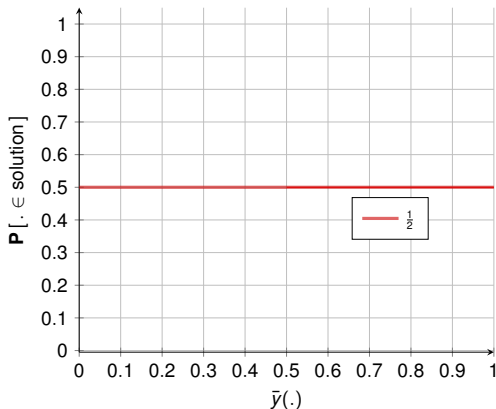


# Randomised Algorithms

## Lecture 11: Spectral Graph Theory

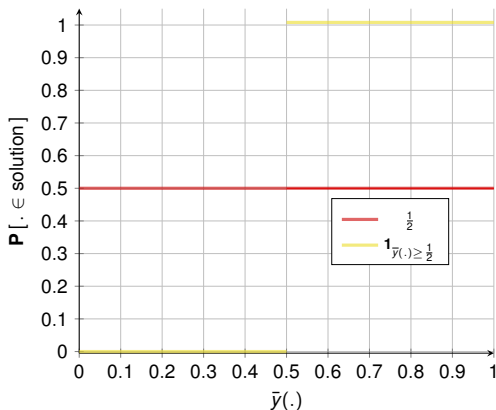
Thomas Sauerwald (tms41@cam.ac.uk)

## Recap: Different Rounding Procedures



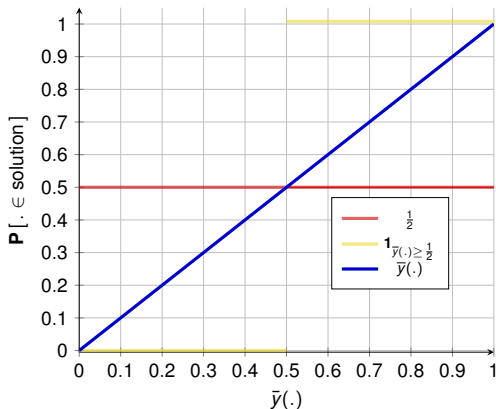
- MAX-CUT, MAX-3-CNF, MAX-CNF: Uniform Guessing

## Recap: Different Rounding Procedures



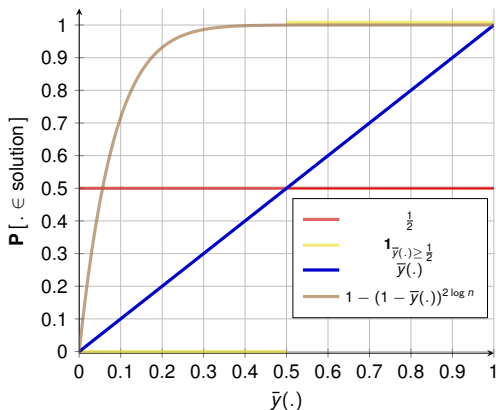
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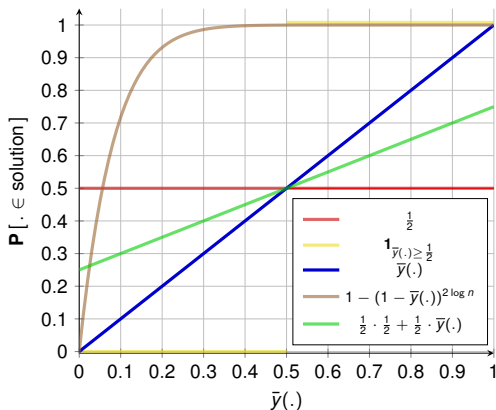
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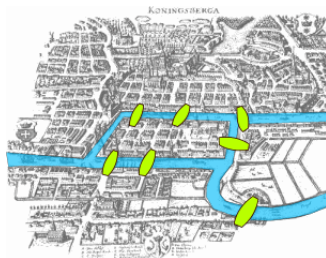
Introduction to (Spectral) Graph Theory and Clustering

Matrices, Spectrum and Structure

A Simplified Clustering Problem

# Origin of Graph Theory

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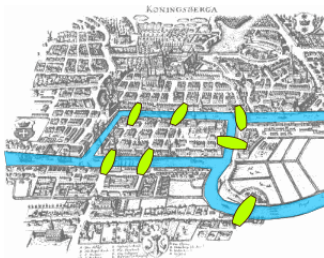


Source: Wikipedia

## Seven Bridges at Königsberg 1737



# Origin of Graph Theory



Source: Wikipedia

Seven Bridges at Königsberg 1737

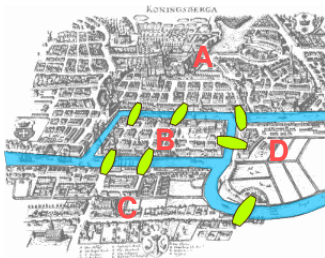


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Leonhard Euler (1707-1783)

Is there a tour which crosses  
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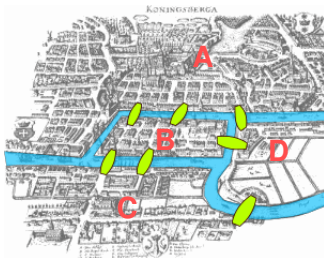


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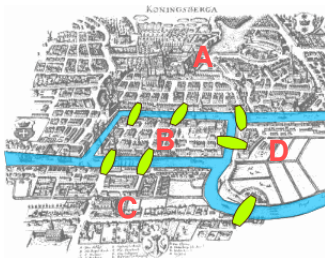
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D

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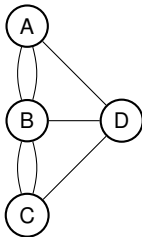
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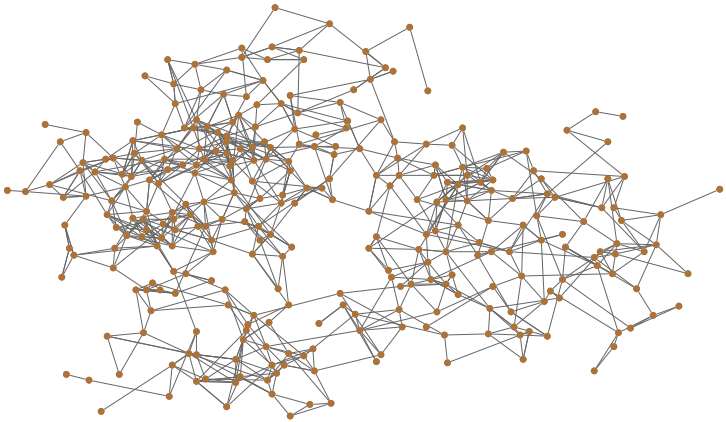
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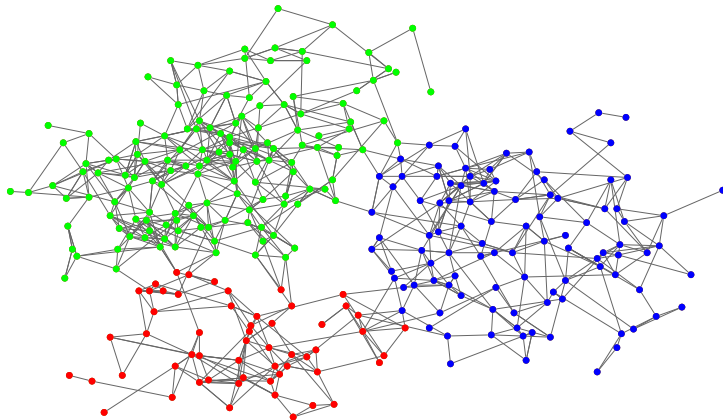
# Graphs Nowadays: Clustering

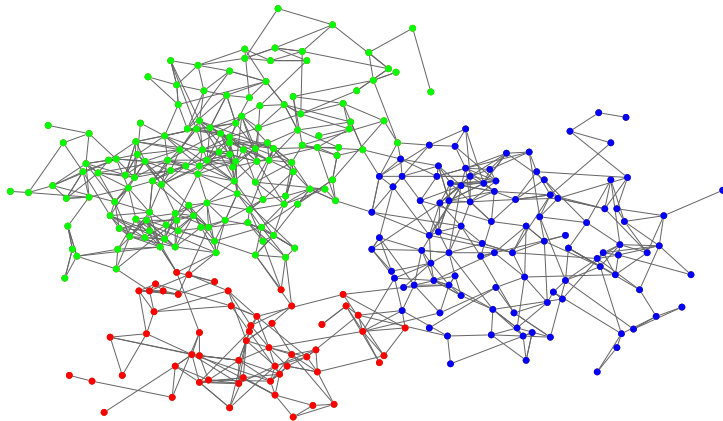
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# Graphs Nowadays: Clustering

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**Goal:** Use spectrum of graphs (unstructured data) to extract clustering (communities) or other structural information.

- Applications of Graph Clustering
  - Community detection
  - Group webpages according to their topics
  - Find proteins performing the same function within a cell
  - Image segmentation
  - Identify bottlenecks in a network
  - ...



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(there is no ground truth (usually), and we cannot learn from mistakes!)

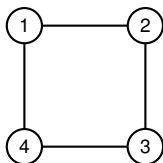
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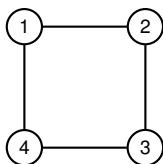
## Graphs



## Matrices

$$\begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

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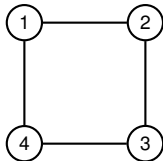
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Introduction to (Spectral) Graph Theory and Clustering

**Matrices, Spectrum and Structure**

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## Adjacency Matrix

---

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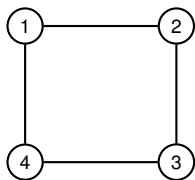
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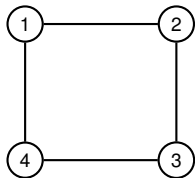
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Properties of  $\mathbf{A}$ :

- The sum of elements in each row/column  $i$  equals the degree of the corresponding vertex  $i$ ,  $\deg(i)$
- Since  $G$  is undirected,  $\mathbf{A}$  is symmetric

# Eigenvalues and Graph Spectrum of A

---

## Eigenvalues and Eigenvectors

Let  $\mathbf{M} \in \mathbb{R}^{n \times n}$ ,  $\lambda \in \mathbb{C}$  is an **eigenvalue** of  $\mathbf{M}$  if and only if there exists  $x \in \mathbb{C}^n \setminus \{\mathbf{0}\}$  such that

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= **orthogonal** and **normalised**

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Remark: For **symmetric** matrices we have **algebraic multiplicity** = **geometric multiplicity** (otherwise  $\geq$ )

## Example 1

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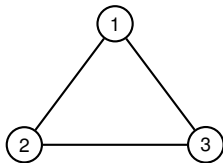
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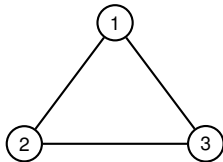
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**Bonus:** Can you find a short-cut to  $\det(\mathbf{A} - \lambda \cdot \mathbf{I})$ ?

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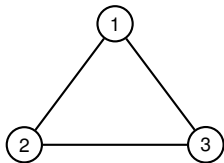
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$$\mathbf{A} = \begin{pmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix}$$

Solution:

- The three eigenvalues are  $\lambda_1 = \lambda_2 = -1, \lambda_3 = 2$ .
- The three eigenvectors are (for example):

$$f_1 = \begin{pmatrix} 1 \\ 0 \\ -1 \end{pmatrix}, \quad f_2 = \begin{pmatrix} -\frac{1}{2} \\ 1 \\ -\frac{1}{2} \end{pmatrix}, \quad f_3 = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}.$$

## Laplacian Matrix

---

### Laplacian Matrix

Let  $G = (V, E)$  be a  $d$ -regular undirected graph. The (normalised) Laplacian matrix of  $G$  is the  $n$  by  $n$  matrix  $\mathbf{L}$  defined as

$$\mathbf{L} = \mathbf{I} - \frac{1}{d}\mathbf{A},$$

where  $\mathbf{I}$  is the  $n \times n$  identity matrix.



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**Question:** How does the matrix  $\mathbf{I} - \frac{1}{d} \cdot \mathbf{A}$  look like?

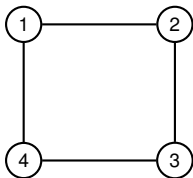
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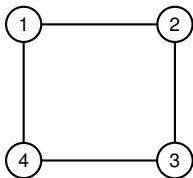
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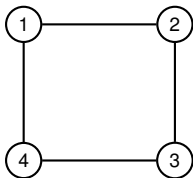
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Properties of  $\mathbf{L}$ :

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- $\mathbf{L}$  is symmetric

## Relating Spectrum of Adjacency Matrix and Laplacian Matrix

---

Correspondence between Adjacency and Laplacian Matrix

**A** and **L** have the same set of eigenvectors.



**Exercise:** Prove this correspondence. Hint: Use that  $\mathbf{L} = \mathbf{I} - \frac{1}{d}\mathbf{A}$ .

*[Exercise 11/12.1]*

# Eigenvalues and Graph Spectrum of L

---

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— Lemma —

Let  $\mathbf{L}$  be the Laplacian matrix of an undirected, regular graph  $G = (V, E)$  with eigenvalues  $\lambda_1 \leq \dots \leq \lambda_n$ .

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3.  $\lambda_n \leq 2$
4.  $\lambda_n = 2$  iff there exists a bipartite connected component.



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3.  $\lambda_n \leq 2$
4.  $\lambda_n = 2$  iff there exists a bipartite connected component.

The proof of these properties is based on a powerful characterisation of eigenvalues/vectors!

## A Min-Max Characterisation of Eigenvalues and Eigenvectors

### Courant-Fischer Min-Max Formula (non-examinable)

Let  $\mathbf{M}$  be an  $n$  by  $n$  symmetric matrix with eigenvalues  $\lambda_1 \leq \dots \leq \lambda_n$ . Then,

$$\lambda_k = \min_{S: \dim(S)=k} \max_{x \in S, x \neq 0} \frac{x^T \mathbf{M} x}{x^T x},$$

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— Lemma —

Let  $\mathbf{L}$  be the Laplacian matrix of a  $d$ -regular graph  $G = (V, E)$  with  $n$  vertices. For any  $x \in \mathbb{R}^n$ ,

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Proof:

$$\begin{aligned} x^T \mathbf{L} x &= x^T \left( \mathbf{I} - \frac{1}{d} \mathbf{A} \right) x = x^T x - \frac{1}{d} x^T \mathbf{A} x \\ &= \sum_{u \in V} x_u^2 - \frac{2}{d} \sum_{\{u,v\} \in E} x_u x_v \\ &= \frac{1}{d} \sum_{\{u,v\} \in E} (x_u^2 + x_v^2 - 2x_u x_v) \\ &= \sum_{\{u,v\} \in E} \frac{(x_u - x_v)^2}{d}. \end{aligned}$$



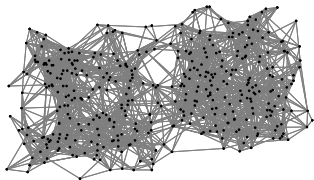
## Visualising a Graph

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**Question:** How can we visualize a complicated object like an unknown graph with many vertices in low-dimensional space?

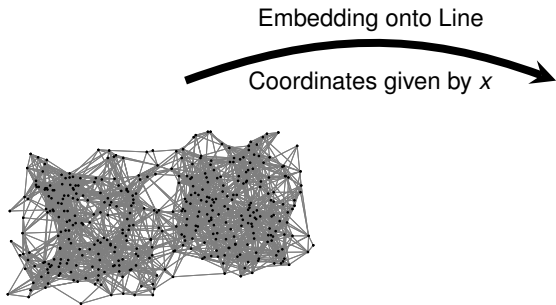
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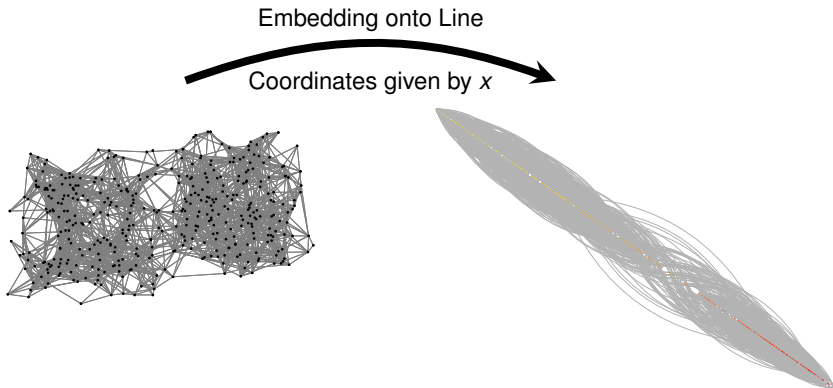
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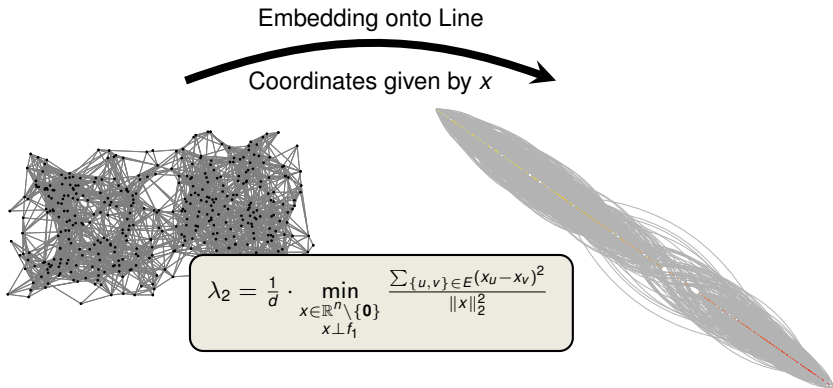
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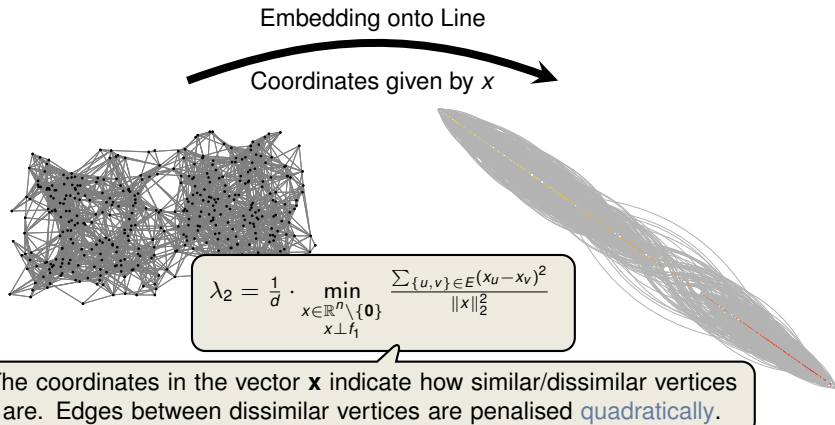
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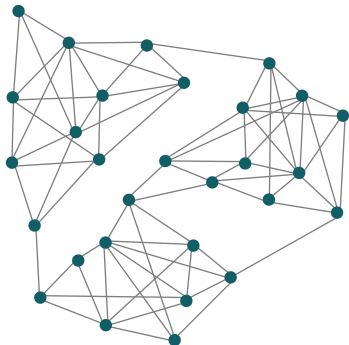
Introduction to (Spectral) Graph Theory and Clustering

Matrices, Spectrum and Structure

**A Simplified Clustering Problem**

## A Simplified Clustering Problem

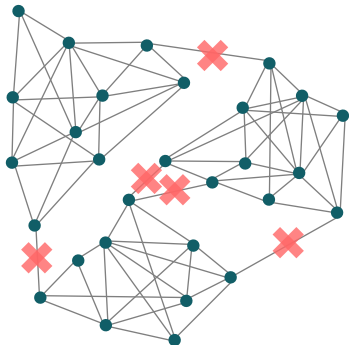
*Partition the graph into **connected components** so that any pair of vertices in the same component is connected, but vertices in different components are not.*





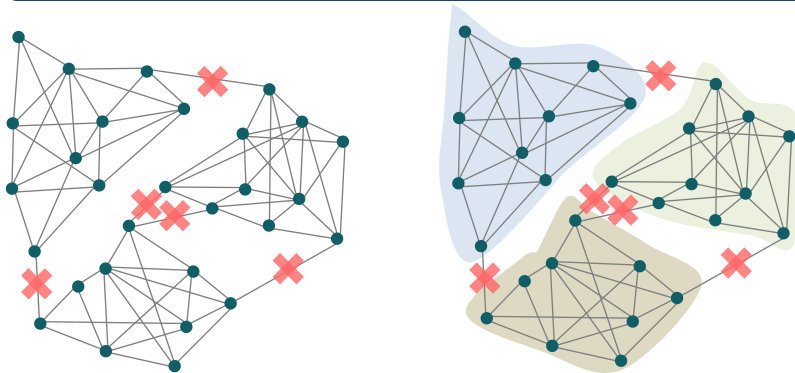
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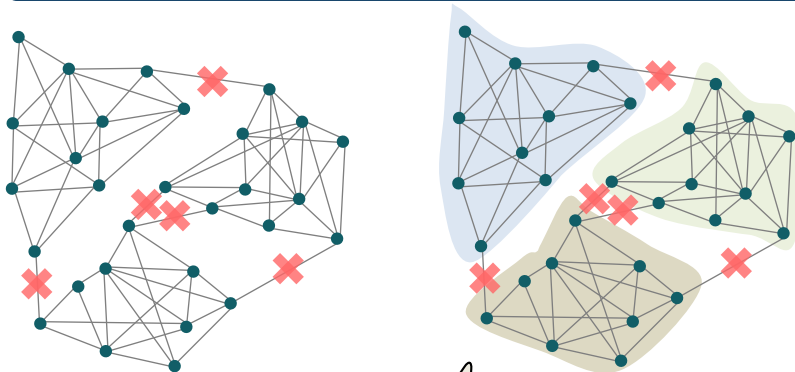
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We could obviously solve this easily using DFS/BFS, but let's see how we can tackle this using the **spectrum of  $L$** !

## Example 2

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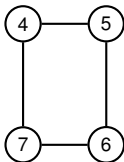
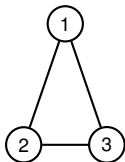


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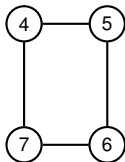
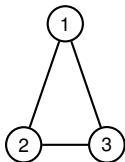


$$\mathbf{A} = \begin{pmatrix} 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 \end{pmatrix}$$

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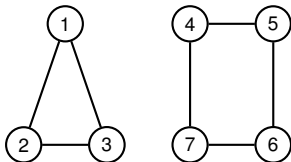
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**Solution:**

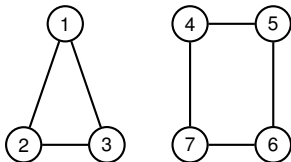
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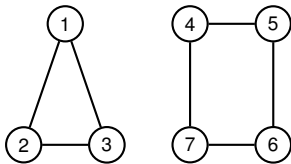
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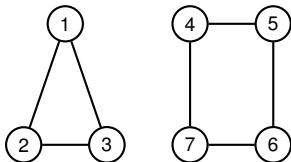
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Next Lecture: A fine-grained approach works even if the clusters are **sparsely** connected!

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- there exist  $f_1, \dots, f_k$  orthonormal such that  $\sum_{\{u,v\} \in E} (f_i(u) - f_i(v))^2 = 0$

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## Proof of Lemma, 2nd statement (non-examinable)

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