

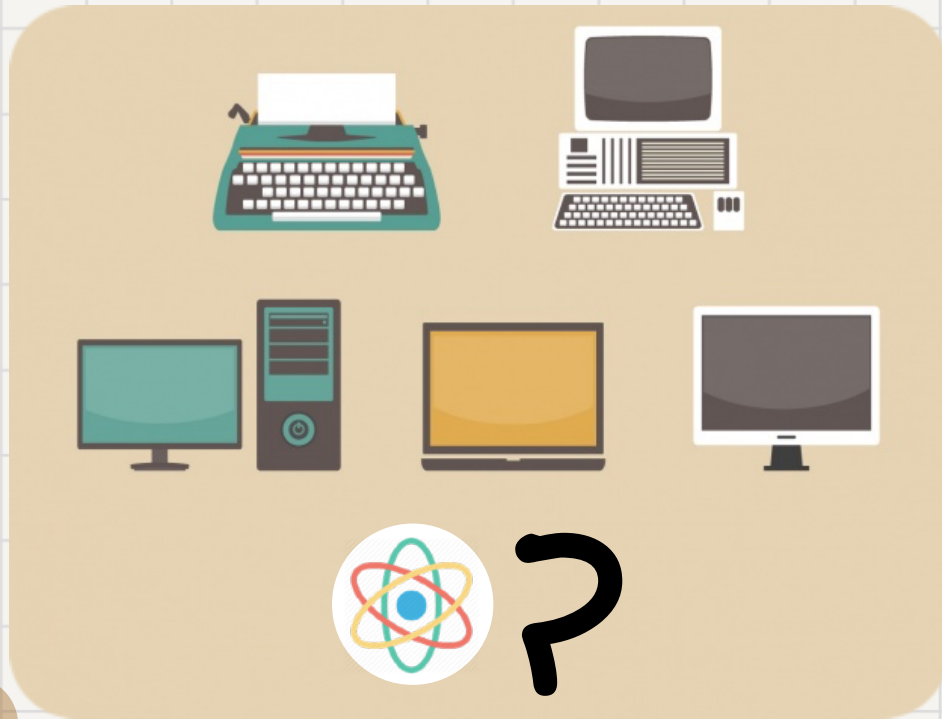
Complexity Theory



Lecture 12

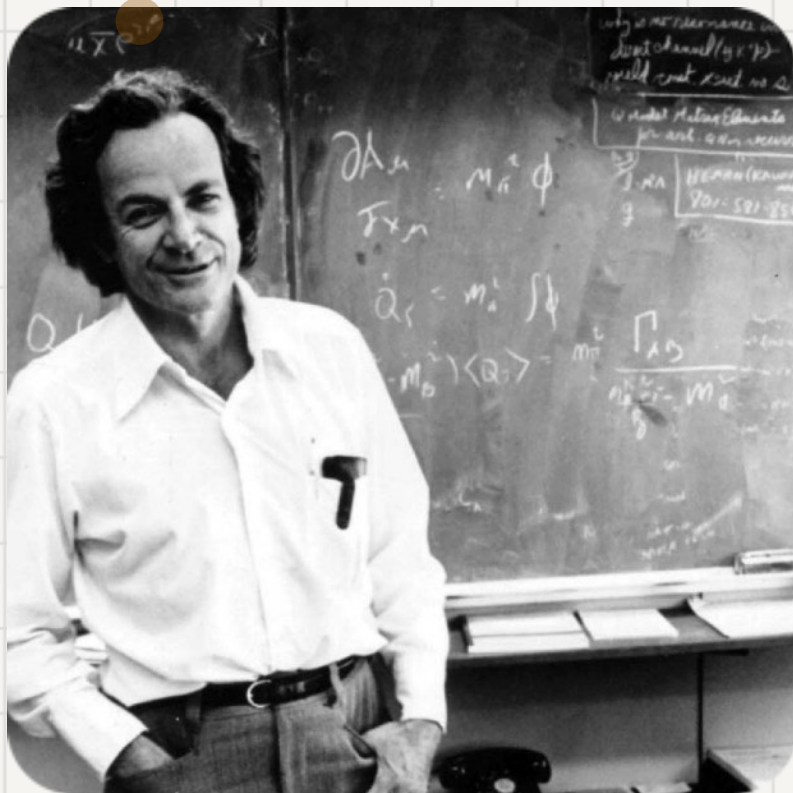
Quantum Complexity

What is quantum computing



The big idea:

Computers that rely on the
Powerful, but *unintuitive*,
principles of quantum mechanics

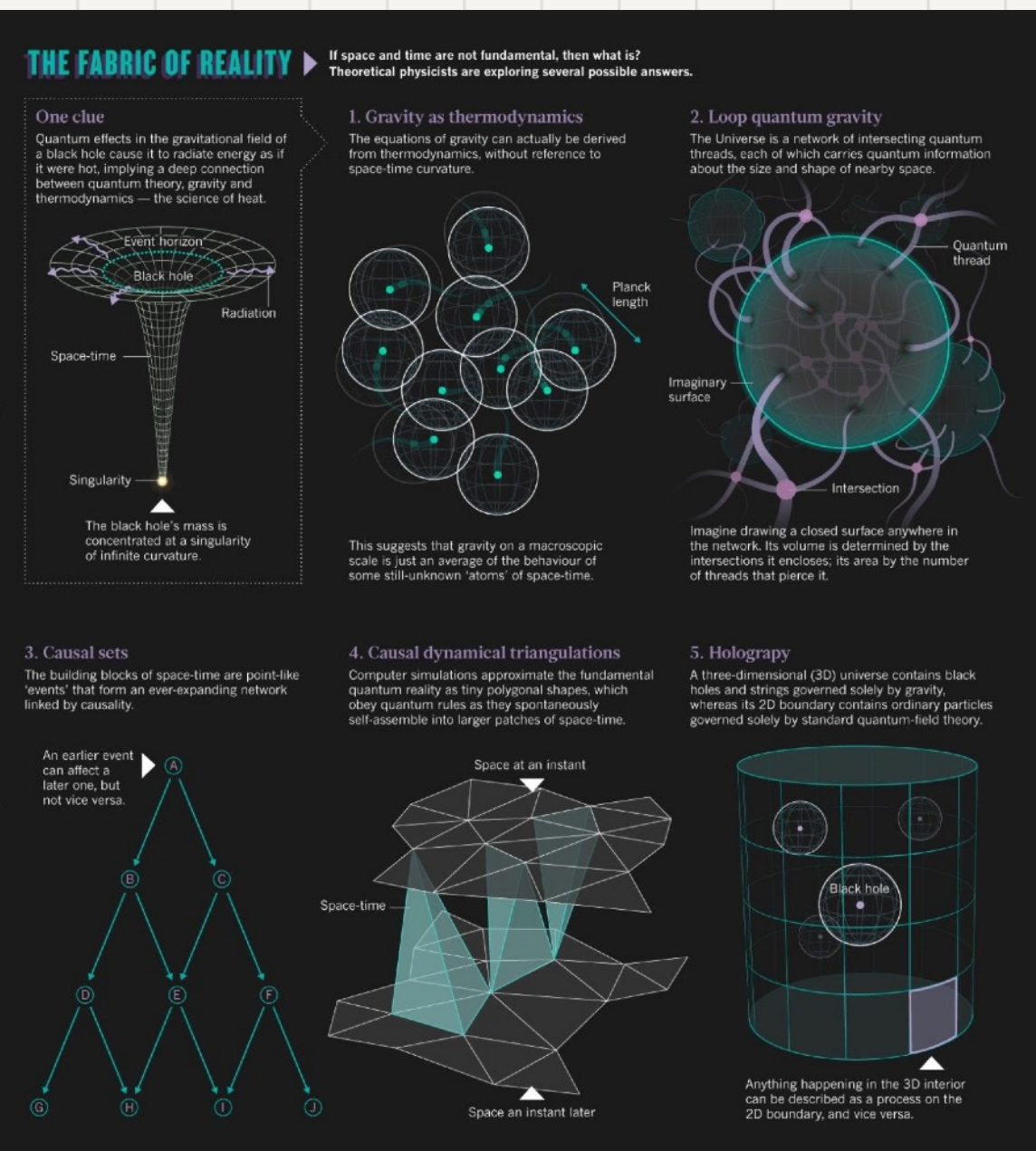


Quantum computing is... nothing less than a distinctively new way of harnessing nature... it will be the first technology that allows useful tasks to be performed in collaboration between parallel universes.

—Fabric of reality



David Deutsch



"if you teach an introductory course on quantum mechanics, and the students don't have nightmares for weeks, tear their hair out, wander around with bloodshot eyes, etc., then you probably didn't get the point across."

Time travel

Parallel universes

Teleportation

Cloning

Faster-than-light
communication

Determinism

Chaos

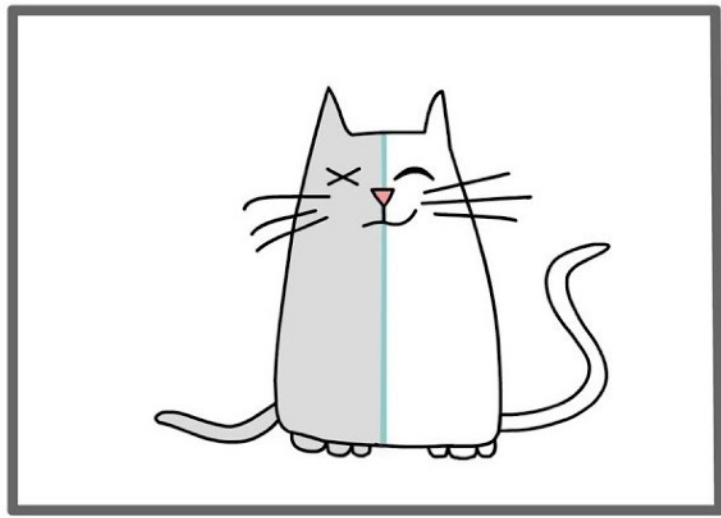
QUANTUM COMPUTING SINCE DEMOCRITUS



SCOTT AARONSON

Superposition, parallel worlds, and cats

Schrödinger's Cat

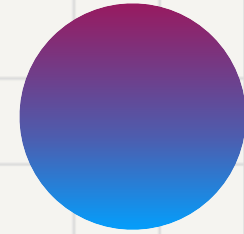


Bit

1 ●

0 ●

Qubit



$$\alpha|0\rangle + \beta|1\rangle$$

$$\frac{1}{\sqrt{2}}|\text{cat}\rangle + \frac{1}{\sqrt{2}}|\text{dead cat}\rangle$$



$2^{20} = 1,048,576$
configurations at once

A crash course on quantum computing

Quantum mechanics is a \mathbb{C} probability theory

Fully captured by 4 postulates:

① Superposition

A quantum state

is a vector $v \in \mathbb{C}^n$

s.t. $\|v\|_2^2 = 1$.

Ex: a qubit $\begin{pmatrix} \alpha \\ \beta \end{pmatrix}$, $|\alpha|^2 + |\beta|^2 = 1$

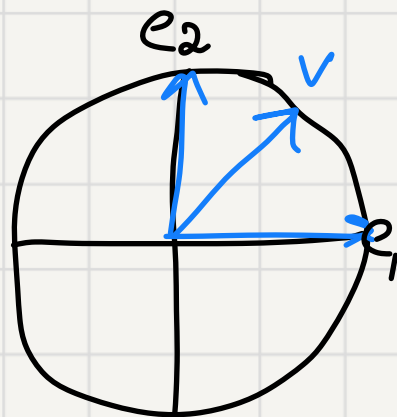
② Measurement

Measuring state $(\alpha_1, \dots, \alpha_n)$ in

the computational basis

collapses it to e_i

w.p. $|\alpha_i|^2$



Mathematical abstraction of quantum computing

III Evolution

A quantum state $v \in \mathbb{C}^n$

evolves to $v' \in \mathbb{C}^n$

via a unitary map

$$\begin{pmatrix} | \\ | \\ v' \\ | \end{pmatrix} = \begin{pmatrix} | \\ | \\ U \\ | \end{pmatrix} \cdot \begin{pmatrix} | \\ | \\ v \\ | \end{pmatrix}$$

IV Entanglement

States $u, v \in \mathbb{C}^n$ are

composed via the tensor

product $u \otimes v$.

Ex: $\begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} =$

Entangled states

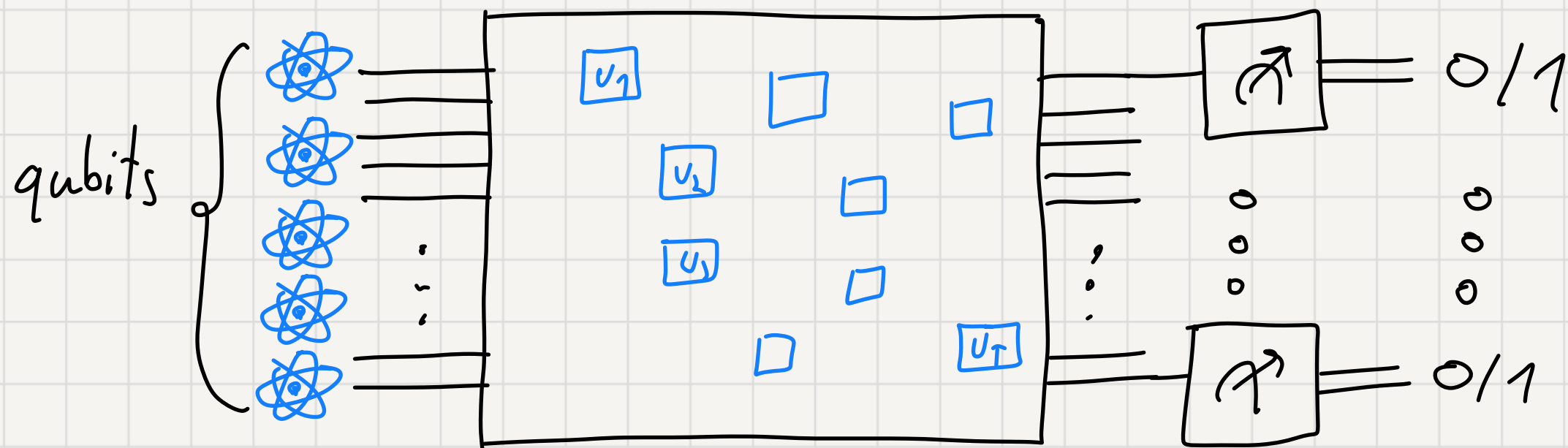
are not of that form!

$$\begin{pmatrix} 1 & 000 \\ 0 & 001 \\ 0 & 010 \\ 0 & 011 \\ 0 & 100 \\ 0 & 101 \\ 0 & 110 \\ 0 & 111 \end{pmatrix}$$

Quantum algorithms

Quantum Turing machines ...are not very nice to work with...

Instead, we typically work with quantum circuits.



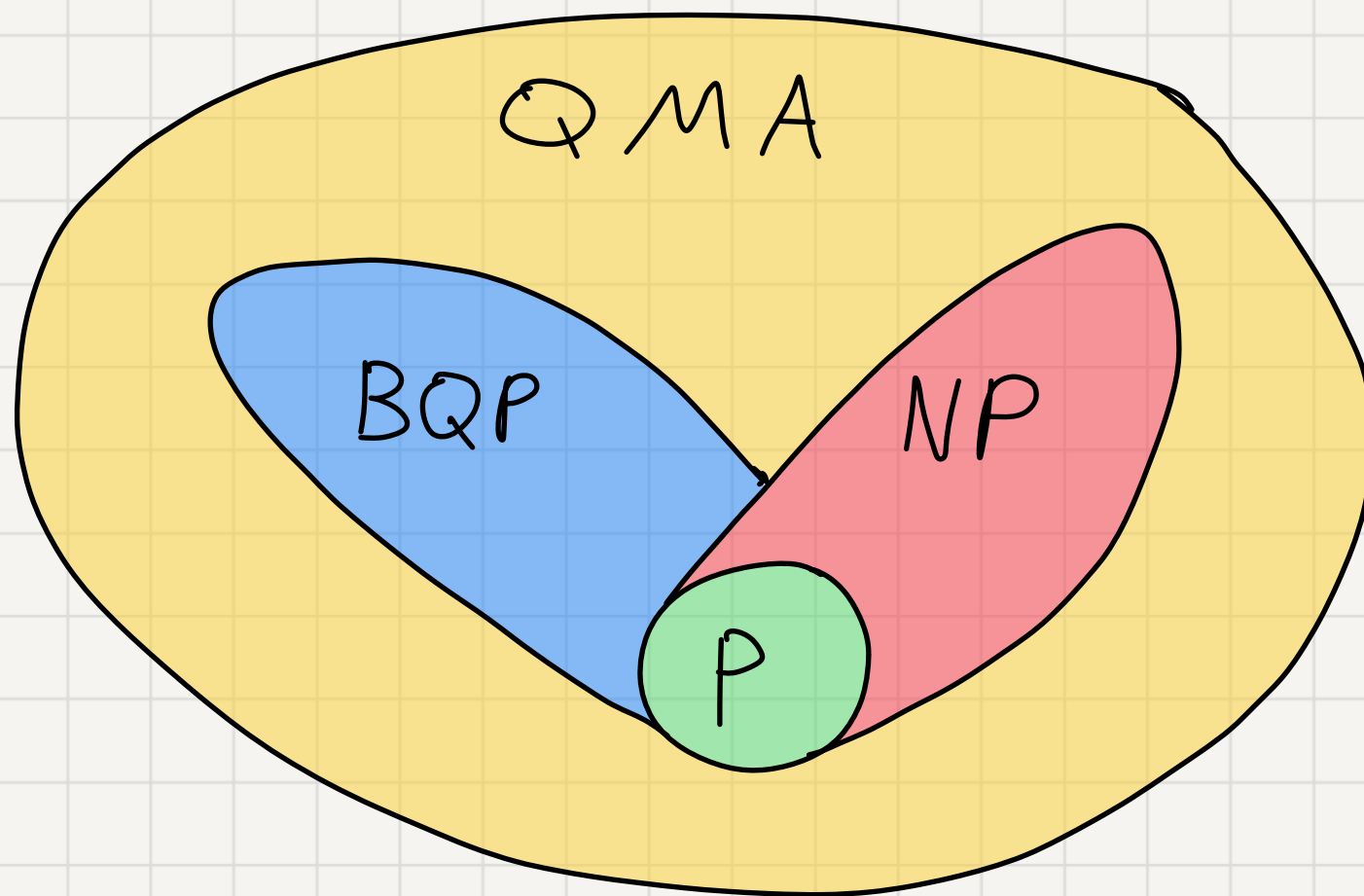
We apply quantum gates (unitary maps) to qubits.

Measure at the end to get classical bits.

Quantum complexity classes

BQP = "quantum P"

QMA = "quantum NP"



BQP

The set of all problems solvable by a
poly-time uniform quantum circuits
 $(C_n)_{n \in \mathbb{N}}$ of polynomial size, w.p. $\geq 2/3$
(i.e., $\forall x \in \{0,1\}^n \Pr[C_n(x) = \mathbb{1}_L(x)] \geq 2/3$)

$T(n)$ -uniformity: C_n can be generated
in $T(n)$ time

Circuit size: #gates \rightarrow time complexity

Error reduction

Let A be a q -algo for computing f such that $\Pr[A(x) = f(x)] \geq 2/3 \quad \forall x$.

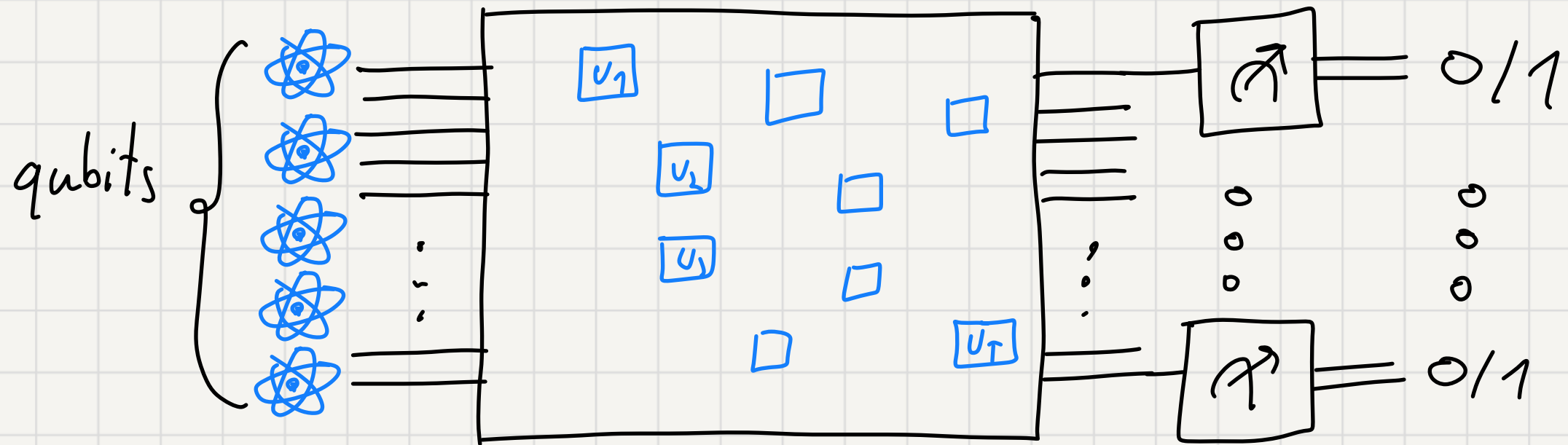
$1/3$ error prob. can be reduced to ϵ !

Repeat A : $A_1, A_2, \dots, A_{\underbrace{\log(1/\epsilon)}_+}$, rule by Maj.

Chernoff bound:

$$\Pr\left[\left|\frac{\sum_{i \in [t]} A_i}{t} - \mathbb{E}[X_i]\right| \leq \frac{1}{\delta}\right] \leq \exp(-t)$$

Quantum algorithms



3 examples where quantum algorithms excel:

- ① Finding sub-group structure (Shor's factoring)
- ② Rapid mixing of Markov chains (Grover's search)
- ③ Computing Fourier Transforms (QFT)

Factoring

Given $n \in \mathbb{N}$, output primes p_1, \dots, p_n s.t.

$$n = p_1 \cdot p_2 \cdot \dots \cdot p_n$$

Decision problem $\text{Factor}(n, k) = 1$

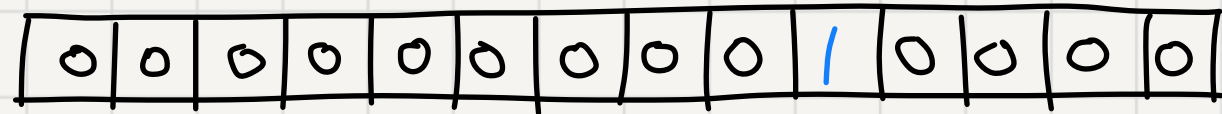
iff n has a prime factor $\leq k$

Shor's algorithm $\text{Factor} \in \text{BQP}$

We know $\text{Factor} \in \text{NP} \cap \text{coNP}$.

Grover's search

Given a string $x \in \{0,1\}^n$, output $i \in [n]$ such that $x_i = 1$



Classical complexity? $\Omega(n)$

Quantum complexity $\Omega(\sqrt{n})$

Quantum Fourier Transform

Given $(f_1, f_2, \dots, f_N) \in \mathbb{C}^N$, output the DFT $(\hat{f}_1, \hat{f}_2, \dots, \hat{f}_N)$

Classical complexity? $O(N \log N)$ Quantum complexity $\tilde{O}(\log N)$

Ask Me Anything