

Foundations of Computer Science

Lecture #11: Procedural Programming

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Intro

```
let rec addLen n = function
| [] -> n
| x :: xs -> addLen (n+1) xs
```

Example:

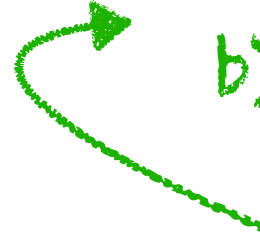
```
addLen 0 [1,2,3]
```

Calling `addLen` with same arguments will *always* produce the same result. We can infer result through function expansion and reduction of expressions. This allows us to:

- Prove algorithm correctness
- Understand and predict algorithm outcome

Procedural Programming

- a) update variable / array
- b) sending / receiving data



Procedural programs can change the machine state.

They can interact with its *environment*.

They use control structures like *branching*, *iteration* and *procedures*.

They use data abstractions of the computer's memory:

- *references* to memory cells
- *arrays*: blocks of memory cells
- *linked structures*, especially *linked lists*

concept: memory cells that are mutable

What are References?

In functional programming:

The store is an ***invisible*** device inside the computer

In procedural / imperative programming:

The store is ***visible***

What are References?

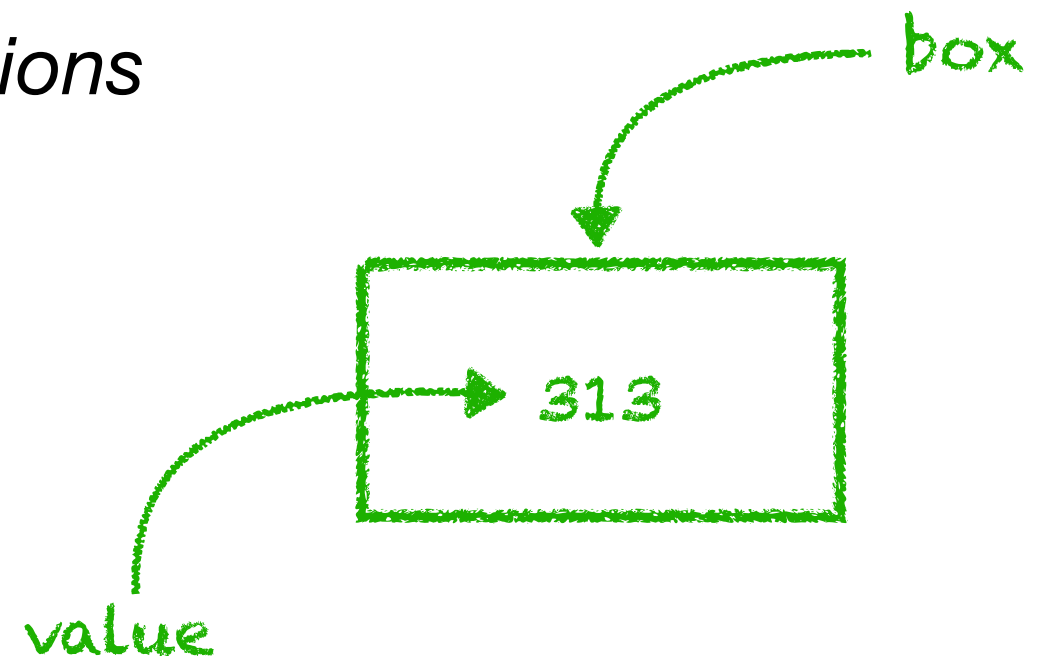
In functional programming:

The store is an ***invisible*** device inside the computer

In procedural / imperative programming:

The store is ***visible***

- References are *storage locations*
- They can be:
 - (a) created
 - (b) inspected
 - (c) updated



The box has an address

ML Primitives for References

$\tau \text{ ref}$ *type* of references to type τ

$\text{ref } E$ *create* a reference
initial contents = the value of E

$! P$ return the *current contents* of reference P 'dereferencing'

$P := E$ *update* the contents of P to the value of E

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P for 'pointer'

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$P := E$

update the contents of P to the value of E

pointer to a 'box'

contents of that 'box'

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Three new ML functions / operators:

ref : 'a \rightarrow 'a ref

! : 'a ref \rightarrow 'a

:= : 'a ref \rightarrow 'a \rightarrow unit

(a) create box

(b) inspect box content

(c) update box content

Trying Out References

```
# let p = ref 5 (* create a reference *)
```

```
val p : int ref = {contents = 5}
```

```
# p := !p + 1    (* p now holds value 6 *)
```

```
- : unit = ()
```

```
# let ps = [ ref 77; p ]
```

```
val ps : int ref list = [{contents = 77}; {contents = 6}]
```

```
# List.hd ps := 3
```

```
- : unit = ()
```

```
# ps
```

```
- : int ref list = [{contents = 3}; {contents = 6}]
```

Trying Out References

```
# let p = ref 5 (* create a reference *)  
val p : int ref = {contents = 5}  
  
# let z = p  
val z : int ref = {contents = 5}  
  
# p := !p + 1    (* p now holds value 6 *)  
- : unit = ()  
  
# p  
- : int ref = {contents = 6}  
  
# z  
- : int ref = {contents = 6}
```

Aliasing: two values refer to the same mutable cell

Commands: Expressions with Effects

- Basic commands update references, write to files, etc.
- $C_1; \dots; C_n$ causes a series of expressions to be evaluated and returns the value of C_n .
- A typical command returns the empty tuple: $()$
- if B then C_1 else C_2 behaves like the traditional control structure if C_1 and C_2 have effects.
- Other ML constructs behave naturally with commands, including match expressions and recursive functions.

Commands: Expressions with Effects

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- Other ML constructs behave naturally with commands, including match expressions and recursive functions.

Example:

$> 1 + (\text{print_newline } "abc"; 3; 101);$

abc

$- : int = 102$

Example: length without Mutability

```
let rec addLen n = function
| [] -> n
| x :: xs -> addLen (n+1) xs
```

```
addLen 0 [1,2,3]
addLen 1 [2,3]
addLen 2 [3]
addLen 3 []
==> returns 3
```

Iteration: the `while` Command

```
# let tlopt = function
| [] -> None
| _::xs -> Some xs
val tlopt : 'a list -> 'a list option = <fun>

# let length xs =
  let lp = ref xs in (* list of uncounted elements *)
  let np = ref 0 in (* accumulated count *)
  let fin = ref false in
  while not !fin do
    match tlopt !lp with
    | None -> fin := true
    | Some xs ->
      lp := xs;
      np := 1 + !np
  done;
  !np (* the final count is returned *)
val length : 'a list -> int = <fun>
```

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  done;
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val length : 'a list -> int = <fun>
```

Example: length with Mutability

evaluation steps:

```
length([1;2;3])  
==> lp = ref [1,2,3]
```

```
tlopt [1;2;3] != None ==> true  
lp := [2,3]; np := 1+0;
```

```
tlopt [2;3] != None ==> true  
lp := [3]; np := 1+1
```

```
tlopt [3] != None ==> true  
lp := []; np := 1+2
```

```
tlopt [] != None ==> false  
fin := true  
==> return !np  
==> returns 3
```

```
let tlopt = function
  | [] -> None
  | _ :: xs -> Some xs
```

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let length xs =
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  done;
  !np
val length : 'a list -> int = <fun>
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```
let rec addLen n =
  function
  | [] -> n
  | x :: xs ->
    addLen (n+1) xs
```

Private, Persistent References

```
# exception TooMuch of int
exception TooMuch of int
# let makeAccount initBalance =
  let balance = ref initBalance in
  let withdraw amt =
    if amt > !balance then
      raise (TooMuch (amt - !balance))
    else begin
      balance := !balance - amt;
      !balance
    end
  in
  withdraw
val makeAccount : int -> int -> int = <fun>
```

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returns a function that
returns contents of
'balance', not the cell itself

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  in
  withdraw
val makeAccount : int -> int -> int = <fun>
```

balance never escapes the
definition of makeAccount

returns a function that
returns contents of
'balance', not the cell itself

Private, Persistent References

```
let my_account = makeAccount 30;
```

```
my_account : int -> int = <fun>
```

```
let my_new_balance = my_account 10;
```

```
my_new_balance : int = 20
```

```
let my_new_balance = my_account ~10;
```

```
my_new_balance : int = 30
```


Two Bank Accounts

```
# let student = makeAccount 500
val student : int -> int = <fun>

# let director = makeAccount 4000000
val director : int -> int = <fun>

# student 5          (* coach fare *)
- : int = 495

# director 150000    (* Tesla *)
- : int = 3850000

# student 500        (* oh oh *)
Exception: TooMuch 5.
```

ML Primitives for Arrays

```
# [| "a"; "b"; "c" |]
  (* allocate a fresh string array *)
- : string array = [| "a"; "b"; "c" |]

# Array.make 3 'a'
  (* array of size 3 with cell containing 'a' *)
- : char array = [| 'a'; 'a'; 'a' |]

# let aa = Array.init 5 (fun i -> i * 10)
  (* array of size 5 initialised to (fun i) *)
val aa : int array = [| 0; 10; 20; 30; 40 |]

# Array.get aa 3
  (* retrieve the 4th cell in the array *)
- : int = 30

# Array.set aa 3 42
  (* set the 4th cell's value to 42 *)
- : unit = ()
```

Array Examples

Array.make

- : *int -> 'a -> 'a array = <fun>*

Array.init

- : *int -> (int -> 'a) -> 'a array = <fun>*

Array.get

- : *'a array -> int -> 'a = <fun>*

Array.set

- : *'a array -> int -> 'a -> unit = <fun>*

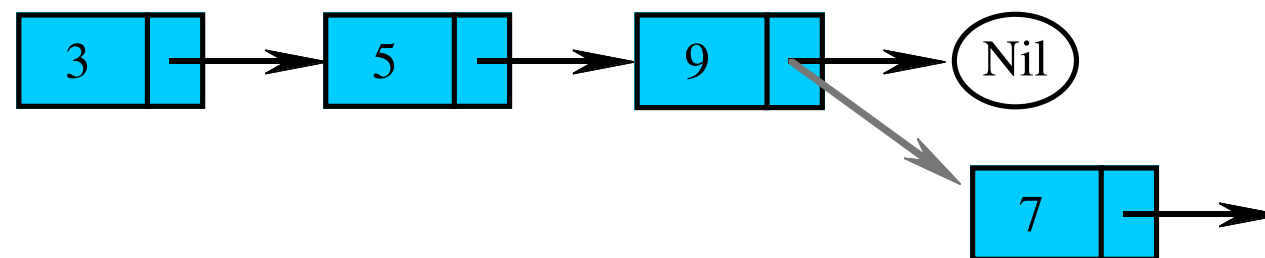
References: ML Versus Conventional Languages

- We must write $!p$ to get the *contents* of p
- We write just p for the address of p
- We can store *private* reference cells in functions; simulating object oriented programming
- OCaml's assignment syntax is $V := E$ instead of $V = E$
- OCaml has similar control structures: `while/done`, `for/done` and `match/with`
- OCaml has short syntax for updating arrays $x.(1)$ and the access is safe against buffer overflows

What More Is There to ML?

With references, we can now make mutable linked lists

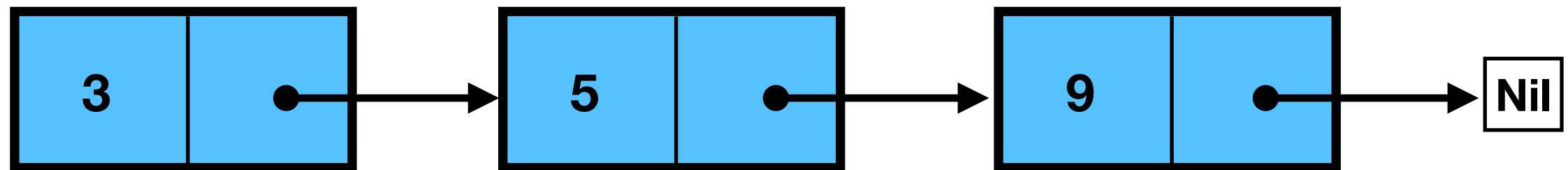
```
# type 'a mlist =  
  | Nil  
  | Cons of 'a * 'a mlist ref  
type 'a mlist = Nil | Cons of 'a * 'a mlist ref
```



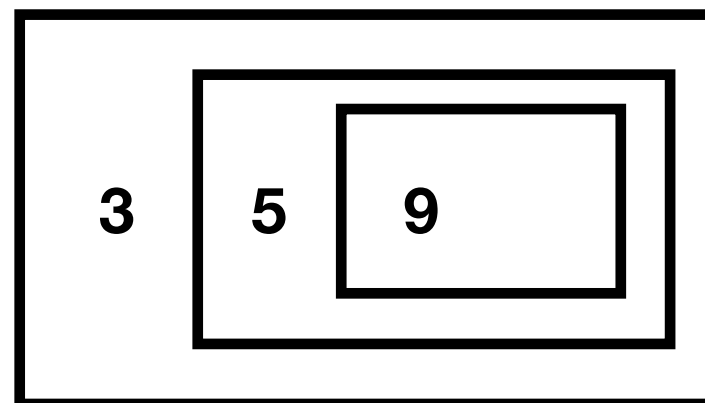
References to References

Two ways to visualize references to references:

(1) Using pointers:



(2) Using nested boxes:



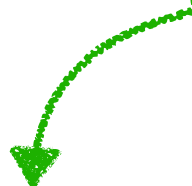
Linked (Mutable) Lists

```
# type 'a mlist =  
  | Nil  
  | Cons of 'a * 'a mlist ref  
type 'a mlist = Nil / Cons of 'a * 'a mlist ref
```

→ The tail can be redirected!

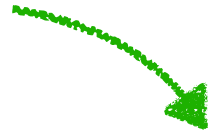
```
# let rec mlistOf = function  
  | [] -> Nil  
  | x :: l -> Cons (x, ref (mlistOf l))  
mlist : 'a list -> 'a mlist = <fun>
```

creates a new pointer to rest of mlist



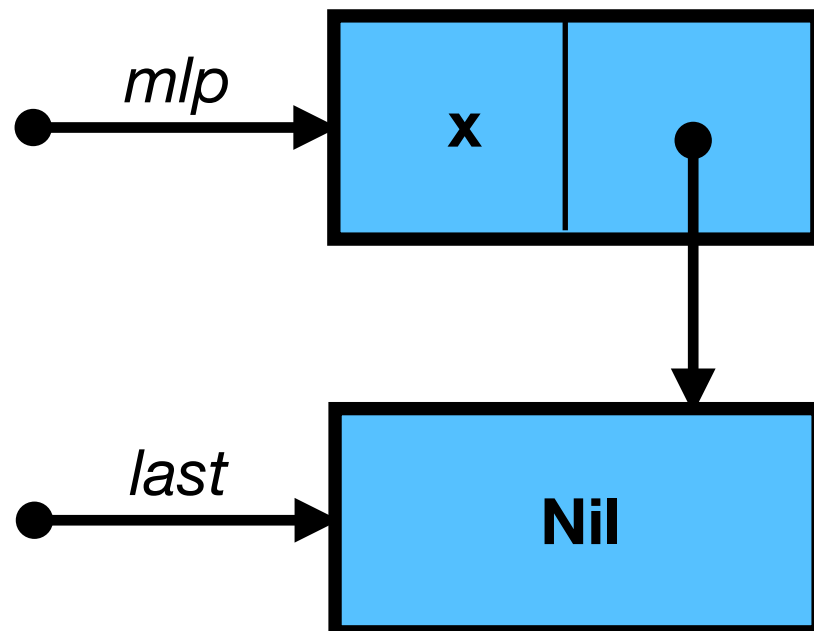
Extending a List to the Rear

pointing to a 'box'



```
# let extend mlp x =  
  let last = ref Nil in  
  mlp := Cons (x, last);  
  last
```

```
> val extend = fn : 'a mlist ref * 'a -> 'a mlist ref
```



Example of Extending a List

```
# let mlp = ref (Nil: string mlist);;  
val mlp : string mlist ref = {contents = Nil}
```

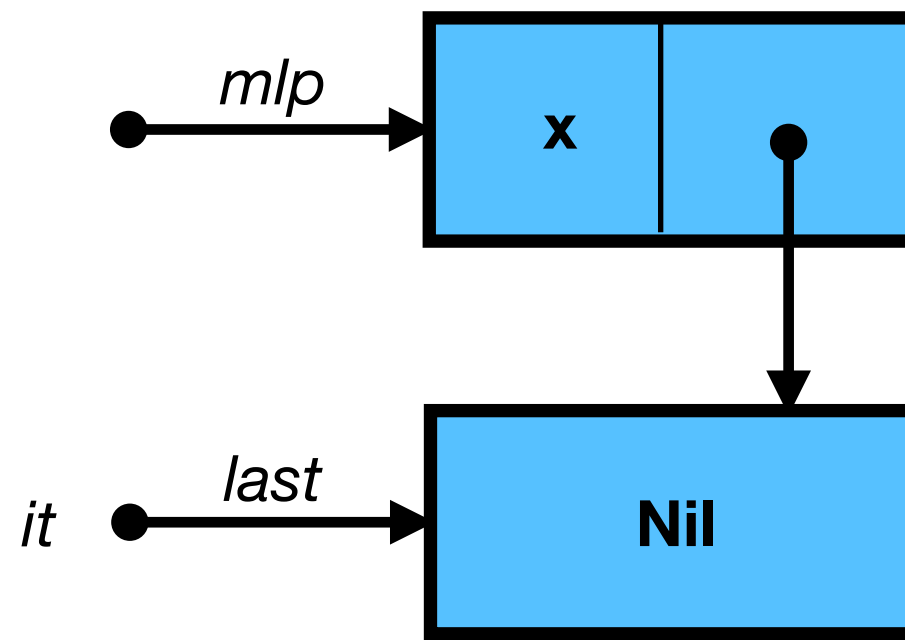
```
# extend mlp "a";;  
- : string mlist ref = {contents = Nil}
```

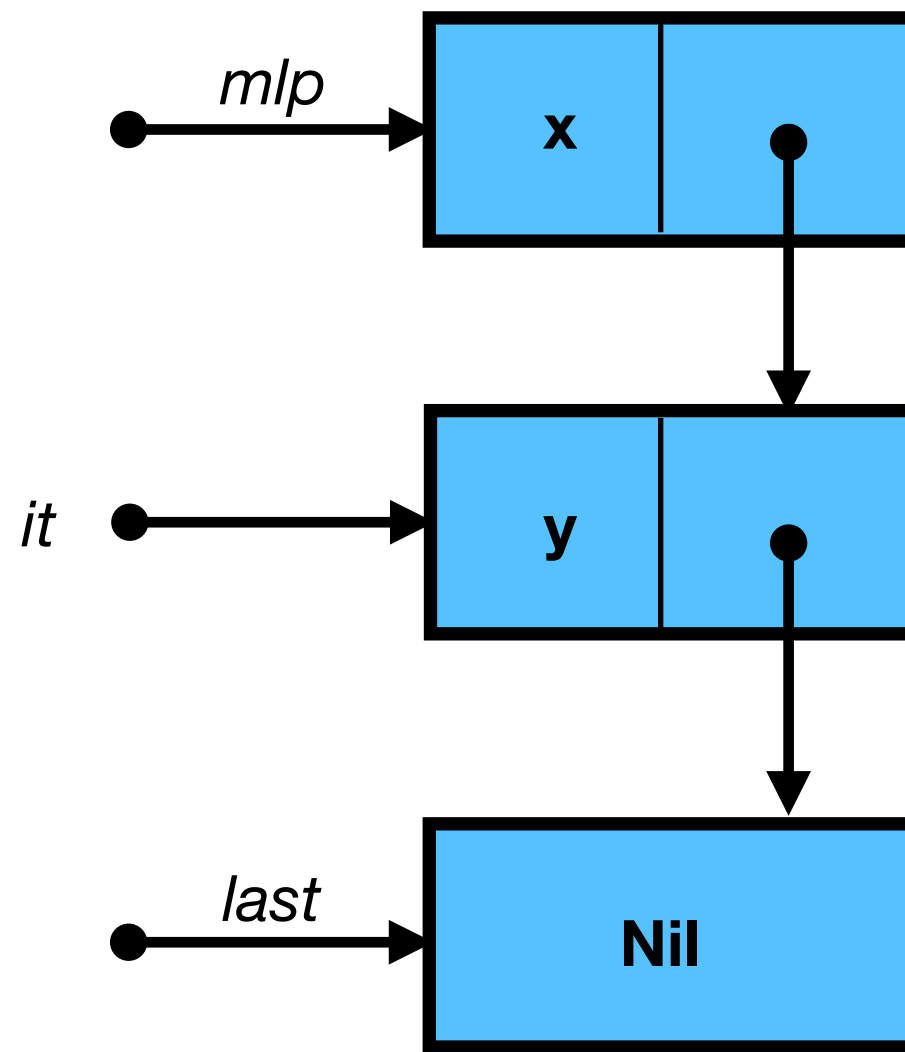
```
# let mlp = ref (Nil : string mlist);;  
val mlp : string mlist ref = {contents = Nil}
```

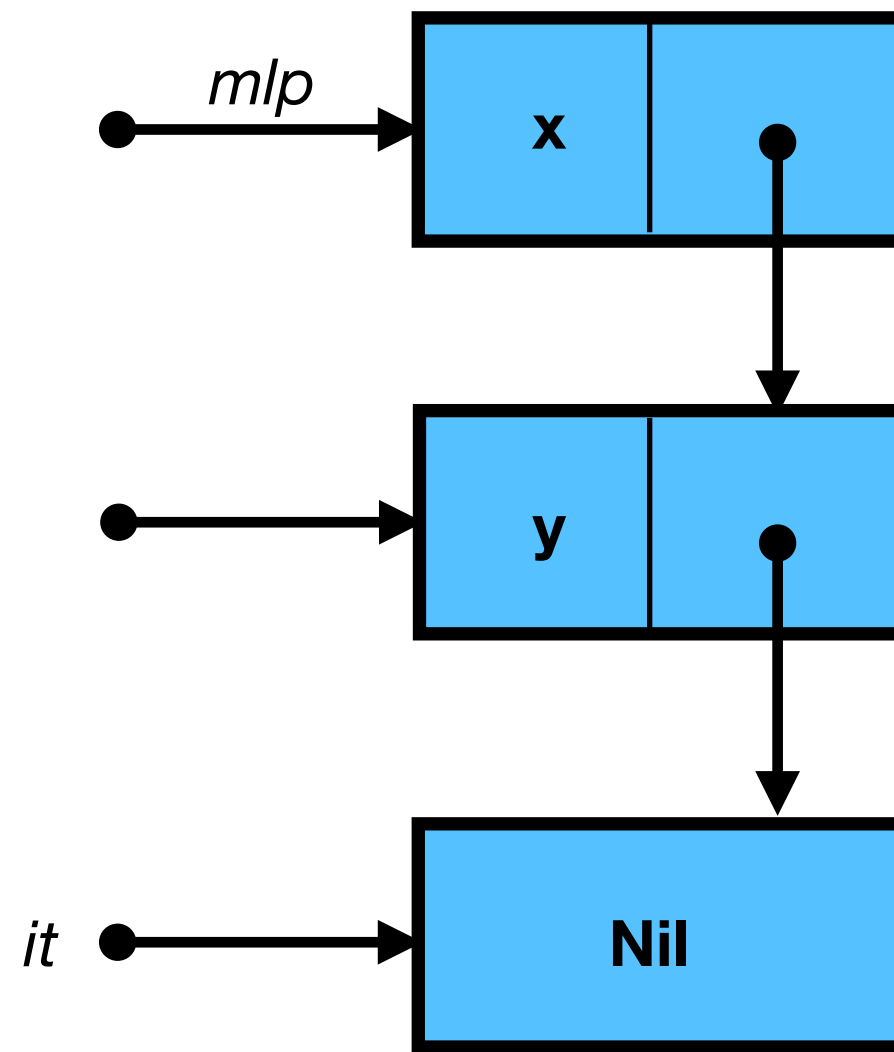
```
# let it = extend mlp "a" ;;  
val it : string mlist ref = {contents = Nil}
```

```
# extend it "b" ;;  
- : string mlist ref = {contents = Nil}
```

```
# mlp ;;  
- : string mlist ref =  
{contents = Cons ("a",  
  {contents = Cons ("b", {contents = Nil})})}
```







$ref (Cons (x, ref (Cons (y, ref Nil))))$

Destructive Concatenation

pointing to a 'box'

contents of a 'box'

```
# let rec joining mlp m12 =  
  match !mlp with  
  | Nil -> mlp := m12  
  | Cons (_, mlp1) -> joining mlp1 m12  
val joining : 'a mlist ref * 'a mlist -> unit = <fun>
```

```
# let join m11 m12 =  
  let mlp = ref m11 in  
  joining mlp m12;  
  !mlp  
val join : 'a mlist -> 'a mlist -> 'a mlist = <fun>
```

Side-Effects

```
# let m11 = mlistOf ["a"];;  
val m11 : string mlist = Cons ("a", {contents = Nil})  
# let m12 = mlistOf ["b";"c"];;  
val m12 : string mlist =  
    Cons ("b", {contents = Cons ("c", {contents = Nil})})  
# join m11 m12 ;;
```

What does this return?

```
- : string mlist =  
Cons ("a",  
    {contents = Cons ("b",  
        {contents = Cons ("c", {contents = Nil})})})
```