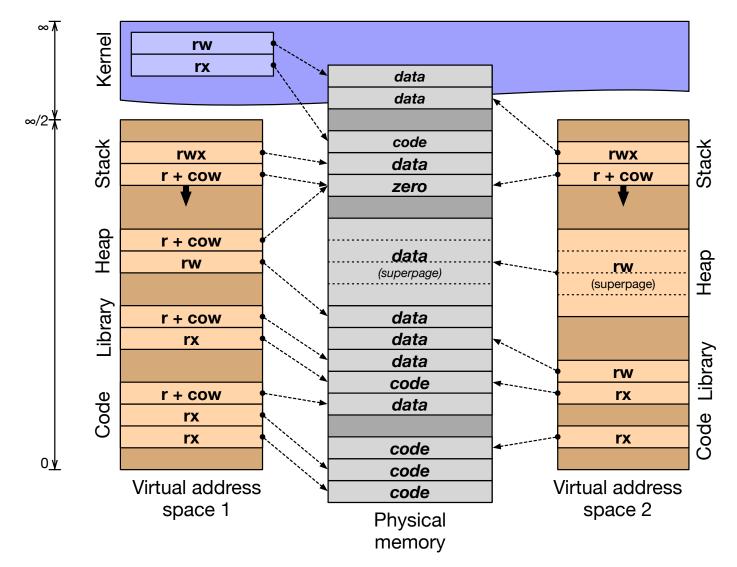
The Process Model (2)

Lecture 4, Part 3: More on VM

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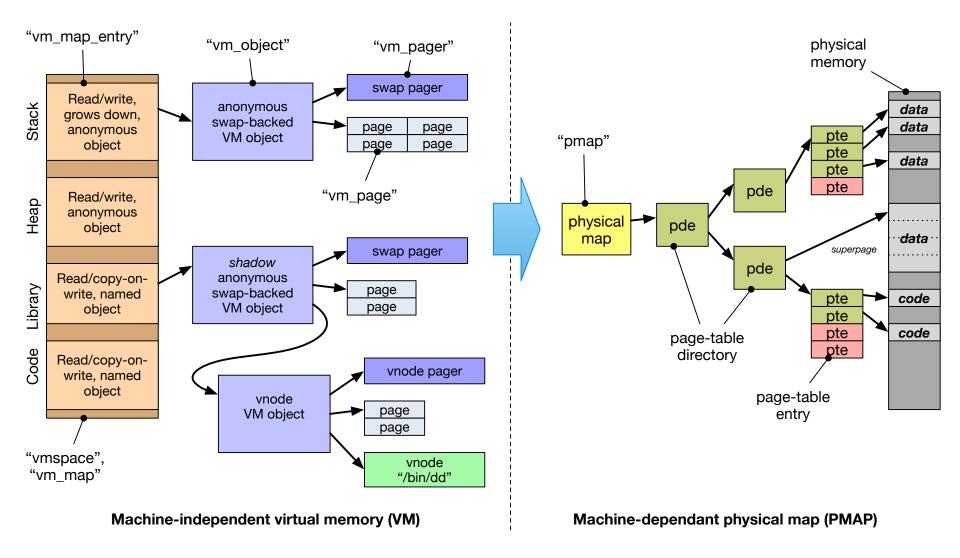
Virtual memory (quick, painful)



So: back to Virtual Memory (VM)

- The process model's isolation guarantees incur real expense
- The VM subsystem works quite hard to avoid expense
 - Shared memory, copy-on-write, page flipping
 - Background page zeroing
 - Superpages to improve TLB efficiency
- VM avoids work, but also manages memory footprint
 - Memory as a cache of secondary storage (files, swap)
 - Demand paging vs. I/O clustering
 - LRU / preemptive swapping to maintain free-page pool
 - Recently: memory compression and deduplication
- These ideas were known before Mach, but...
 - Acetta, et al. impose principled design, turn them into an art form
 - Provide a model beyond V→P mappings in page tables
 - And ideas such as the **message-passing—shared-memory duality**

Kernel programmer view of VM



Mach VM in other operating systems

- Mach: VM mappings, objects, pages, etc., are first-class kernel services exposed via system calls
- In two directly derived systems, quite different stories:
- Mac OS XAlthough not a microkernel, Mach's VM/IPC ApplicationProgramming Interfaces (APIs) are available to user programs,
and widely used for IPC, debugging, ...
- FreeBSDMach VM is used as a foundation for UNIX APIs, but isavailable for use only as a Kernel Programming Interface (KPI)
- In FreeBSD, Mach is used:
 - To efficiently implement UNIX's fork() and execve()
 - For memory-management APIs e.g., mmap() and mprotect()
 - By VM-optimised IPC e.g., pipe() and sendfile()
 - By the filesystem to implement a merged VM-buffer cache
 - By **device drivers** that manage memory in interesting ways (e.g., GPU drivers mapping pages into user processes)
 - By a set of VM worker threads, such as the page daemon, swapper, syncer, and page-zeroing thread