

Lemma 58 For all positive integers m and n ,

$$\text{CD}(m, n) = \begin{cases} D(n) & , \text{ if } n \mid m \\ \text{CD}(n, \text{rem}(m, n)) & , \text{ otherwise} \end{cases}$$

Since a positive integer n is the greatest divisor in $D(n)$, the lemma suggests a recursive procedure:

$$\text{gcd}(m, n) = \begin{cases} n & , \text{ if } n \mid m \\ \text{gcd}(n, \text{rem}(m, n)) & , \text{ otherwise} \end{cases}$$

for computing the *greatest common divisor*, of two positive integers m and n . This is

Euclid's Algorithm

NB: If $\text{gcd}(m, n)$ terminates say with output R ,
Then $\underline{CD}(m, n) = \underline{D}(R)$.

gcd

```
fun gcd( m , n )  
  = let  
    val ( q , r ) = divalg( m , n )  
  in  
    if r = 0 then n  
    else gcd( n , r )  
  end
```

Example 59 ($\gcd(13, 34) = 1$)

$$\begin{aligned}\gcd(13, 34) &= \gcd(34, 13) \\ &= \gcd(13, 8) \\ &= \gcd(8, 5) \\ &= \gcd(5, 3) \\ &= \gcd(3, 2) \\ &= \gcd(2, 1) \\ &= 1\end{aligned}$$

$$\underline{CD}(m, n) = \underline{D}(k)$$

$$\underline{CD}(m, n) = \{d \in \mathbb{N} : d|m \wedge d|n\}$$

$$\underline{D}(k) = \{d \in \mathbb{N} : d|k\}$$

$$\Leftrightarrow \left[\begin{array}{l} \forall d \in \mathbb{N}. \\ (d|m \wedge d|n) \Leftrightarrow d|k \end{array} \right]$$

$$\Leftrightarrow \left[\begin{array}{l} (1) k|m \wedge k|n \\ \wedge (2) \forall d \in \mathbb{N}. d|m \wedge d|n \Rightarrow d|k \end{array} \right]$$

Proper tries (1) and (2) unique by character k .

Suppose $\left[\begin{array}{l} (1)_1, k_1 | m \sim k_1 | n \\ (2)_1, \forall d. d | m \sim d | n \Rightarrow d | k_1 \end{array} \right]$

and $\left[\begin{array}{l} (1)_2, k_2 | m \sim k_2 | n \\ (2)_2, \forall d. d | m \sim d | n \Rightarrow d | k_2 \end{array} \right]$

Then, we claim, $k_1 = k_2$.

Theorem 60 *Euclid's Algorithm \gcd terminates on all pairs of positive integers and, for such m and n , $\gcd(m, n)$ is the greatest common divisor of m and n in the sense that the following two properties hold:*

- (i) *both $\gcd(m, n) \mid m$ and $\gcd(m, n) \mid n$, and*
- (ii) *for all positive integers d such that $d \mid m$ and $d \mid n$ it necessarily follows that $d \mid \gcd(m, n)$.*

PROOF:

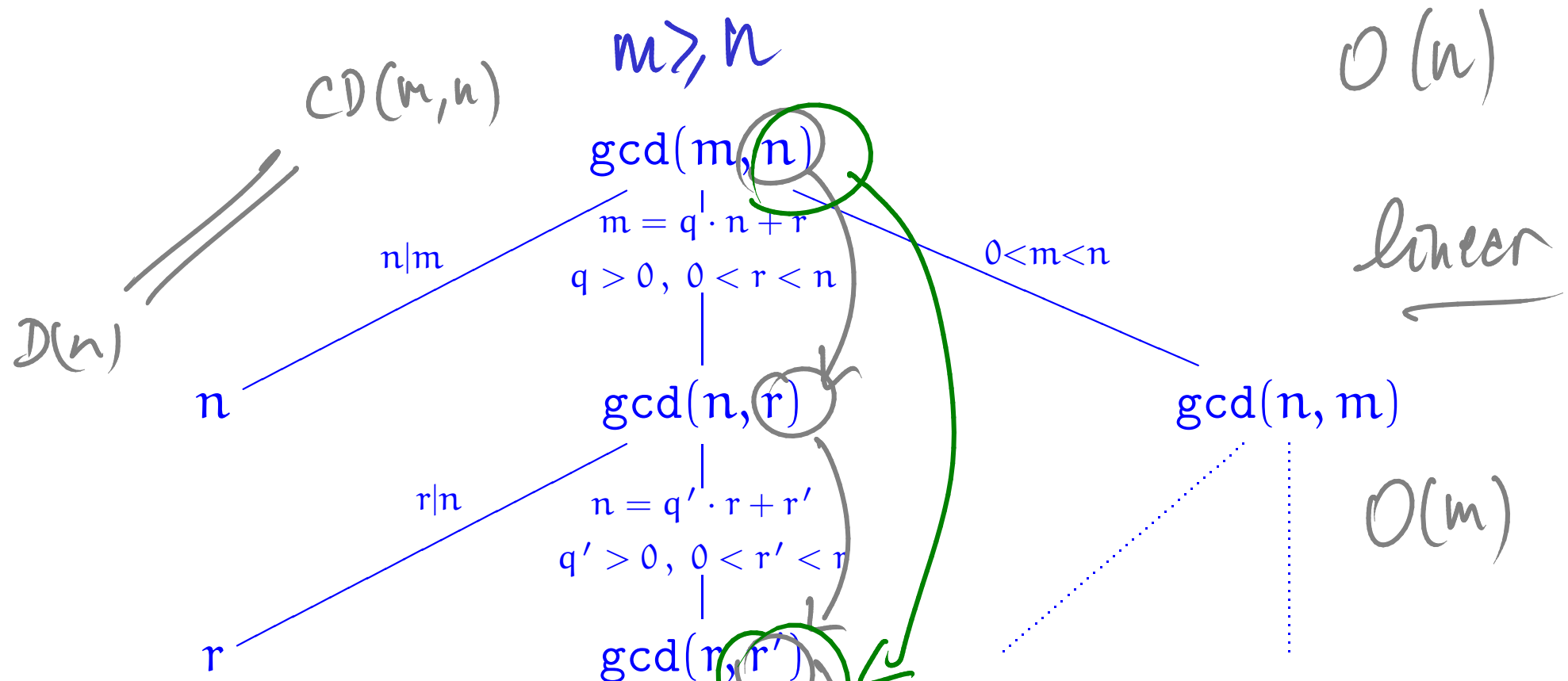
PROOF PRINCIPLE

To show that some k is the gcd of m and n

go on to show that

$$(1) \quad k \mid m \text{ and } k \mid n$$

$$\text{and } (2) \quad \forall d. \ d \mid m \wedge d \mid n \Rightarrow d \mid k.$$



$n = q' \cdot r + r' \Rightarrow r + r' > 2r'$
 $\Rightarrow r' < \frac{n}{2}$

decreases while remaining positive.

running time is $O(\log(\max(m, n)))$.

Fractions in lowest terms

```
fun lowterms( m , n )  
  = let  
    val gcdval = gcd( m , n )  
  in  
    ( m div gcdval , n div gcdval )  
  end
```


Some fundamental properties of gcds

Lemma 62 For all positive integers l , m , and n ,

1. **(Commutativity)** $\gcd(m, n) = \gcd(n, m)$,

2. **(Associativity)** $\gcd(l, \gcd(m, n)) = \gcd(\gcd(l, m), n)$,

3. **(Linearity)^a** $\gcd(l \cdot m, l \cdot n) = \underline{l \cdot \gcd(m, n)}$.

PROOF:

We show (1) $l \cdot \gcd(m, n) \mid lm$ and $l \cdot \gcd(m, n) \mid ln$

and (2) $\forall d. d \mid lm$ and $d \mid ln$
then $d \mid l \cdot \gcd(m, n)$.

^aAka (Distributivity).

$$(1) \text{ RTP : } l \cdot \underline{\text{gcd}}(m, n) \mid l \cdot m$$

Since $\underline{\text{gcd}}(m, n) \mid m$ then $l \cdot \underline{\text{gcd}}(m, n) \mid l \cdot m$

$$[\text{Lemma } a \mid b \Rightarrow a \cdot c \mid b \cdot c]$$

Analogously $l \cdot \underline{\text{gcd}}(m, n) \mid l \cdot n$.

(2) RTP : $\forall d$ if $d \mid lm$ and $d \mid ln$ then $d \mid l \cdot \underline{\text{gcd}}(m, n)$

Let d be an arbitrary pos. int. such that

$$\textcircled{1} \quad \underline{d \mid l \cdot m} \text{ and } \textcircled{2} \quad \underline{d \mid l \cdot n} \quad (*) \quad [(a \mid b \wedge b \mid c) \Rightarrow a \mid c]$$

$$\text{RTP : } d \mid l \cdot \underline{\text{gcd}}(m, n)$$

From $\textcircled{1}$ and $\textcircled{2}$, we have $d \mid \underline{\text{gcd}}(l \cdot m, l \cdot n)$

If $\underline{\text{gcd}}(l \cdot m, l \cdot n) \mid l \cdot \underline{\text{gcd}}(m, n)$ by $(*)$ we will be done.

We want to show

$$\left. \begin{array}{l} \text{gcd}(l \cdot m, l \cdot n) \\ \text{gcd}(m, n) \end{array} \right\} \mid l \cdot \text{gcd}(m, n)$$

Exercise.

Euclid's Theorem

Theorem 63 For positive integers k , m , and n , if $k \mid (m \cdot n)$ and $\gcd(k, m) = 1$ then $k \mid n$.

PROOF: Assume $k \mid (m \cdot n)$ and $\gcd(k, m) = 1$

\Downarrow \Downarrow

$l \cdot k = m \cdot n$ $n \cdot \gcd(k, m) = n$

for some l \parallel

$\underbrace{\hspace{15em}}$
 $\gcd(nk, nm)$

$$k \cdot \gcd(n, l) = \gcd(nk, l \cdot k) = n \implies k \mid n. \quad \square$$

Corollary 64 (Euclid's Theorem) For positive integers m and n , and prime p , if $p \mid (m \cdot n)$ then $p \mid m$ or $p \mid n$.

Now, the second part of Fermat's Little Theorem follows as a corollary of the first part and Euclid's Theorem.

PROOF: Assume $p \mid (m \cdot n)$

case ① if $p \mid m$ then we're done.

case ② if $p \nmid m$ then $\gcd(p, m) = 1$
and so $p \mid n$.



FLT

$$i^p \equiv i \pmod{p}$$

suppose $i \not\equiv 0 \pmod{p}$

then $p \mid (i^p - i) = (i^{p-1} - 1)i$

and by Euclid's Thm.

$$p \mid i^{p-1} - 1$$

That is, $i^{p-1} \equiv 1 \pmod{p}$.

NB:

$$p \mid \binom{p}{m}$$

$$0 < m < p$$

(prime p)

}
Exercise.