## Introduction to Probability

Lecture 12: Online Algorithms Mateja Jamnik, <u>Thomas Sauerwald</u>

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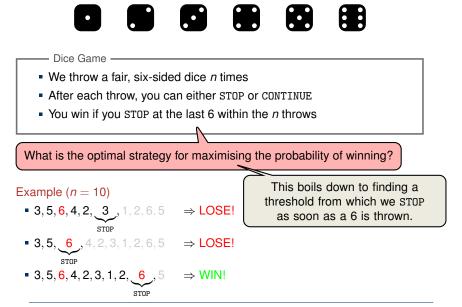


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Stopping Problem 2: The Secretary Problem

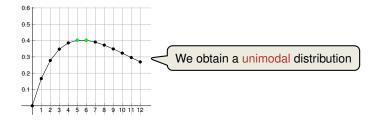
A Generalisation: The Odds Algorithm (non-examinable)

#### Introduction: Dice Game



## **Dice Game (Solution)**

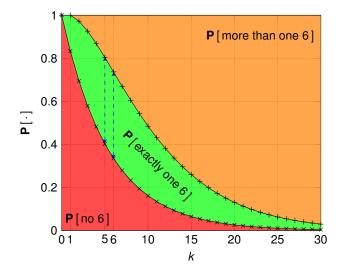
**P** [obtain exactly one 6 in last k throws] = 
$$\binom{k}{1} \cdot \frac{1}{6} \cdot \left(\frac{5}{6}\right)^{k-1} = \frac{k}{6} \cdot \left(\frac{5}{6}\right)^{k-1}$$



- This is maximised for k = 6 (or k = 5)  $\Rightarrow$  best strategy: wait until we have 6 (5) throws left, and then STOP at the first 6
- Probability of success is:

$$\left(\frac{5}{6}\right)^5 \approx 0.40.$$

Intro to Probability



#### Stopping Problem 2: The Secretary Problem

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## The Secretary Problem

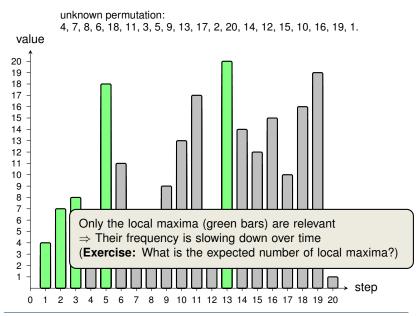
The Problem

- We are interviewing *n* candidates for one job in a sequential, random order
- A candidate must be accepted (STOP) or rejected immediately after the interview and cannot be recalled
- Goal: maximise the probability of hiring the best candidate

also known as marriage problem (Kepler 1613), hiring problem or best choice problem.

Further Remarks -

- After seeing candidate *i*, we only know the relative order among the first *i* candidates.
- ⇒ For our problem we may as well assume that the only information we have when interviewing candidate *i* is whether that candidate is best among {1,...,*i*} or not.



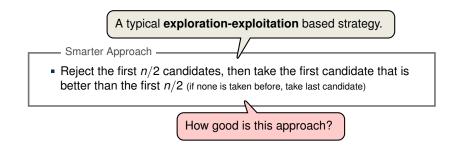
Intro to Probability

Stopping Problem 2: The Secretary Problem

Naive Approach -

- Always pick the first (or any other) candidate
- Probability for success is:

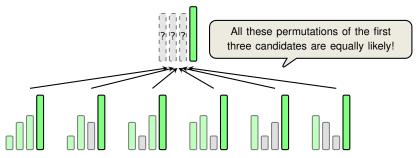
**P**[hire best candidate] = 
$$\frac{1}{n}$$
.



Example 1	
	ound on the success probability of the refined approach st candidate better than the first $n/2$ ).
	Answer

## Finding the Optimal Strategy (1/2)

 Observation 1: At interview *i*, it only matters if current candidate is best so far (i.e., no benefit in counting how many "best-so-far" candidates we had).



 Observation 2: If at interview *i*, the best strategy is to accept the candidate (if it is "best-so-far"), then the same holds for interview *i* + 1

Optimal Strategy

- Explore but reject the first x 1 candidates
- Accept first candidate  $i \ge x$  which is better than all candidates before

Example 2

Find x which maximises the probability of hiring the best candidate.

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Suppose n = 50: P[success] For n = 50,  $x = 50/e \approx 18$ , we obtain a success probability of  $\approx 1/e^{-1}$ 0.5 0.4 0 Х 10 1820 30 40 50

"The Postdoc Variant of the Secretary Problem" (Vanderbei'80) -

- same setup as in the secretary problem before
- difference: we want to pick the second-best ("the best [postdoc] is going to Harvard")
- Success probability of the optimal strategy is:

$$\frac{0.25n^2}{n(n-1)} \quad \stackrel{n \to \infty}{\longrightarrow} \quad \frac{1}{4}$$

Thus it is easier to pick the best than the second-best(!)

Stopping Problem 2: The Secretary Problem

A Generalisation: The Odds Algorithm (non-examinable)

## Details of the Odds Algorithm

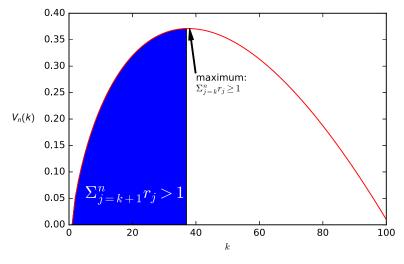
- Let  $I_1, I_2, \ldots, I_n$  be a sequence of independent indicators and let  $p_i = \mathbf{E}[I_i]$ • Let  $r_j := \frac{p_j}{1-p_i}$  (the odds) and  $p_j \in (0, 1)$  for all j = 1, 2, ..., n
- Example 3 What is the probability that after trial k, there is exactly one success?  $\mathbf{P}\left[\sum_{i=k}^{n} l_{i}=1\right] = \sum_{i=k}^{n} p_{i} \cdot \prod_{\substack{k \leq i \leq n, i \neq i}}^{n} (1-p_{i}) = \sum_{i=k}^{n} r_{i} \cdot \left(\prod_{i=k}^{n} (1-p_{i})\right)\right]$ • One can prove that  $\mathbf{P}\left[\sum_{j=k}^{n} I_{j} = 1\right]$  is unimodal in  $k \Rightarrow$  there is an ideal point from which on we should STOP at the first success!

Odds Algorithm ("Sum the Odds to One and Stop", F. Thomas Bruss, 2000)

- 1. Let  $k^*$  be the largest k such that  $\sum_{j=k}^{n} r_j \ge 1$ 2. Ignore everything before the  $k^*$ -th trial, then STOP at the first success.

• The success probability is  $\sum_{j=k^*}^n r_j \cdot (\prod_{i=k^*}^n (1 - p_i))$ .

This algorithm always executes the optimal strategy!



Source: Group Fibonado

Example 4	•
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#### Use the Odds Algorithm to analyse the Secretary Problem.

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Stopping Problem 2: The Secretary Problem

A Generalisation: The Odds Algorithm (non-examinable)

#### Part I: Introduction to Probability

Lecture 1: Conditional probabilities and Bayes' theorem

#### Part II: Random Variables

- Lecture 2: Random variables, probability mass function, expectation
- Lecture 3: Expectation properties, variance, discrete distributions
- Lecture 4: More discrete distributions: Poisson, Geometric, Negative Binomial
- Lecture 5: Continuous random variables
- Lecture 6: Marginals and Joint Distributions
- Lecture 7: Independence, Covariance and Correlation

#### Part III: Moments and Limit Theorems

- Lecture 8: Basic Inequalities and Law of Large Numbers
- Lecture 9: Central Limit Theorem

#### Part IV: Applications and Statistics

- Lecture 10: Estimators (Part I)
- Lecture 11: Estimators (Part II)
- Lecture 12: Online Algorithms

#### Very Important:

- Bernoulli, Binomial, Poisson
- (Continuous) Uniform, Normal, Exponential

#### (Somewhat Less) Important:

Geometric, Negative Binomial, Hypergeometric, Discrete Uniform

#### Not used or not defined in this course (and thus not examinable):

- Cauchy, Gamma, bivariate Normal
- Beta

# Thank you and Best Wishes for the Exam!