University of Cambridge 2024/25 Part II / Part III / MPhil ACS *Category Theory* Exercise Sheet 3 by Andrew Pitts and Marcelo Fiore

CCCs

- 1. Show that the category **Preord** of preorders and monotone functions is a cartesian closed category.
- 2. Show that for any objects X and Y in a cartesian closed category C, there are functions

$$f \in \mathcal{C}(X, Y) \mapsto \ulcorner f \urcorner \in \mathcal{C}(1, Y^X)$$
$$g \in \mathcal{C}(1, Y^X) \mapsto \overline{g} \in \mathcal{C}(X, Y)$$

that give a bijection between the set C(X, Y) of C-morphisms from X to Y and the set $C(1, Y^X)$ of global elements of the exponential Y^X .

- 3. Show that for any objects X and Y in a cartesian closed category C, the morphism app : $Y^X \times X \rightarrow Y$ satisfies $cur(app) = id_{Y^X}$.
- 4. Suppose $f: Y \times X \to Z$ and $g: W \to Y$ are morphisms in a cartesian closed category C. Prove that

$$\operatorname{cur}(f \circ (g \times \operatorname{id}_X)) = (\operatorname{cur} f) \circ g \in \operatorname{C}(W, Z^{\mathsf{A}})$$

5. Let C be a cartesian closed category. For each C-object X and C-morphism $f : Y \rightarrow Z$, define

$$f^X \triangleq \operatorname{cur}(Y^X \times X \xrightarrow{\operatorname{app}} Y \xrightarrow{J} Z) \in \mathbf{C}(Y^X, Z^X)$$

- (a) Prove that $(id_Y)^X = id_{Y^X}$.
- (b) Given $f \in C(Y \times X, Z)$ and $g \in C(Z, W)$, prove that

$$\operatorname{cur}(g \circ f) = g^X \circ \operatorname{cur} f \in \operatorname{C}(Y, W^X)$$

- (c) Deduce that if $u \in C(Y, Z)$ and $v \in C(Z, W)$, then $(v \circ u)^X = v^X \circ u^X \in C(Y^X, W^X)$.
- 6. Let C be a cartesian closed category. For each C-object X and C-morphism $f : Y \rightarrow Z$, define

$$X^{f} \triangleq \operatorname{cur}(X^{Z} \times Y \xrightarrow{\operatorname{id} \times f} X^{Z} \times Z \xrightarrow{\operatorname{app}} X) \in \operatorname{C}(X^{Z}, X^{Y})$$

- (a) Prove that $X^{id_Y} = id_{X^Y}$.
- (b) Given $g \in C(W, X)$ and $f \in C(Y \times X, Z)$, prove that

$$\operatorname{cur}(f \circ (\operatorname{id}_Y \times g)) = Z^g \circ \operatorname{cur} f \in \mathcal{C}(Y, Z^W)$$

(c) Deduce that if $u \in C(Y, Z)$ and $v \in C(Z, W)$, then $X^{(v \circ u)} = X^u \circ X^v \in C(X^W, X^Y)$.

7. For $f : B \to A$ and $g : X \to Y$ in a cartesian closed category show that

$$X^A \xrightarrow{X^f} X^B \xrightarrow{g^B} Y^B = X^A \xrightarrow{g^A} Y^A \xrightarrow{Y^f} Y^B$$

Show further that these composites equal

$$g^f \triangleq \operatorname{cur}(g \circ \operatorname{app} \circ (\operatorname{id}_{X^A} \times f)) : X^A \to Y^B$$

- 8. Let C be a cartesian closed category in which every pair of objects X and Y possesses a binary coproduct $X \xrightarrow{inl_{X,Y}} X + Y \xleftarrow{inr_{X,Y}} Y$. For all objects $X, Y, Z \in C$ construct an isomorphism $(Y \times X) + (Z \times X) \longrightarrow (Y + Z) \times X$.
- 9. (a) State, without proof, what the product in **Set**^{op} of two objects is.
 - (b) Show by example that there are objects X and Y in Set^{op} for which there is no exponential and hence that Set^{op} is not a cartesian closed category.

IPL and STLC

- 1. Using the natural deduction rules for Intuitionistic Propositional Logic, give proofs of the following judgements. In each case write down a corresponding typing judgement of the Simply Typed Lambda Calculus.
 - $\begin{array}{ll} (a) \ \diamond, \psi \vdash (\varphi \Rightarrow \psi) \Rightarrow \psi \\ (b) \ \diamond, \varphi \vdash (\varphi \Rightarrow \psi) \Rightarrow \psi \\ (c) \ \diamond, ((\varphi \Rightarrow \psi) \Rightarrow \psi) \Rightarrow \psi \vdash \varphi \Rightarrow \psi \end{array}$
- 2. (a) Given simple types *A*, *B*, *C*, give terms *s* and *t* of the Simply Typed Lambda Calculus that satisfy the following typing and $\beta\eta$ -equality judgements:

$$\diamond, x : (A \times B) \to C \vdash s : A \to (B \to C)$$

$$\diamond, y : A \to (B \to C) \vdash t : (A \times B) \to C$$

$$\diamond, x : (A \times B) \to C \vdash t[s/y] =_{\beta\eta} x : (A \times B) \to C$$

$$\diamond, y : A \to (B \to C) \vdash s[t/x] =_{\beta\eta} y : A \to (B \to C)$$

(b) Explain why the above implies that for any three objects *X*, *Y*, *Z* in a cartesian closed category C, there are morphisms

$$f: Z^{(X \times Y)} \to (Z^Y)^X$$
$$q: (Z^Y)^X \to Z^{(X \times Y)}$$

that give an isomorphism $Z^{(X \times Y)} \cong (Z^Y)^X$ in C.

- 3. Make up and solve a question like the above ending with an isomorphism $X^1 \cong X$ for any object X in a cartesian closed category with terminal object 1.
- 4. Given types A', A, B, B' in Simply Typed Lambda Calculus, give a term t satisfying

$$\diamond \vdash t : (A' \to A) \to (B \to B') \to (A \to B) \to (A' \to B')$$

If the semantics in a cartesian closed category of A', A, B and B' are the objects X', X, Y, Y' respectively, what is the semantics of t?