

Proposition 63 Let m be a positive integer. A modular integer k in \mathbb{Z}_m has a reciprocal if, and only if, there exist integers i and j such that $k \cdot i + m \cdot j = 1$.

PROOF: Let m be a pos. int.

Let k be an int s.t. $0 \leq k < m$.

RTP: k has a reciprocal

iff \exists int i, j . $k \cdot i + m \cdot j = 1$

(\Rightarrow) Assume k has a reciprocal; that is, there is k an int. s.t. $0 \leq \bar{k} < m$ with $k \cdot \bar{k} \equiv 1 \pmod{m}$.
Then $k \cdot \bar{k} - 1 = l \cdot m$ for an int l , and 1 is the int. linear comb. $k \cdot \bar{k} + (-l) \cdot m$.

(\Leftarrow) Assume: $\exists i, j \text{ int. } k \cdot i + m \cdot j = 1$

RTP: $\exists \bar{k}$ s.t. $0 \leq \bar{k} < m$ and $k \cdot \bar{k} \equiv 1 \pmod{m}$.

Let i_0 and j_0 be int. s.t. $k \cdot i_0 + m \cdot j_0 = 1$

Then $k \cdot i_0 - 1$ is a multiple of m ; and so

$k \cdot i_0 \equiv 1 \pmod{m}$. Also $i_0 \equiv [i_0]_m \pmod{m}$

with $0 \leq [i_0]_m < m$. Thus, $k \cdot [i_0]_m \equiv 1 \pmod{m}$



Integer linear combinations

Definition 64 An integer r is said to be a linear combination of a pair of integers m and n whenever there are integers s and t such that $s \cdot m + t \cdot n = r$.

Proposition 65 Let m be a positive integer. A modular integer k in \mathbb{Z}_m has a reciprocal if, and only if, 1 is an integer linear combination of m and k .

Proposition Let a and b be integers.

For all integers d , the following are equivalent:

1. $d|a$ and $d|b$

2. for all integers i and j , $d|(ai + bj)$

PROOF: Let a and b be int. Let d be int.

(\Rightarrow) Assume: ^① $d|a$ and ^② $d|b$.

Let i, j int. From ①, $a = d \cdot k$ for an int k , from ②
 $b = d \cdot l$ for an int l . Consider $ai + bj = k \cdot i \cdot d + l \cdot j \cdot d$
 $= (k \cdot i + l \cdot j) \cdot d$. Then $d|ai + bj$.

(\Leftarrow) Assume \forall int. i, j . $d \mid (ai + bj)$

In particular, this is the case instantiating $i=1$ and $j=0$, that is, $d \mid a$; analogously, instantiating $i=0$ and $j=1$, we have $d \mid b$.

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Important mathematical jargon: Sets

Very roughly, sets are the mathematicians' data structures. Informally, we will consider a set as a (well-defined, unordered) collection of mathematical objects, called the elements (or members) of the set.

Set membership

The symbol ' \in ' known as the *set membership* predicate is central to the theory of sets, and its purpose is to build statements of the form

$$x \in A$$

that are true whenever it is the case that the object x is an element of the set A , and false otherwise.

Defining sets

The set	of even primes	is	{2}
	of booleans		{true, false}
	[-2..3]		{-2, -1, 0, 1, 2, 3}

Set comprehension

The basic idea behind set comprehension is to define a set by means of a property that precisely characterises all the elements of the set.

$$a \in \{x \in A \mid P(x)\} \Leftrightarrow (a \in A \wedge P(a))$$

Notations:

$$\{x \in A \mid P(x)\} \quad , \quad \{x \in A : P(x)\}$$

Set equality

Two sets are equal precisely when they have the same elements

Examples:

▶ $\{x \in \mathbb{N} : 2 \mid x \wedge x \text{ is prime}\} = \{2\}$

▶ For a positive integer m ,

$$\{x \in \mathbb{Z} : m \mid x\} = \{x \in \mathbb{Z} : x \equiv 0 \pmod{m}\}$$

▶ $\{d \in \mathbb{N} : d \mid 0\} = \mathbb{N}$

Equivalent predicates specify equal sets:

$$\{x \in A \mid P(x)\} = \{x \in A \mid Q(x)\}$$

iff

$$\forall x. P(x) \iff Q(x)$$

Example: For a positive integer m ,

$$\begin{aligned} & \{x \in \mathbb{Z}_m \mid x \text{ has a reciprocal in } \mathbb{Z}_m\} \\ = & \{x \in \mathbb{Z}_m \mid 1 \text{ is an integer linear combination of } m \text{ and } x\} \end{aligned}$$

Greatest common divisor

Given a natural number n , the set of its *divisors* is defined by set comprehension as follows

$$D(n) = \{ d \in \mathbb{N} : d \mid n \} .$$

Example 67

1. $D(0) = \mathbb{N}$

2. $D(1224) = \left\{ \begin{array}{l} 1, 2, 3, 4, 6, 8, 9, 12, 17, 18, 24, 34, 36, 51, 68, \\ 72, 102, 136, 153, 204, 306, 408, 612, 1224 \end{array} \right\}$

Remark Sets of divisors are hard to compute. However, the computation of the greatest divisor is straightforward. :)

Going a step further, what about the *common divisors* of pairs of natural numbers? That is, the set

$$\text{CD}(m, n) = \{ d \in \mathbb{N} : d \mid m \wedge d \mid n \}$$

for $m, n \in \mathbb{N}$.

Example 68

$$\text{CD}(1224, 660) = \{ 1, 2, 3, 4, 6, 12 \}$$

Since $\text{CD}(n, n) = D(n)$, the computation of common divisors is as hard as that of divisors. But, what about the computation of the *greatest common divisor*?

Lemma 71 (Key Lemma) Let m and m' be natural numbers and let n be a positive integer such that $m \equiv m' \pmod{n}$. Then,

$$CD(m, n) = CD(m', n) .$$

PROOF:

$$m \equiv \underline{\text{rem}}(m, n) \pmod{n}$$

$$\begin{aligned} CD(m, n) \\ &= CD(\underline{\text{rem}}(m, n), n) \end{aligned}$$

$$m \equiv m + in \pmod{n}$$

$$= CD(m + n, n)$$

$$= CD(m - n, n)$$

Assume: ① $m \equiv m' \pmod{n}$

$$\text{CD}(m, n) \stackrel{?}{=} \text{CD}(m', n)$$

equiv

$$\forall d. (d|m \wedge d|n) \Leftrightarrow (d|m' \wedge d|n)$$

(\Rightarrow) Assume: ② $d|m$ and ③ $d|n$

So $d|n$ and RTP: $d|m'$

From ① $m - m' = i \cdot n$ for an int. i .

So $m' = m + (-i)n$ and from ② and ③, d divides

any int. lin. comb. of m and n ; in particular, m'

(\Leftarrow) Analogously.



Lemma 73 For all positive integers m and n ,

$$\text{CD}(m, n) = \begin{cases} D(n) & , \text{ if } n \mid m \\ \text{CD}(n, \text{rem}(m, n)) & , \text{ otherwise} \end{cases}$$

Since a positive integer n is the greatest divisor in $D(n)$, the lemma suggests a recursive procedure:

$$\text{gcd}(m, n) = \begin{cases} n & , \text{ if } n \mid m \\ \text{gcd}(n, \text{rem}(m, n)) & , \text{ otherwise} \end{cases}$$

for computing the *greatest common divisor*, of two positive integers m and n . This is

Euclid's Algorithm

gcd

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fun gcd( m , n )  
  = let  
    val ( q , r ) = divalg( m , n )  
  in  
    if r = 0 then n  
    else gcd( n , r )  
  end
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Example 74 ($\text{gcd}(13, 34) = 1$)

$$\begin{aligned}\text{gcd}(13, 34) &= \text{gcd}(34, 13) \\ &= \text{gcd}(13, 8) \\ &= \text{gcd}(8, 5) \\ &= \text{gcd}(5, 3) \\ &= \text{gcd}(3, 2) \\ &= \text{gcd}(2, 1) \\ &= 1\end{aligned}$$

NB If gcd terminates on input (m, n) with output $\text{gcd}(m, n)$ then $\text{CD}(m, n) = D(\text{gcd}(m, n))$.

Proposition 75 For all natural numbers m, n and a, b , if $CD(m, n) = D(a)$ and $CD(m, n) = D(b)$ then $a = b$.

Proposition 76 For all natural numbers m, n and k , the following statements are equivalent:

1. $CD(m, n) = D(k)$.

2. $\blacktriangleright k \mid m \wedge k \mid n$, and

\blacktriangleright for all natural numbers d , $d \mid m \wedge d \mid n \implies d \mid k$.

Exercise

Definition 77 For natural numbers m, n the unique natural number k such that

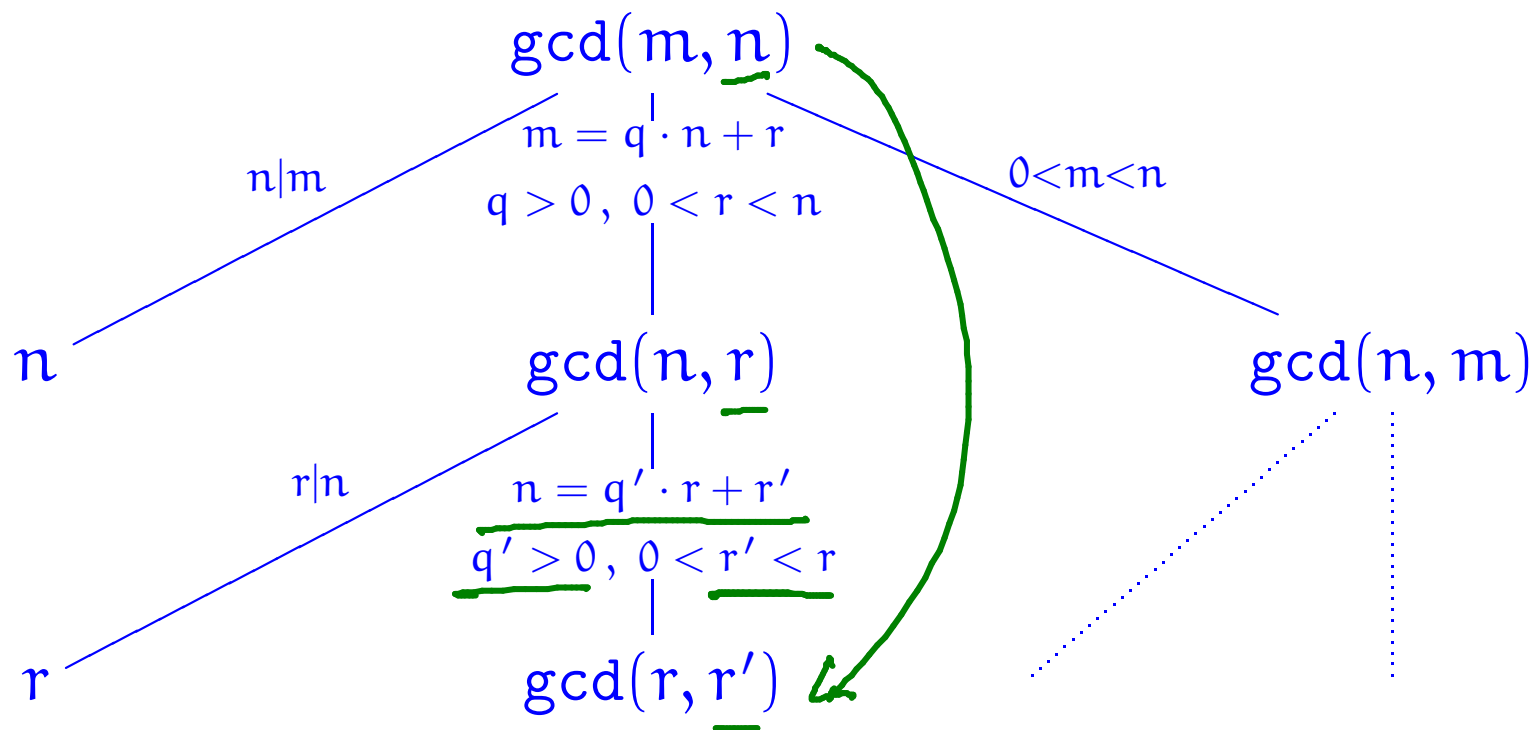
- ▶ $k \mid m \wedge k \mid n$, and
- ▶ for all natural numbers d , $d \mid m \wedge d \mid n \implies d \mid k$.

is called the **greatest common divisor** of m and n , and denoted $\gcd(m, n)$.

Theorem 78 *Euclid's Algorithm \gcd terminates on all pairs of positive integers and, for such m and n , the positive integer $\gcd(m, n)$ is the greatest common divisor of m and n in the sense that the following two properties hold:*

- (i) both $\gcd(m, n) \mid m$ and $\gcd(m, n) \mid n$, and*
- (ii) for all positive integers d such that $d \mid m$ and $d \mid n$ it necessarily follows that $d \mid \gcd(m, n)$.*

PROOF:



$$n = q' \cdot r + r' > q' \cdot r' + r' = (q' + 1) \cdot r' \geq 2 \cdot r'$$

$$\Rightarrow r' < n/2$$