

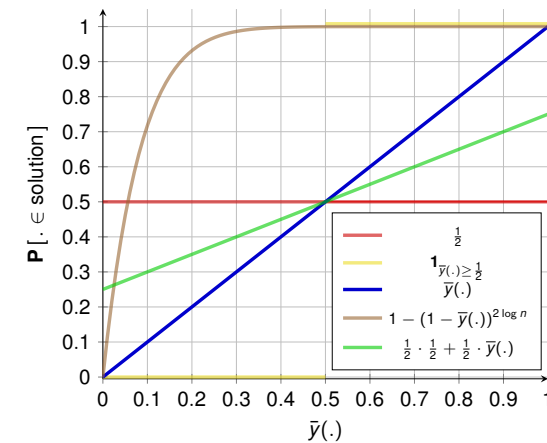
## Randomised Algorithms

### Lecture 11: Spectral Graph Theory

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Lent 2025

## Recap: Different Rounding Procedures



- **MAX-CUT, MAX-3-CNF, MAX-CNF: Uniform Guessing**
- **Vertex-Cover: Deterministic Rounding with Threshold 1/2**
- **Set-Cover (First Try), MAX-CNF: Linear Randomised Rounding**
- **Set-Cover (Final Algorithm): Non-Linear Randomised Rounding**
- **MAX-CNF (Hybrid Algorithm): Guessing + Linear Randomised Rounding**

11. Spectral Graph Theory © T. Sauerwald

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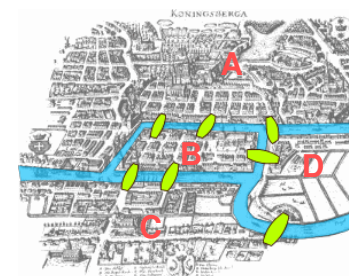
## Outline

Introduction to (Spectral) Graph Theory and Clustering

Matrices, Spectrum and Structure

A Simplified Clustering Problem

## Origin of Graph Theory



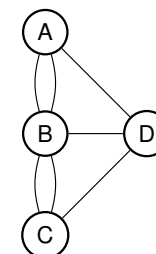
Source: Wikipedia

Seven Bridges at Königsberg 1737

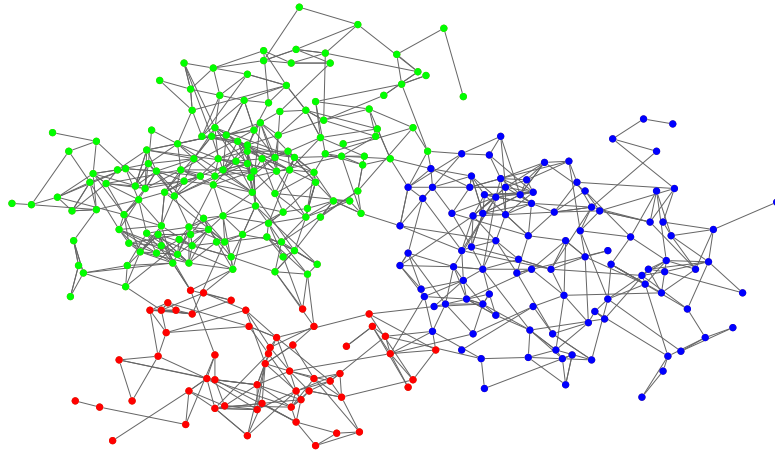


Source: Wikipedia

Leonhard Euler (1707-1783)



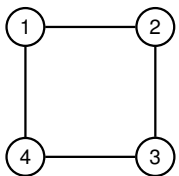
Is there a tour which crosses each bridge **exactly** once?



**Goal:** Use spectrum of graphs (unstructured data) to extract clustering (communities) or other structural information.

- Applications of Graph Clustering
  - Community detection
  - Group webpages according to their topics
  - Find proteins performing the same function within a cell
  - Image segmentation
  - Identify bottlenecks in a network
  - ...
- **Unsupervised** learning method  
(there is no ground truth (usually), and we cannot learn from mistakes!)
- Different formalisations for different applications
  - **Geometric Clustering:** partition points in a Euclidean space
    - $k$ -means,  $k$ -medians,  $k$ -centres, etc.
  - **Graph Clustering:** partition vertices in a graph
    - modularity, **conductance**, min-cut, etc.

### Graphs



- Connectivity
- Bipartiteness
- Number of triangles
- Graph Clustering
- Graph isomorphism
- Maximum Flow
- Shortest Paths
- ...

### Matrices

$$\begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

- Eigenvalues
- Eigenvectors
- Inverse
- Determinant
- Matrix-powers
- ...

Introduction to (Spectral) Graph Theory and Clustering

Matrices, Spectrum and Structure

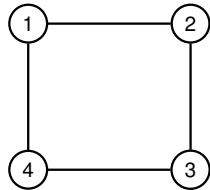
A Simplified Clustering Problem

## Adjacency Matrix

### Adjacency matrix

Let  $G = (V, E)$  be an **undirected** graph. The **adjacency matrix** of  $G$  is the  $n$  by  $n$  matrix  $\mathbf{A}$  defined as

$$\mathbf{A}_{u,v} = \begin{cases} 1 & \text{if } \{u, v\} \in E \\ 0 & \text{otherwise.} \end{cases}$$



$$\mathbf{A} = \begin{pmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{pmatrix}$$

Properties of  $\mathbf{A}$ :

- The sum of elements in each row/column  $i$  equals the **degree** of the corresponding vertex  $i$ ,  $\deg(i)$
- Since  $G$  is **undirected**,  $\mathbf{A}$  is **symmetric**

## Eigenvalues and Graph Spectrum of $\mathbf{A}$

### Eigenvalues and Eigenvectors

Let  $\mathbf{M} \in \mathbb{R}^{n \times n}$ ,  $\lambda \in \mathbb{C}$  is an **eigenvalue** of  $\mathbf{M}$  if and only if there exists  $x \in \mathbb{C}^n \setminus \{\mathbf{0}\}$  such that

$$\mathbf{M}x = \lambda x.$$

We call  $x$  an **eigenvector** of  $\mathbf{M}$  corresponding to the eigenvalue  $\lambda$ .

An **undirected** graph  $G$  is  **$d$ -regular** if every degree is  $d$ , i.e., every vertex has exactly  $d$  connections.

### Graph Spectrum

Let  $\mathbf{A}$  be the adjacency matrix of a  **$d$ -regular** graph  $G$  with  $n$  vertices. Then,  $\mathbf{A}$  has  $n$  **real eigenvalues**  $\lambda_1 \leq \dots \leq \lambda_n$  and  $n$  corresponding **orthonormal eigenvectors**  $f_1, \dots, f_n$ . These eigenvalues associated with their **multiplicities** constitute the **spectrum** of  $G$ .

= **orthogonal** and **normalised**

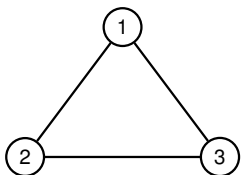
Remark: For **symmetric** matrices we have **algebraic multiplicity** = **geometric multiplicity** (otherwise  $\geq$ )

## Example 1

**Bonus:** Can you find a short-cut to  $\det(\mathbf{A} - \lambda \cdot \mathbf{I})$ ?



**Question:** What are the Eigenvalues and Eigenvectors?



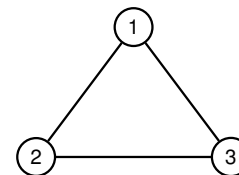
$$\mathbf{A} = \begin{pmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix}$$

## Example 1

**Bonus:** Can you find a short-cut to  $\det(\mathbf{A} - \lambda \cdot \mathbf{I})$ ?



**Question:** What are the Eigenvalues and Eigenvectors?



$$\mathbf{A} = \begin{pmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix}$$

**Solution:**

- The three eigenvalues are  $\lambda_1 = \lambda_2 = -1, \lambda_3 = 2$ .
- The three eigenvectors are (for example):

$$f_1 = \begin{pmatrix} 1 \\ 0 \\ -1 \end{pmatrix}, \quad f_2 = \begin{pmatrix} -\frac{1}{2} \\ 1 \\ -\frac{1}{2} \end{pmatrix}, \quad f_3 = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}.$$

## Laplacian Matrix

Laplacian Matrix

Let  $G = (V, E)$  be a  $d$ -regular undirected graph. The (normalised) Laplacian matrix of  $G$  is the  $n$  by  $n$  matrix  $\mathbf{L}$  defined as

$$\mathbf{L} = \mathbf{I} - \frac{1}{d} \mathbf{A},$$

where  $\mathbf{I}$  is the  $n \times n$  identity matrix.



**Question:** How does the matrix  $\mathbf{I} - \frac{1}{d} \cdot \mathbf{A}$  look like?

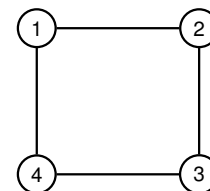
## Laplacian Matrix

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where  $\mathbf{I}$  is the  $n \times n$  identity matrix.



$$\mathbf{L} = \begin{pmatrix} 1 & -1/2 & 0 & -1/2 \\ -1/2 & 1 & -1/2 & 0 \\ 0 & -1/2 & 1 & -1/2 \\ -1/2 & 0 & -1/2 & 1 \end{pmatrix}$$

Properties of  $\mathbf{L}$ :

- The sum of elements in each row/column equals zero
- $\mathbf{L}$  is symmetric

## Relating Spectrum of Adjacency Matrix and Laplacian Matrix

Correspondence between Adjacency and Laplacian Matrix

$\mathbf{A}$  and  $\mathbf{L}$  have the same set of eigenvectors.



**Exercise:** Prove this correspondence. Hint: Use that  $\mathbf{L} = \mathbf{I} - \frac{1}{d} \mathbf{A}$ .  
[Exercise 11/12.1]

## Eigenvalues and Graph Spectrum of $\mathbf{L}$

Eigenvalues and eigenvectors

Let  $\mathbf{M} \in \mathbb{R}^{n \times n}$ ,  $\lambda \in \mathbb{C}$  is an eigenvalue of  $\mathbf{M}$  if and only if there exists  $x \in \mathbb{C}^n \setminus \{0\}$  such that

$$\mathbf{M}x = \lambda x.$$

We call  $x$  an eigenvector of  $\mathbf{M}$  corresponding to the eigenvalue  $\lambda$ .

Graph Spectrum

Let  $\mathbf{L}$  be the Laplacian matrix of a  $d$ -regular graph  $G$  with  $n$  vertices. Then,  $\mathbf{L}$  has  $n$  real eigenvalues  $\lambda_1 \leq \dots \leq \lambda_n$  and  $n$  corresponding orthonormal eigenvectors  $f_1, \dots, f_n$ . These eigenvalues associated with their multiplicities constitute the spectrum of  $G$ .

## Useful Facts of Graph Spectrum

Lemma

Let  $\mathbf{L}$  be the Laplacian matrix of an undirected, regular graph  $G = (V, E)$  with eigenvalues  $\lambda_1 \leq \dots \leq \lambda_n$ .

1.  $\lambda_1 = 0$  with eigenvector  $\mathbf{1}$
2. the multiplicity of the eigenvalue 0 is equal to the number of connected components in  $G$
3.  $\lambda_n \leq 2$
4.  $\lambda_n = 2$  iff there exists a bipartite connected component.

The proof of these properties is based on a powerful characterisation of eigenvalues/vectors!

## A Min-Max Characterisation of Eigenvalues and Eigenvectors

Courant-Fischer Min-Max Formula (non-examinable)

Let  $\mathbf{M}$  be an  $n$  by  $n$  symmetric matrix with eigenvalues  $\lambda_1 \leq \dots \leq \lambda_n$ . Then,

$$\lambda_k = \min_{S: \dim(S)=k} \max_{x \in S, x \neq 0} \frac{x^T \mathbf{M} x}{x^T x},$$

where  $S$  is a subspace of  $\mathbb{R}^n$ . The eigenvectors corresponding to  $\lambda_1, \dots, \lambda_k$  minimise such expression.

$$\lambda_1 = \min_{x \in \mathbb{R}^n \setminus \{0\}} \frac{x^T \mathbf{M} x}{x^T x}$$

minimised by an eigenvector  $f_1$  for  $\lambda_1$

$$\lambda_2 = \min_{\substack{x \in \mathbb{R}^n \setminus \{0\} \\ x \perp f_1}} \frac{x^T \mathbf{M} x}{x^T x}$$

minimised by  $f_2$

## Quadratic Forms of the Laplacian

Lemma

Let  $\mathbf{L}$  be the Laplacian matrix of a  $d$ -regular graph  $G = (V, E)$  with  $n$  vertices. For any  $x \in \mathbb{R}^n$ ,

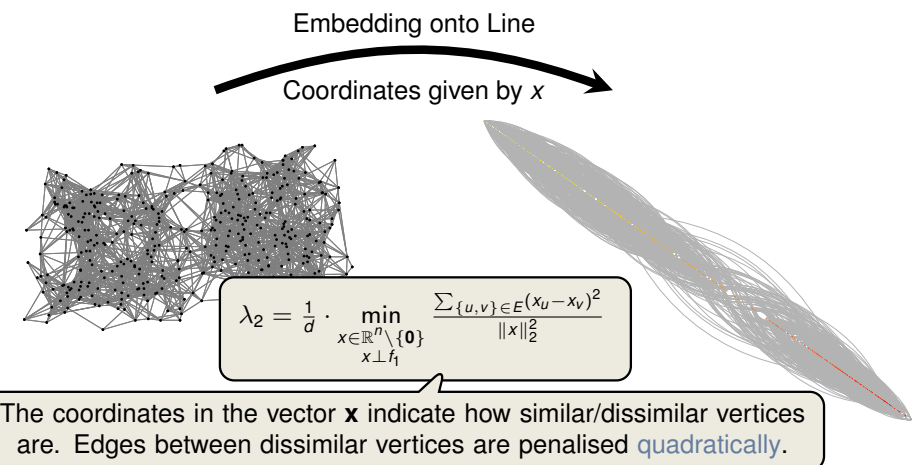
$$x^T \mathbf{L} x = \sum_{\{u,v\} \in E} \frac{(x_u - x_v)^2}{d}.$$

Proof:

$$\begin{aligned} x^T \mathbf{L} x &= x^T \left( \mathbf{I} - \frac{1}{d} \mathbf{A} \right) x = x^T x - \frac{1}{d} x^T \mathbf{A} x \\ &= \sum_{u \in V} x_u^2 - \frac{2}{d} \sum_{\{u,v\} \in E} x_u x_v \\ &= \frac{1}{d} \sum_{\{u,v\} \in E} (x_u^2 + x_v^2 - 2x_u x_v) \\ &= \sum_{\{u,v\} \in E} \frac{(x_u - x_v)^2}{d}. \end{aligned}$$

## Visualising a Graph

**Question:** How can we visualize a complicated object like an unknown graph with many vertices in low-dimensional space?

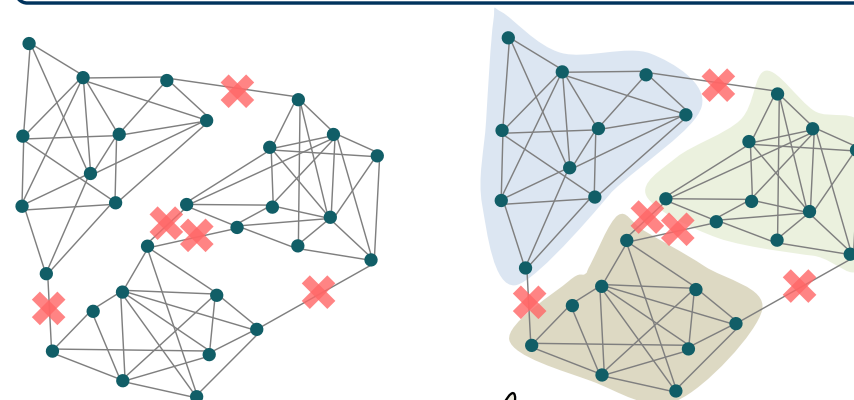


Introduction to (Spectral) Graph Theory and Clustering

Matrices, Spectrum and Structure

A Simplified Clustering Problem

Partition the graph into **connected components** so that any pair of vertices in the same component is connected, but vertices in different components are not.

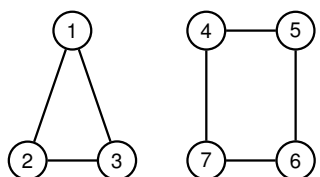


We could obviously solve this easily using DFS/BFS, but let's see how we can tackle this using the **spectrum of L**!

## Example 2



**Question:** What are the Eigenvectors with Eigenvalue 0 of  $L$ ?



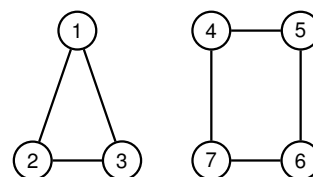
$$A = \begin{pmatrix} 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 \end{pmatrix}$$

$$L = \begin{pmatrix} 1 & -\frac{1}{2} & -\frac{1}{2} & 0 & 0 & 0 & 0 \\ -\frac{1}{2} & 1 & -\frac{1}{2} & 0 & 0 & 0 & 0 \\ -\frac{1}{2} & -\frac{1}{2} & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & -\frac{1}{2} & 0 & -\frac{1}{2} \\ 0 & 0 & 0 & -\frac{1}{2} & 1 & -\frac{1}{2} & 0 \\ 0 & 0 & 0 & 0 & -\frac{1}{2} & 1 & -\frac{1}{2} \\ 0 & 0 & 0 & -\frac{1}{2} & 0 & -\frac{1}{2} & 1 \end{pmatrix}$$

## Example 2



**Question:** What are the Eigenvectors with Eigenvalue 0 of  $L$ ?



$$A = \begin{pmatrix} 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 \end{pmatrix}$$

$$L = \begin{pmatrix} 1 & -\frac{1}{2} & -\frac{1}{2} & 0 & 0 & 0 & 0 \\ -\frac{1}{2} & 1 & -\frac{1}{2} & 0 & 0 & 0 & 0 \\ -\frac{1}{2} & -\frac{1}{2} & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & -\frac{1}{2} & 0 & -\frac{1}{2} \\ 0 & 0 & 0 & -\frac{1}{2} & 1 & -\frac{1}{2} & 0 \\ 0 & 0 & 0 & 0 & -\frac{1}{2} & 1 & -\frac{1}{2} \\ 0 & 0 & 0 & -\frac{1}{2} & 0 & -\frac{1}{2} & 1 \end{pmatrix}$$

**Solution:**

- Two smallest eigenvalues are  $\lambda_1 = \lambda_2 = 0$ .
- The corresponding two eigenvectors are:

$$f_1 = \begin{pmatrix} 1 \\ 1 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}, \quad f_2 = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 1 \\ 1 \end{pmatrix} \quad (\text{or } f_1 = \begin{pmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{pmatrix}, \quad f_2 = \begin{pmatrix} -1/3 \\ -1/3 \\ -1/3 \\ 1/4 \\ 1/4 \\ 1/4 \\ 1/4 \end{pmatrix})$$

Thus we can easily solve the simplified clustering problem by computing the eigenvectors with eigenvalue 0

Next Lecture: A fine-grained approach works even if the clusters are **sparsely** connected!

## Proof of Lemma, 2nd statement (non-examinable)

Let us generalise and formalise the previous example!

Proof (multiplicity of 0 equals the no. of connected components):

1. (" $\implies$ "  $cc(G) \leq \text{mult}(0)$ ). We will show:

$G$  has exactly  $k$  connected comp.  $C_1, \dots, C_k \implies \lambda_1 = \dots = \lambda_k = 0$

- Take  $\chi_{C_i} \in \{0, 1\}^n$  such that  $\chi_{C_i}(u) = \mathbf{1}_{u \in C_i}$  for all  $u \in V$
- Clearly, the  $\chi_{C_i}$ 's are **orthogonal**
- $\chi_{C_i}^T \mathbf{L} \chi_{C_i} = \frac{1}{d} \cdot \sum_{\{u,v\} \in E} (\chi_{C_i}(u) - \chi_{C_i}(v))^2 = 0 \implies \lambda_1 = \dots = \lambda_k = 0$

2. (" $\impliedby$ "  $cc(G) \geq \text{mult}(0)$ ). We will show:

$\lambda_1 = \dots = \lambda_k = 0 \implies G$  has at least  $k$  connected comp.  $C_1, \dots, C_k$

- there exist  $f_1, \dots, f_k$  orthonormal such that  $\sum_{\{u,v\} \in E} (f_i(u) - f_i(v))^2 = 0$
- $\implies f_1, \dots, f_k$  constant on connected components
- as  $f_1, \dots, f_k$  are pairwise orthogonal,  $G$  must have  $k$  different connected components.

□