Complexity Theory

Easter 2023 Suggested Exercises 3

- 1. Show that a language L is in co-NP if, and only if, there is a nondeterministic Turing machine M and a polynomial p such that M halts in time p(n) for all inputs of length x, and L is exactly the set of strings x such that all computations of M on input x end in an accepting state.
- 2. Define a *strong* nondeterministic Turing machine as one where each computation has three possible outcomes: accept, reject or maybe. If M is such a machine, we say that it accepts L, if for every $x \in L$, every computation path of M on x ends in either accept or maybe, with at least one accept and for $x \notin L$, every computation path of M on x ends in reject or maybe, with at least one such path ending in reject.

Show that if L is decided by a strong nondeterministic Turing machine running in polynomial time, then $L \in \mathsf{NP} \cap \mathsf{co}\text{-}\mathsf{NP}$.

- 3. We saw in the lectures that if there is a one-way function, then there is a language L in UP that is not in P. Suppose that the RSA function described in the lecture notes (page 38) is a one-way function. What is the language L that can then be proved to be in UP \ P?
- 4. Consider the algorithm presented in the lecture which establishes that Reachability is in $\mathsf{SPACE}((\log n)^2)$. What is the time complexity of this algorithm?

Can you generalise the time bound to the entire complexity class? That is, give a class of functions F, such that

$$\mathsf{SPACE}((\log n)^2) \subseteq \bigcup_{f \in F} \mathsf{TIME}(f)$$

5. Show that, for every nondeterministic machine M which uses $O(\log n)$ work space, there is a machine R with three tapes (input, work and output) which works as follows. On input x, R produces on its output tape a description of the configuration graph for M, x, and R uses $O(\log |x|)$ space on its work tape.

Explain why this means that if Reachability is in L, then L = NL.

6. Consider the language L in the alphabet $\{a,b\}$ given by $L = \{a^nb^n \mid n \in \mathbb{N}\}$. $L \notin \mathsf{SPACE}(c)$ for any constant c. Why?

- 7. On page 42 of the notes, a number of functions are listed as being constructible. Show that this is the case by giving, for each one, a description of an appropriate Turing machine.
 - Prove that if f and g are constructible functions and $f(n) \ge n$, then so are f(g), f + g, $f \cdot g$ and 2^f .
- 8. For any constructible function f, and any language $L \in \mathsf{NTIME}(f(n))$, there is a nondeterministic machine M that accepts L and all of whose computations terminate in time O(f(n)) for all inputs of length n. Give a detailed argument for this statement, describing how M might be obtained from a machine accepting L in time f(n).
- 9. In the lecture, a proof of the Time Hierarchy Theorem was sketched. Give a similar argument for the following Space Hierarchy Theorem:

Space Hierarchy. For every constructible function f, there is a language in $\mathsf{SPACE}(f(n) \cdot \log f(n))$ that is not in $\mathsf{SPACE}(f(n))$.

- 10. Show that, if $\mathsf{SPACE}((\log n)^2) \subseteq \mathsf{P}$, then $\mathsf{L} \neq \mathsf{P}$. (Hint: use the Space Hierarchy Theorem from Exercise 9 above.)
- 11. POLYLOGSPACE is the complexity class

$$\bigcup_k \mathsf{SPACE}((\log n)^k).$$

- (a) Show that, for any k, if $A \in \mathsf{SPACE}((\log n)^k)$ and $B \leq_L A$, then $B \in \mathsf{SPACE}((\log n)^k)$.
- (b) Show that there are no POLYLOGSPACE-complete problems with respect to \leq_L . (Hint: use (a) and the space hierarchy theorem).
- (c) Which of the following might be true: $P \subseteq POLYLOGSPACE$, $P \supseteq POLYLOGSPACE$, P = POLYLOGSPACE?
- (d) What is the relationship between the class POLYLOGSPACE and the class Quasi-P (see Exercise Sheet 1, Question 7)?