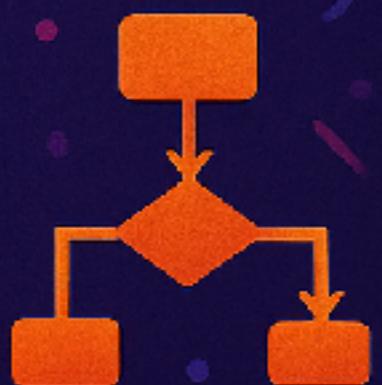
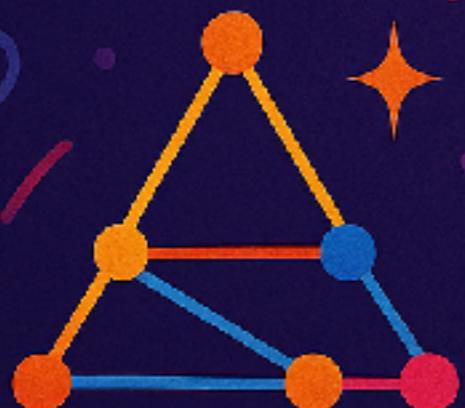


COMPLEXITY THEORY

$P \neq NP$



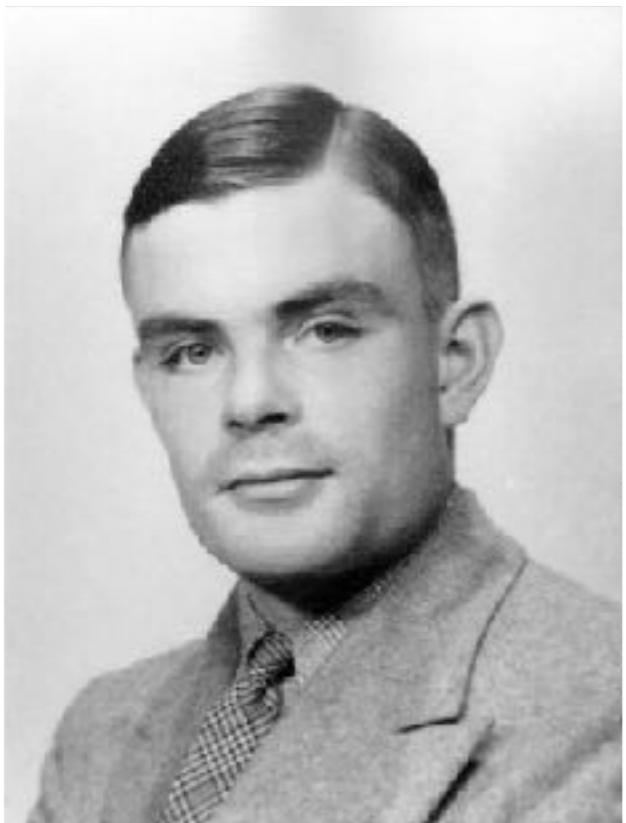
Introduction

Tom Gur

The story begins here in Cambridge

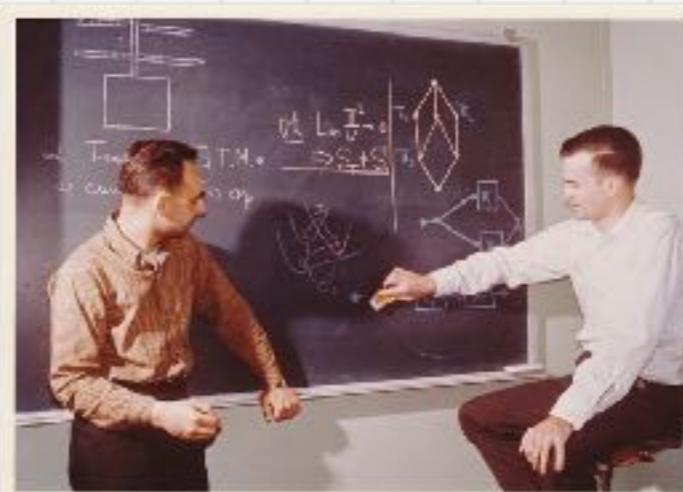


From Computation Theory to Complexity Theory



Finite or infinite, that is the question.

Turing, 1936



Or is it?

Hartmanis and Stearns, 1965

Numbers and scale

10 - Fingers

100 - a full lecture theatre

1000 = 2^{10} - a rock concert

1000000 = 2^{20} - 5x population of Cambridge, 1M seconds = 11.5 days

1 Billion = 2^{30} - 1B seconds = 31.5 years, 1GHz = 1/billion sec, iPhone

1 Trillion = 2^{40} - 1T seconds = 34,842.1 years, FLOPS of PlayStation

2^{50} = 35,678,377.2 years, FLOPS of a supercomputer

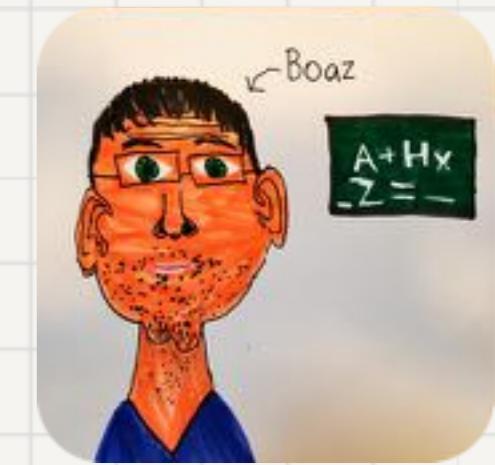
2^{60} - 36 Billion years, 3 times the age of the universe

2^{150} - all super computers working throughout the universe's lifespan...

$2^{1000000}$ complexity of an exponential-time algorithm on a small input...



Ryan O'Donnell
CMU



Boaz Barak
Harvard

What is Complexity Theory

Complexity theory characterises **tractable** computation!

Why care?

Practice: Learn how to recognise intractable problems.

Theory: Understanding the power and limitations of computation,
inspiring new paradigms

Interdisciplinary: Deep connections to physics and mathematics

Bonus: Ways to make progress on big philosophical questions:

randomness, quantum, free will, and beyond



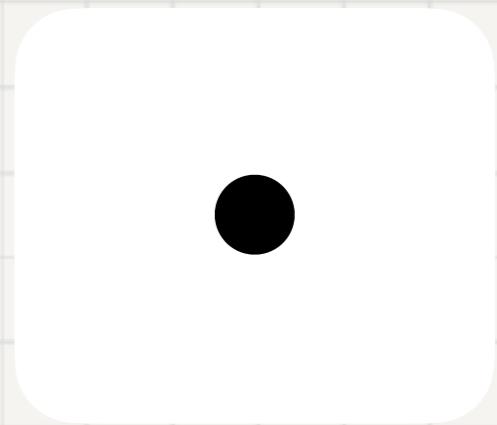
Scott Aaronson
UT Austin

The power of abstraction

Avi Wigderson
Princeton



Without enough abstraction, it is hard to make progress



With excessive abstraction, we lose complexity

Asymptotics

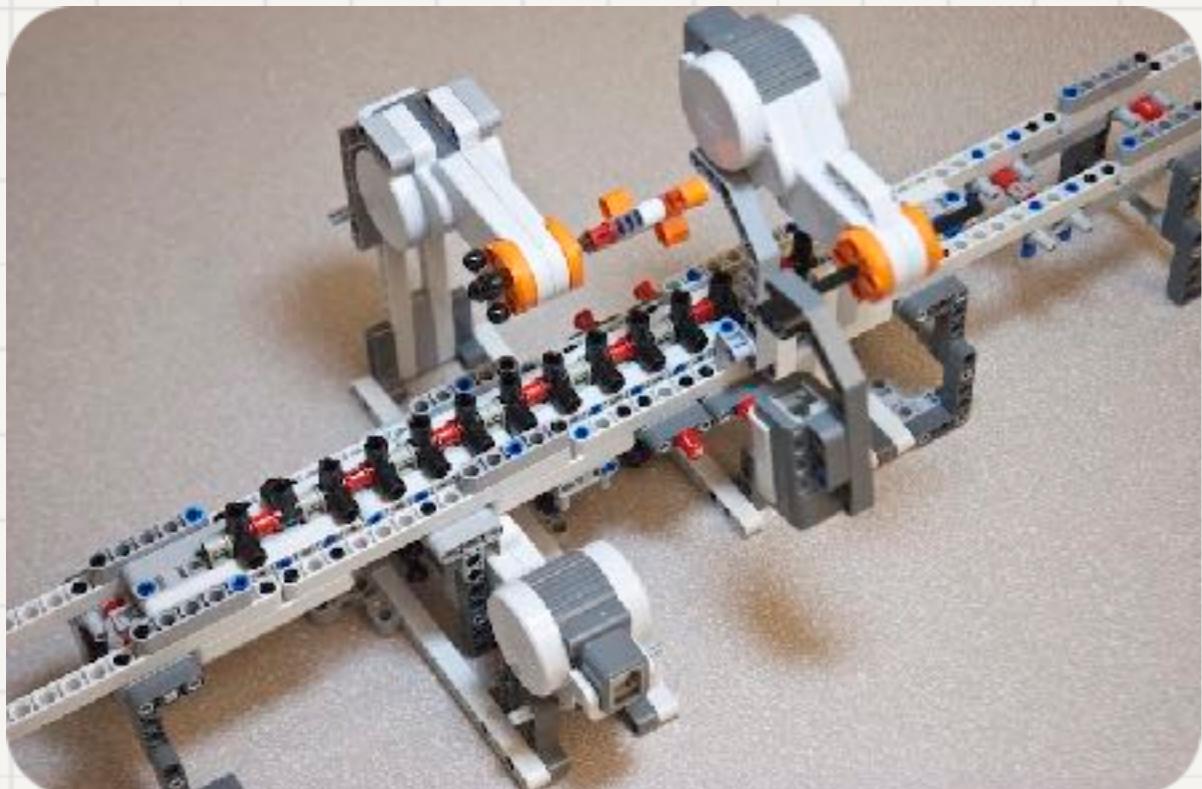
Definition

For functions $f : \mathbb{N} \rightarrow \mathbb{N}$ and $g : \mathbb{N} \rightarrow \mathbb{N}$, we say that:

- $f = O(g)$, if there is an $n_0 \in \mathbb{N}$ and a constant c such that for all $n > n_0$, $f(n) \leq cg(n)$;
- $f = \Omega(g)$, if there is an $n_0 \in \mathbb{N}$ and a constant c such that for all $n > n_0$, $f(n) \geq cg(n)$.
- $f = \theta(g)$ if $f = O(g)$ and $f = \Omega(g)$.

The power of abstraction

Turing Machines



Why TM?

Church-Turing thesis: a function can be computed iff it is TM-computable

Extended Church-Turing thesis: adds efficiency to the picture

(Quantum Computing is the only widely accepted contender to ECT)

- Q — a finite set of states;
- Σ — a finite set of symbols, including \sqcup and \triangleright .
- $s \in Q$ — an initial state;
- $\delta : (Q \times \Sigma) \rightarrow (Q \cup \{\text{acc}, \text{rej}\}) \times \Sigma \times \{L, R, S\}$

$$(q, w, u) \xrightarrow{M} (q', w', u')$$

if

- $w = va$;
- $\delta(q, a) = (q', b, D)$; and
- either $D = L$ and $w' = v$ and $u' = bu$
or $D = S$ and $w' = vb$ and $u' = u$
or $D = R$ and $w' = vbc$ and $u' = x$, where $u = cx$. If u is empty, then $w' = vb\sqcup$ and u' is empty.

The language $L(M) \subseteq \Sigma^*$ accepted by the machine M is the set of strings

$$\{x \mid (s, \triangleright, x) \xrightarrow{M}^* (\text{acc}, w, u) \text{ for some } w \text{ and } u\}$$

Turing fact of the day

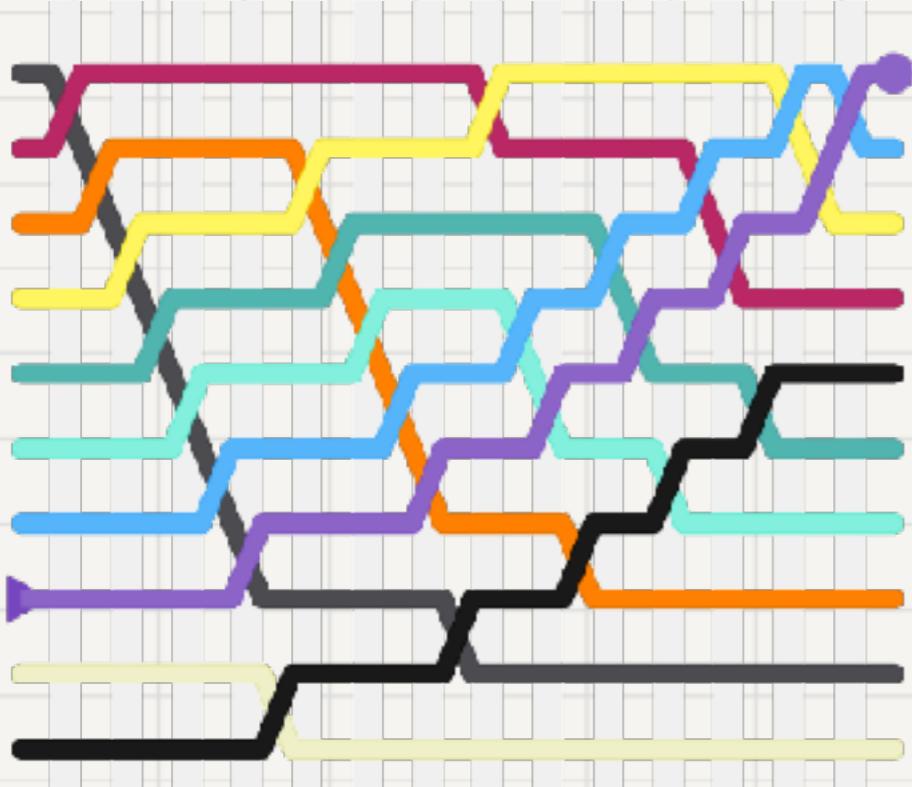


Ran a marathon in 2:46:03.

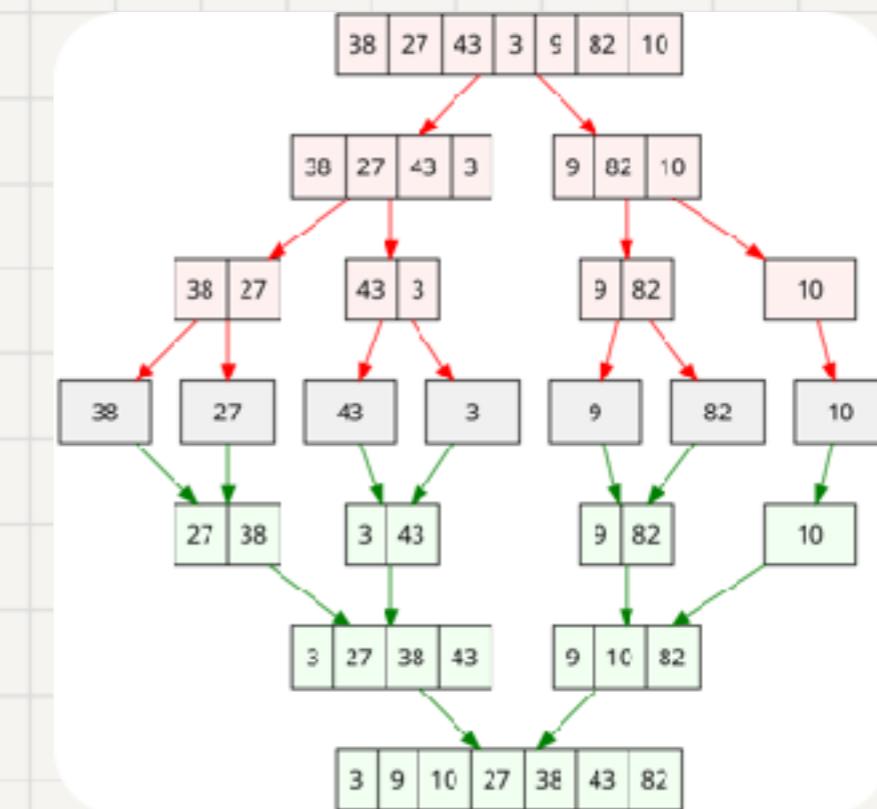
Considered representing England in the 1948 London Olympics.

Algorithms & Problems

Bubble Sort runs in time $O(n^2)$



Merge Sort in $O(n \log n)$



We compare the two algorithms as they solve the same problem.

But what is the complexity of the **sorting problem**?

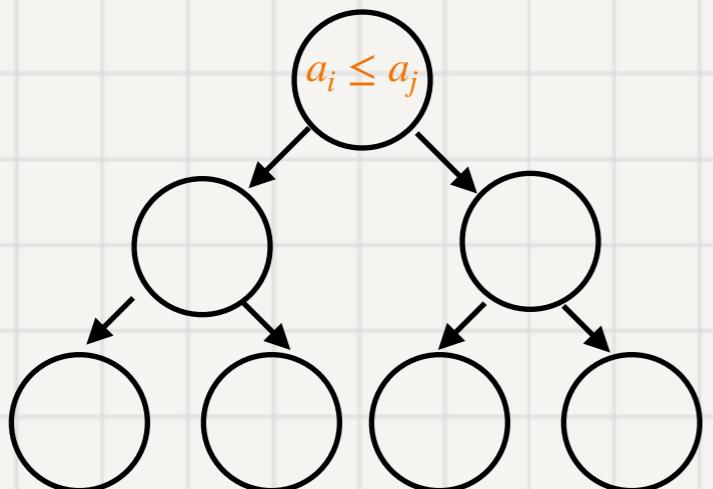
Upper bounds & lower bounds

What is the time **complexity** of the sorting problem?

Merge Sort gives us an **upper bound** of $O(n \log n)$

Showing a **lower bound** requires arguing about all algorithms

Further abstraction: Decision Trees



A comparison algorithm sorts **n** numbers

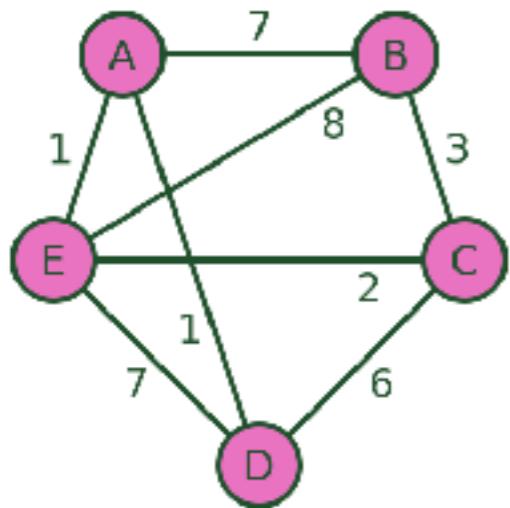
a_1, \dots, a_n by comparing pairs

There are **$n!$** permutations, hence

tree height $\geq \log(n!) = \Omega(n \log n)$

Bonus: break the abstraction!

Travelling Salesman Problem (TSP)



Given a weighted graph with vertices v_1, \dots, v_n ,

Is there an ordering $\sigma \in S_n$ s.t.

$$\sum_{i=1}^n w(\sigma(i), \sigma(i+1)) \leq \tau$$

Complexity of brute force: $O(n!)$

Best upper bound: $O(n^2 2^n)$

Best lower bound*: $\Omega(n)$

Between these two is the

chasm of our ignorance

About the course

Textbooks

“Computational Complexity: A Modern Approach” by Arora and Barak

“Computational Complexity: A Conceptual Perspective” by Oded Goldreich

Exam

Everything covered in the lectures

How to ace Complexity Theory

Internalising ideas vs problem solving

Participating in class

Active learning

Anonymous feedback

Anonymous Feedback

Complexity Theory, Cambridge

tg508@cam.ac.uk [Switch accounts](#) 

✉ Not shared

Please feel free to leave any comments, suggestions, and requests. If things are going well, a good word is always appreciated. If you have ideas on improving the course, please let me know (in a kind and respectful way). I hope you enjoy the course!

Your answer

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<https://www.cl.cam.ac.uk/teaching/2425/Complexity/materials.html>