; 1])

An Application: Making Change



- Till has unlimited supply of coins, for certain coin values
- List of coins `till` is given in descending order
- Larger coins preferred (tried first)

An Application: Making Change

List of possible coin values



```
let rec change till amt =  
  if amt = 0 then  
    []  
  else  
    match till with  
    | [] -> raise (Failure "no more coins!")  
    | c::till ->  
      if amt < c then  
        change till amt  
      else  
        c :: change (c::till) (amt - c)
```

- The recursion *terminates* when $\text{amt} = 0$.
- Tries the *largest coin first* to use large coins.
- The algorithm is **greedy**, and it CAN FAIL!

An Application: Making Change

```
let till = [50; 20; 10; 5; 2; 1];;  
change till 43;;
```

~~50~~ 20 (amt=23) 20 (amt=3) ~~10~~ ~~5~~ 2 (amt=1) 1 (amt=0)

```
- : int list = [20; 20; 2; 1]
```

```
let till = [5; 2];;  
change till 16;;
```

5 (amt=11) 5 (amt=6) 5 (amt=1) ~~2~~ ? amt≠0, till=[]

Exception: Failure "no more coins!"

An Application: Making Change

```
let rec change till amt =  
  if amt = 0 then  
    []  
  else  
    match till with  
    | [] -> raise (Failure "no more coins!")  
    | c::till ->  
      if amt < c then  
        change till amt  
      else  
        c :: change (c::till) (amt - c)
```

? amt≠0, till=[]

ALL Ways of Making Change

Disclaimer: This is kind of hard.

```
let rec change till amt =
  if amt = 0 then
    [ [] ]
  else
    match till with
    | [] -> []
    | c::till ->
      if amt < c then
        change till amt
      else
        let rec allc = function
          | [] -> []
          | cs :: css -> (c::cs) :: allc css
        in
          allc (change (c::till) (amt - c)) @
            change till amt
```

success (zero)

failure

generates all possible solutions


Out: val change : int list -> int -> int list list = <fun>

ALL Ways of Making Change

...

declaration

```
let rec allc = function  
| [] -> []  
| cs :: css -> (c::cs) :: allc css  
in
```



expression

```
allc (change (c::till) (amt - c)) @  
change till amt
```

use coin c

@

don't use coin c

$c::[\dots], c::[\dots], \dots$

$[\dots], [\dots], \dots$

cons c to solutions for amt-c

solutions for amt

ALL Ways of Making Change

```
In : let till = [5; 3; 2];;
```

```
In : change till 6;;
```

```
Out: - : int list list = [[3; 3]; [2; 2; 2]]
```

```
In : let till = [5; 2];;
```

```
In : change till 16;;
```

```
Out: - : int list list =  
      [[2; 2; 2; 5; 5]; [2; 2; 2; 2; 2; 2; 2; 2]]
```

ALL Ways of Making Change – Faster!

```
let rec change till amt chg chgs =  
  if amt = 0 then  
    chg::chgs  
  else  
    match till with  
    | [] -> chgs  
    | c::till ->  
      if amt < 0 then  
        chgs  
      else  
        change (c::till) (amt - c) (c::chg) use coin  
        change till amt chg chgs  
        solutions that don't use coin
```

We've added **another accumulating parameter!**

Repeatedly improving simple code is called **stepwise refinement.**

ALL Ways of Making Change – Faster!

In : `change [5;3;2] 6 [] [];`

Out: `- : int list list = [[3; 3]; [2; 2; 2]]`