

## Ordered pairing

**Notation:**

$(a, b)$  or  $\langle a, b \rangle$

**Fundamental property:**

$$(a, b) = (x, y) \implies a = x \wedge b = y$$

**A construction:**

For every pair  $a$  and  $b$ , three applications of the pairing axiom provide the set

$$\langle a, b \rangle = \{ \{ a \}, \{ a, b \} \}$$

which defines an ordered pairing of  $a$  and  $b$ .

**Proposition 109 (Fundamental property of ordered pairing)**

For all  $a, b, x, y$ ,

$$\langle a, b \rangle = \langle x, y \rangle \iff (a = x \wedge b = y) .$$

YOUR PROOF:

$$\{ \{ \underline{a} \}, \{ \underline{a, b} \} \} = \{ \{ \underline{x} \}, \{ \underline{x, y} \} \}$$

$$\iff a = x \wedge b = y$$

← easy →

$$\left( \underline{\{a\} = \{x\}} \vee \underline{\{a\} = \{x, y\}} \right)$$

$$\wedge \left( \underline{\{a, b\} = \{x\}} \vee \underline{\{a, b\} = \{x, y\}} \right)$$

FOUR CASES ——— EX FOR YOU.

MY PROOF: Let  $a, b, x, y$  be arbitrary.

$(\Leftarrow)$  *Vacuous.*

$(\Rightarrow)$  Assume  $\{\{a\}, \{a, b\}\} = \{\{x\}, \{x, y\}\}$ .

Then,  $\{a\} = \{x\} \vee \{a\} = \{x, y\}$ ; and, in either case,  $a = x$ .

Hence,  $\{\{a\}, \{a, b\}\} = \{\{a\}, \{a, y\}\}$  and, by Proposition 108.2 (on page 347),  $\{a, b\} = \{a, y\}$  which, again by Proposition 108.2, implies  $b = y$ .

## Products

The product  $A \times B$  of two sets  $A$  and  $B$  is the set

$$A \times B = \{ x \mid \exists a \in A, b \in B. x = (a, b) \}$$

where

$$\left( \begin{array}{l} \forall a_1, a_2 \in A, b_1, b_2 \in B. \\ (a_1, b_1) = (a_2, b_2) \iff (a_1 = a_2 \wedge b_1 = b_2) \end{array} \right) .$$

Thus,

$$\forall x \in A \times B. \exists! a \in A. \exists! b \in B. x = (a, b) .$$

More generally, for a fixed natural number  $n$  and sets  $A_1, \dots, A_n$ , we have

$$\begin{aligned} \prod_{i=1}^n A_i &= A_1 \times \cdots \times A_n \\ &= \{x \mid \exists a_1 \in A_1, \dots, a_n \in A_n. x = (a_1, \dots, a_n)\} \end{aligned}$$

where

$$\forall a_1, a'_1 \in A_1, \dots, a_n, a'_n \in A_n.$$

$$(a_1, \dots, a_n) = (a'_1, \dots, a'_n) \iff (a_1 = a'_1 \wedge \cdots \wedge a_n = a'_n) .$$

**NB** Cunningly enough, the definition is such that  $\prod_{i=1}^0 A_i = \{()\}$ .

**Notation 110** For a natural number  $n$  and a set  $A$ , one typically writes  $A^n$  for  $\prod_{i=1}^n A$ .

## Pattern-matching notation

**Example:** The subset of ordered pairs from a set  $A$  with equal components is formally

$$\{x \in A \times A \mid \exists a_1 \in A. \exists a_2 \in A. x = (a_1, a_2) \wedge a_1 = a_2\}$$

but often abbreviated using *pattern-matching notation* as

$$\{(a_1, a_2) \in A \times A \mid a_1 = a_2\}$$

**Notation:** For a property  $P(a, b)$  with  $a$  ranging over a set  $A$  and  $b$  ranging over a set  $B$ ,

$$\{(a, b) \in A \times B \mid P(a, b)\}$$

abbreviates

$$\{x \in A \times B \mid \exists a \in A. \exists b \in B. x = (a, b) \wedge P(a, b)\}$$

**Proposition 111** For all finite sets  $A$  and  $B$ ,

$$\#(A \times B) = \#A \cdot \#B .$$

PROOF IDEA <sup>a</sup> :

$$\begin{aligned} A &= \{a_1 \text{ --- } a_n\} \\ B &= \{b_1 \text{ --- } b_j\} \\ A \times B &= \{ (a_1, b_1) \text{ --- } (a_1, b_j) \\ &\quad \vdots \\ &\quad (a_n, b_1) \text{ --- } (a_n, b_j) \} \end{aligned}$$

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<sup>a</sup>See Theorem 162.2 on page 439.



## Sets and logic

$\mathcal{P}(U)$	$\{ \text{false}, \text{true} \}$
$\emptyset$	<b>false</b>
$U$	<b>true</b>
$\cup$	$\vee$
$\cap$	$\wedge$
$(\cdot)^c$	$\neg(\cdot)$
$\bigcup$	$\exists$
$\bigcap$	$\forall$

## Big unions

**Example:**

~~$\sum \phi = 0$~~   $\sum \phi = 0$

- Consider the family of sets

$$\mathcal{T} = \left\{ T \subseteq \underline{[5]} \mid \begin{array}{l} \text{the sum of the elements of} \\ T \text{ is less than or equal } 2 \end{array} \right\}$$
$$= \{ \underline{\emptyset}, \underline{\{0\}}, \underline{\{1\}}, \underline{\{0, 1\}}, \underline{\{0, 2\}} \}$$

- The *big union* of the family  $\mathcal{T}$  is the set  $\bigcup \mathcal{T}$  given by the union of the sets in  $\mathcal{T}$ :

$$n \in \bigcup \mathcal{T} \iff \exists T \in \mathcal{T}. n \in T .$$

Hence,  $\bigcup \mathcal{T} = \{0, 1, 2\}$ .



**Examples:**

1. For  $A, A_1, A_2 \in \mathcal{P}(U)$ ,

$$\bigcup \emptyset = \emptyset$$

$$\bigcup \{A\} = A$$

$$\bigcup \{A_1, A_2\} = A_1 \cup A_2$$

$$\bigcup \{A, A_1, A_2\} = A \cup A_1 \cup A_2$$

2. For  $\mathcal{F} \in \mathcal{P}(\mathcal{P}(\mathcal{P}(\mathcal{U})))$ , let us introduce the notation

$$\underbrace{\left\{ \bigcup \mathcal{A} \in \mathcal{P}(\mathcal{U}) \mid \mathcal{A} \in \mathcal{F} \right\}}_{\text{map } \bigcup \mathcal{F}}$$

for the set

$$\left\{ X \in \mathcal{P}(\mathcal{U}) \mid \exists \mathcal{A} \in \mathcal{F}. X = \bigcup \mathcal{A} \right\} \in \mathcal{P}(\mathcal{P}(\mathcal{U}))$$

noticing that this is justified by the fact that, for all  $x \in \mathcal{U}$ ,

$$\begin{aligned} x \in \bigcup \{ X \in \mathcal{P}(\mathcal{U}) \mid \exists \mathcal{A} \in \mathcal{F}. X = \bigcup \mathcal{A} \} \\ \iff \exists X \in \mathcal{P}(\mathcal{U}). \exists \mathcal{A} \in \mathcal{F}. X = \bigcup \mathcal{A} \wedge x \in X \\ \iff \exists \mathcal{A} \in \mathcal{F}. x \in \bigcup \mathcal{A} \end{aligned}$$

$$\mathcal{F} \in \mathcal{P}(\mathcal{P}(\mathcal{P}(U)))$$

$$\cup(\cup \mathcal{F}) = ?$$

$$\mathcal{F} = \{ \dots A, \dots B \dots \}$$

$$A = \{ A_1, A_2, \dots, A_k \}$$

$$B = \{ B_1, B_2, \dots, B_j \}$$

$$\cup(\cup \mathcal{F}) =$$

$$\cup \{ \dots A_1, A_2, \dots, A_k, \dots \\ B_1, B_2, \dots, B_j \dots \}$$

$$= \{ \dots A_1 \cup A_2 \cup \dots \cup A_k \cup \dots \\ \dots \cup B_1 \cup B_2 \cup \dots \cup B_j \cup \dots \}$$

$$= \{ \dots \cup(\cup A) \cup \dots \cup(\cup B) \dots \}$$

$$= \cup(\text{map } \cup \mathcal{F})$$

We then have the following *associativity law*:

**Proposition 113** For all  $\mathcal{F} \in \mathcal{P}(\mathcal{P}(\mathcal{P}(U)))$ ,

$$\underline{\cup(\cup\mathcal{F})} = \underline{\cup\{\cup A \in \mathcal{P}(U) \mid A \in \mathcal{F}\}} \in \mathcal{P}(U) .$$

**Btw** In trying to understand this statement, ponder about the following analogous identity for the ML `list` datatype constructor: for all `F : 'a list list list`,

```
flatten ( flatten F )  
= flatten ( map flatten F ) : 'a list
```

The above two identities are the *associativity law* of a mathematical structure known as a monad, which has become a fundamental tool in functional programming.

MY PROOF: For  $\mathcal{F} \in \mathcal{P}(\mathcal{P}(\mathcal{P}(\mathcal{U})))$  and  $x \in \mathcal{U}$ , one calculates that:

$$x \in \bigcup (\bigcup \mathcal{F})$$

$$\iff \exists X \in \bigcup \mathcal{F}. x \in X$$

$$\iff \exists \mathcal{A} \in \mathcal{F}. \exists X \in \mathcal{A}. x \in X$$

$$\iff \exists \mathcal{A} \in \mathcal{F}. x \in \bigcup \mathcal{A}$$

$$\iff x \in \bigcup \{ \bigcup \mathcal{A} \in \mathcal{P}(\mathcal{U}) \mid \mathcal{A} \in \mathcal{F} \}$$



## Big intersections

### Example:

- ▶ Consider the family of sets

$$\mathcal{S} = \left\{ S \subseteq [5] \mid \begin{array}{l} \text{the sum of the elements of} \\ S \text{ is less than or equal } 6 \end{array} \right\}$$

😊

$$= \{ \{2,4\}, \{0,2,4\}, \{1,2,3\} \}$$

ϕ ? ?

- ▶ The *big intersection* of the family  $\mathcal{S}$  is the set  $\bigcap \mathcal{S}$  given by the intersection of the sets in  $\mathcal{S}$ :

$$n \in \bigcap \mathcal{S} \iff \forall S \in \mathcal{S}. n \in S .$$

Hence,  $\bigcup \mathcal{S} = \{2\}$ .

**Definition 114** Let  $U$  be a set. For a collection of sets  $\mathcal{F} \subseteq \mathcal{P}(U)$ , we let the big intersection (relative to  $U$ ) be defined as

$$\bigcap \mathcal{F} = \{ \underline{x \in U} \mid \forall A \in \mathcal{F}. x \in A \} .$$

**Examples:** For  $A, A_1, A_2 \in \mathcal{P}(U)$ ,

$$\bigcap \emptyset = U$$

$$\bigcap \{A\} = A$$

$$\bigcap \{A_1, A_2\} = A_1 \cap A_2$$

$$\bigcap \{A, A_1, A_2\} = A \cap A_1 \cap A_2$$

**Theorem 115** *Let*

$$\mathcal{F} = \left\{ S \subseteq \mathbb{R} \mid (0 \in S) \wedge (\forall x \in \mathbb{R}. x \in S \implies (x+1) \in S) \right\} .$$

*Then, (i)  $\mathbb{N} \in \mathcal{F}$  and (ii)  $\mathbb{N} \subseteq \bigcap \mathcal{F}$ . Hence,  $\bigcap \mathcal{F} = \mathbb{N}$ .*

**NB** This result is typically interpreted as stating that:

$\mathbb{N}$  is the least set of numbers containing 0 and closed under successors.

? prop. 116?

Marcello will explain



**Proposition 116** Let  $U$  be a set and let  $\mathcal{F} \subseteq \mathcal{P}(U)$  be a family of subsets of  $U$ .

1. For all  $S \in \mathcal{P}(U)$ ,

$$S = \bigcup \mathcal{F}$$

iff

$$[\forall A \in \mathcal{F}. A \subseteq S]$$

$$\wedge [\forall X \in \mathcal{P}(U). (\forall A \in \mathcal{F}. A \subseteq X) \Rightarrow S \subseteq X]$$

$= Q(S)$

2. For all  $T \in \mathcal{P}(U)$ ,

$$T = \bigcap \mathcal{F}$$

iff

$$[\forall A \in \mathcal{F}. T \subseteq A]$$

$$\wedge [\forall Y \in \mathcal{P}(U). (\forall A \in \mathcal{F}. Y \subseteq A) \Rightarrow Y \subseteq T]$$

$Q(X)$

Smallest  
for  $Q$

$Q'(T)$

$Q'(Y)$

biggest

for  $Q'$

## Union axiom

Every collection of sets has a union.

The set whose existence is postulated by the union axiom for a collection  $\mathcal{F}$  is typically denoted

$$\bigcup \mathcal{F}$$

and, in the case  $\mathcal{F} = \{A, B\}$ , abbreviated to

$$A \cup B .$$

Thus,

$$x \in \bigcup \mathcal{F} \iff \exists X \in \mathcal{F} . x \in X ,$$

and hence

$$x \in (A \cup B) \iff (x \in A) \vee (x \in B) .$$

Using the separation and union axioms, for every collection  $\mathcal{F}$ , consider the set

$$\{x \in \bigcup \mathcal{F} \mid \forall X \in \mathcal{F}. x \in X\} .$$

For non-empty  $\mathcal{F}$  this set is denoted

$$\bigcap \mathcal{F}$$

because, in this case,

$$\forall x. x \in \bigcap \mathcal{F} \iff (\forall X \in \mathcal{F}. x \in X) .$$

In particular, for  $\mathcal{F} = \{A, B\}$ , this is abbreviated to

$$A \cap B$$

with defining property

$$\forall x. x \in (A \cap B) \iff (x \in A) \wedge (x \in B) .$$