

Relations

Definition 99 A (binary) relation R from a set A to a set B

$$R : A \dashrightarrow B \quad \text{or} \quad R \in \text{Rel}(A, B) \quad ,$$

is

$$R \subseteq A \times B \quad \text{or} \quad R \in \mathcal{P}(A \times B) \quad .$$

Notation 100 One typically writes $a R b$ for $(a, b) \in R$.

Informal examples:

- ▶ Computation.
- ▶ Typing.
- ▶ Program equivalence.
- ▶ Networks.
- ▶ Databases.

Examples:

- ▶ Empty relation.

$$\emptyset : A \dashrightarrow B$$

$$(a \emptyset b \iff \mathbf{false})$$

- ▶ Full relation.

$$(A \times B) : A \dashrightarrow B$$

$$(a (A \times B) b \iff \mathbf{true})$$

- ▶ Identity (or equality) relation.

$$\text{id}_A = \{ (a, a) \mid a \in A \} : A \dashrightarrow A$$

$$(a \text{id}_A a' \iff a = a')$$

- ▶ Integer square root.

$$R_2 = \{ (m, n) \mid m = n^2 \} : \mathbb{N} \dashrightarrow \mathbb{Z}$$

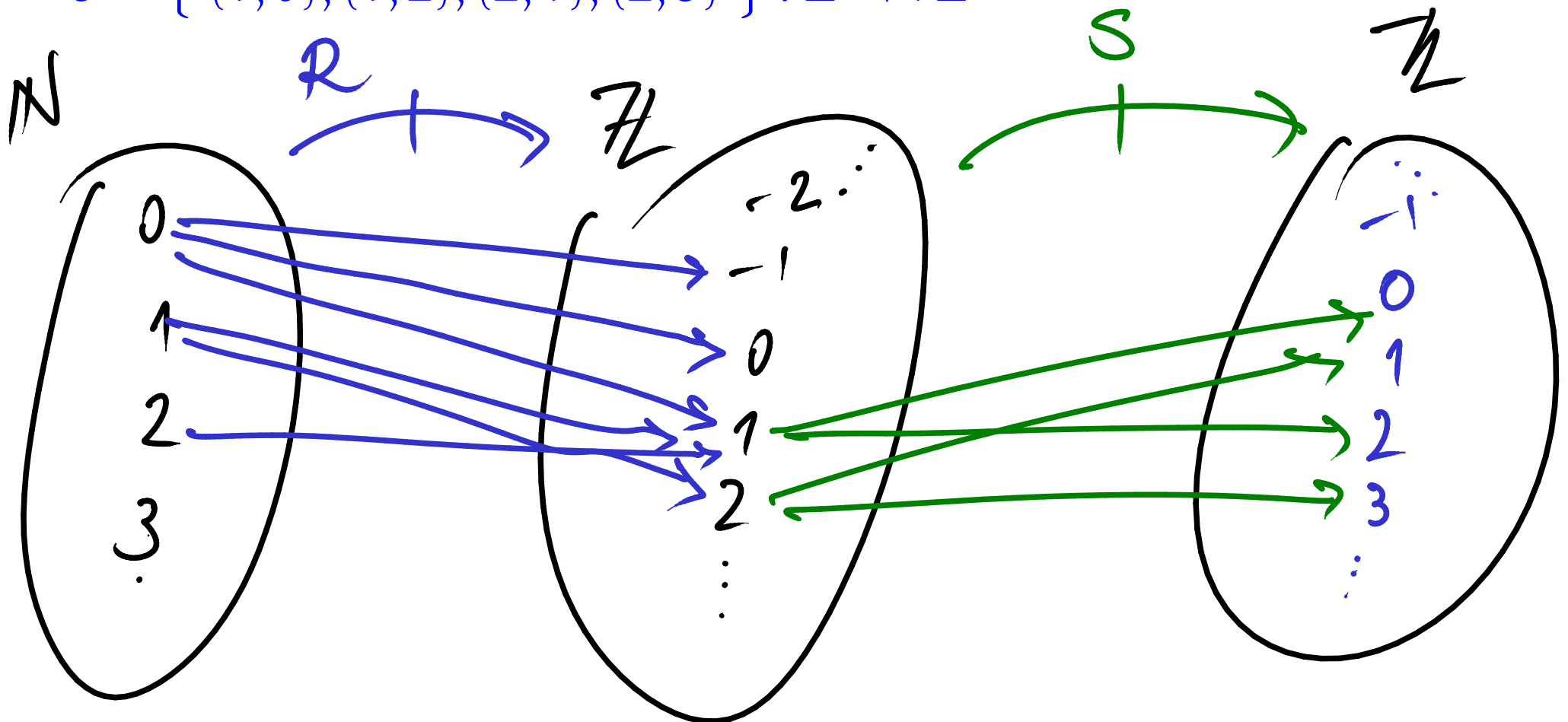
$$(m R_2 n \iff m = n^2)$$

Internal diagrams

Example:

$$R = \{ (0, 0), (0, -1), (0, 1), (1, 2), (1, 1), (2, 1) \} : \mathbb{N} \dashrightarrow \mathbb{Z}$$

$$S = \{ (1, 0), (1, 2), (2, 1), (2, 3) \} : \mathbb{Z} \dashrightarrow \mathbb{Z}$$



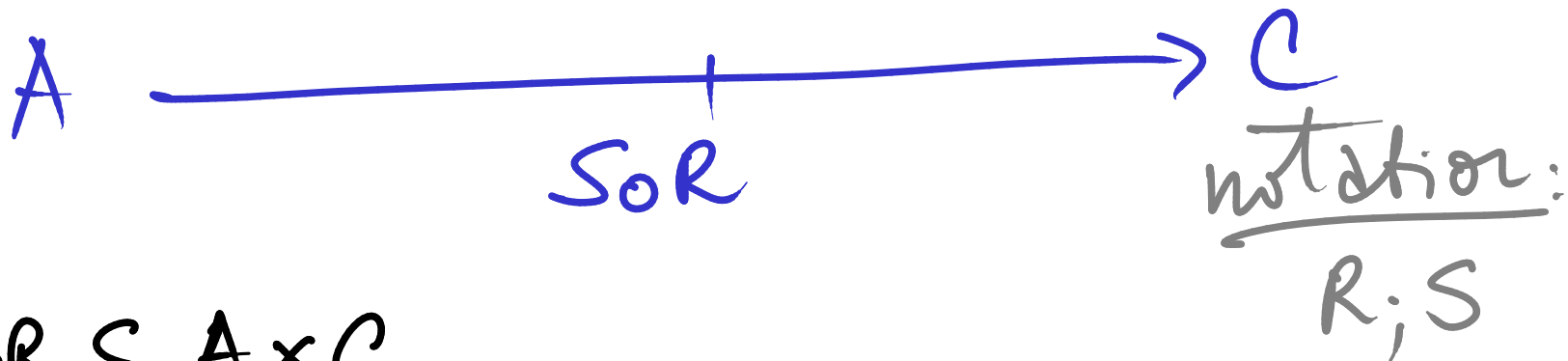
Relational extensionality

$$R = S : A \rightarrow B$$

iff

$$\forall a \in A. \forall b \in B. a R b \iff a S b$$

Relational composition



$$SoR \subseteq A \times C$$

$$\forall a \in A, c \in C.$$

$$a(SoR)c \stackrel{\text{def.}}{\Leftrightarrow} \exists b \in B. aRb \wedge bSc$$

Example

$$\underline{Sq}: \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}$$

$$\underline{Neg}: \mathbb{R} \rightarrow \mathbb{R}$$

$$a \underline{Sq} b \stackrel{\text{def}}{\iff} a = b^2$$

$$x \underline{Neg} y \stackrel{\text{def}}{\iff} y = -x$$

Claim:

$$\underline{Neg} \circ \underline{Sq} = \underline{Sq}: \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}$$

RTP: $\forall s \in \mathbb{R}_{\geq 0}, t \in \mathbb{R}.$

$$\stackrel{\text{def}}{\iff} s (\underline{Neg} \circ \underline{Sq}) t \stackrel{?}{\iff} s \underline{Sq} t \stackrel{\text{by def}}{\iff} s = t^2$$

\swarrow exercise

$$\exists r \in \mathbb{R}. s \underline{Sq} r \wedge r \underline{Neg} t \iff \exists r \in \mathbb{R}. s = r^2 \wedge t = -r$$

Theorem 102 *Relational composition is associative and has the identity relation as neutral element.*

► *Associativity.*

For all $R : A \rightarrow B$, $S : B \rightarrow C$, and $T : C \rightarrow D$,

$$(T \circ S) \circ R = T \circ (S \circ R)$$

NB. Unambiguously
 $T \circ S \circ R$

► *Neutral element.*

For all $R : A \rightarrow B$,

$$R \circ \text{id}_A = R = \text{id}_B \circ R .$$

$$(T \circ S) \circ R = T \circ (S \circ R): A \rightarrow D$$

$\forall a \in A, d \in D.$

$$a((T \circ S) \circ R) d \stackrel{?}{\Leftrightarrow} a(T \circ (S \circ R)) d$$

$$\begin{aligned} &\Downarrow \\ &\exists c \in C. a(S \circ R) c \wedge c T d \end{aligned}$$

$$\begin{aligned} &\Uparrow \\ &\exists b \in B. a R b \wedge b(T \circ S) d \end{aligned}$$

$$\begin{aligned} &\Downarrow \\ &\exists c \in C. \exists b \in B. \end{aligned}$$

$$\begin{aligned} &\Downarrow \\ &\exists b \in B. a R b \wedge \exists c \in C. b S c \wedge c T d \end{aligned}$$

$$a R b \wedge b S c \wedge c T d$$

$$\begin{aligned} &\Downarrow \\ &\exists b \in B. \exists c \in C. a R b \wedge b S c \wedge c T d. \end{aligned}$$



Relations and matrices

Definition 103

1. For positive integers m and n , an $(m \times n)$ -matrix M over a ~~COMM.~~ semiring $(S, 0, \oplus, 1, \odot)$ is given by entries $M_{i,j} \in S$ for all $0 \leq i < m$ and $0 \leq j < n$.

$$M = \begin{bmatrix} 0 & 1 & \dots & j & \dots & n-1 \\ \vdots & & & M_{i,j} & & \\ \vdots & & & & & \\ \vdots & & & & & \\ m-1 & & & & & \end{bmatrix}$$

Theorem 104 Matrix multiplication is associative and has the identity matrix as neutral element.

M ($m \times n$)-matrix
 N ($n \times l$)-matrix
 $N \cdot M$ ($m \times l$)-matrix

I ($m \times m$)-matrix
 $I_{ij} = \begin{cases} 1 & i=j \\ 0 & \text{otherwise} \end{cases}$

$$(N \cdot M)_{i,j} = \sum_R M_{i,R} \odot N_{R,j}$$

iterated \oplus of S

Recall: $[k] = \{0, 1, 2, \dots, k-1\}$ $k \in \mathbb{N}$

Relations from $[m]$ to $[n]$ and $(m \times n)$ -matrices over Booleans provide two alternative views of the same structure.

This carries over to identities and to composition/multiplication.

Given

$$R: [m] \rightarrow [n]$$

Define

$\underline{\text{mat}}(R)$ $(m \times n)$ -matrix

$$\left(\underline{\text{mat}}(R) \right)_{i,j} = \begin{cases} \text{true} & i R j \\ \text{false} & \text{else} \end{cases}$$

$i \in [m], j \in [n]$

$$\text{Bool} = \{ \text{true}, \text{false} \}$$

$$\oplus = \vee$$

$$\odot = \wedge$$

$$R: [3] \rightarrow [2]$$

$$\underline{\text{mat}}(R) = \begin{array}{c} 0 \\ 1 \\ 2 \end{array} \begin{array}{cc} 0 & 1 \\ \left[\begin{array}{cc} \text{true} & \text{false} \\ \text{false} & \text{true} \\ \text{false} & \text{false} \end{array} \right] \end{array}$$

$$R = \{ (0,0), (1,1) \}$$

Given $(m \times n)$ -matrix M , define $\underline{\text{rel}}(M): \widehat{[m]} \rightarrow \widehat{[n]}$

$\forall i \in \widehat{[m]}, j \in \widehat{[n]},$

$$i (\underline{\text{rel}}(M)) j \stackrel{\text{def}}{\iff} (M_{i,j} = \text{true})$$

Proposition:

$$\underline{\text{rel}}(\underline{\text{mat}}(R)) = R$$

$$\underline{\text{mat}}(\underline{\text{rel}}(M)) = M$$

Proposition: $[m] \xrightarrow{R} [n]$, $[n] \xrightarrow{S} [l]$

$\text{mat}(R)$
 $(m \times n)$ -matrix

$\text{mat}(S)$
 $(n \times l)$ -matrix

$$\underline{\text{mat}}(S \circ R) = \underline{\text{mat}}(S) \cdot \underline{\text{mat}}(R) \quad (*)$$

Corollary:

$$S \circ R = \underline{\text{rel}} \left(\underline{\text{mat}}(S) \cdot \underline{\text{mat}}(R) \right)$$

$$\left(\underline{\text{mat}}(S \circ R) \right)_{i,j} = \underline{\text{true}}$$

$$\Leftrightarrow i(S \circ R)j$$

$$\Leftrightarrow \exists k. iRk \wedge kSj$$

$$\left(\underline{\text{mat}}(S) \cdot \underline{\text{mat}}(R) \right)_{i,j}$$

$$= \bigvee_k. \underline{\text{mat}}(R)_{i,k} \wedge \underline{\text{mat}}(S)_{k,j}$$

$$= \bigvee_k (i,k) \in R \wedge (k,j) \in S$$

$$\underline{\text{mat}}(\text{id}) = I$$

Def M, N ($m \times n$) matrices.

$$(M+N)_{i,j} = M_{i,j} \oplus N_{i,j}$$

Prop: $R, S: [m] \rightarrow [n]$

$$\begin{cases} \underline{\text{mat}}(R \cup S) = \underline{\text{mat}}(R) + \underline{\text{mat}}(S) \\ \underline{\text{mat}}(\emptyset) = 0 \end{cases}$$