

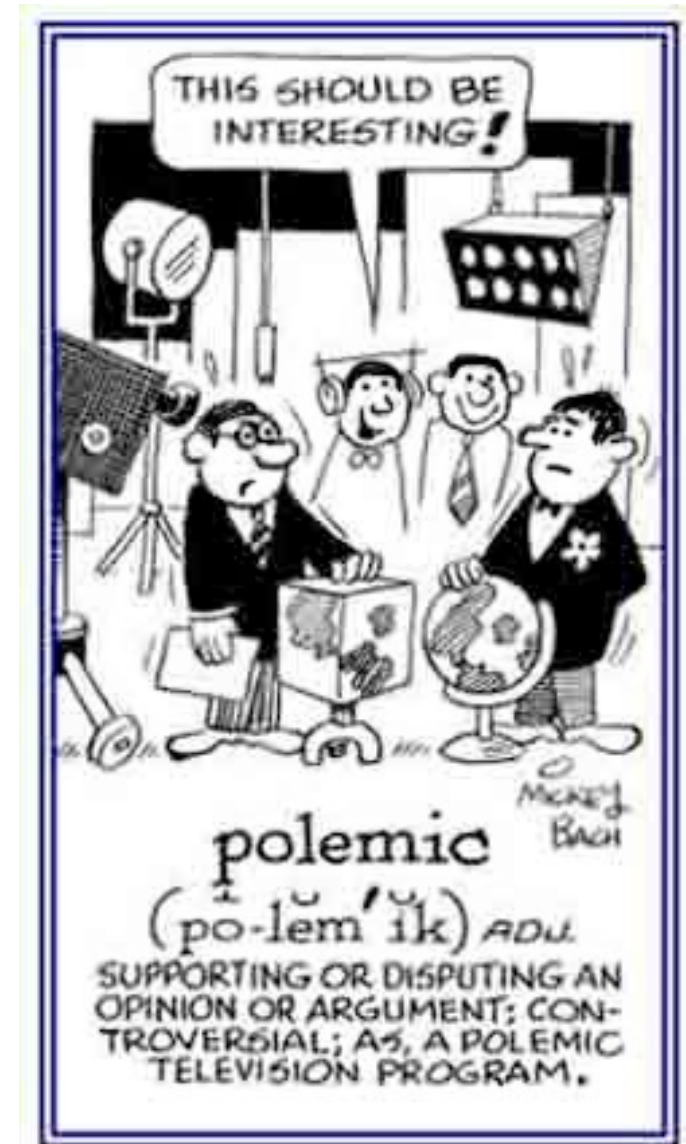
# Topic 2 – Architecture and Philosophy

- Abstraction
- Layering
- Layers and Communications
- Entities and Peers
- What is a protocol?
- Protocol Standardization
- The architects process
  - How to break system into modules
  - Where modules are implemented
  - Where is state stored
- Internet Philosophy and Tensions

# TRIGGER WARNING

- Philosophy,
- Bad Analogies, and
- RANTS verging  
on POLEMIC

Will follow.....



# Abstraction Concept

A mechanism for breaking down a problem

*what not how*

- eg Specification *versus* implementation
- eg Modules in programs

Allows replacement of implementations without affecting system behavior

*Vertical versus Horizontal*

*“Vertical”* what happens in a box *“How does it attach to the network?”*

*“Horizontal”* the communications paths running through the system

**Hint:** paths are built (*“layered”*) on top of other paths

# Computer System Modularity

Partition system into modules & abstractions:

- Well-defined interfaces give flexibility
  - **Hides** implementation - can be freely changed
  - Extend functionality of system by adding new modules
- E.g., libraries encapsulating set of functionality
- E.g., programming language + compiler abstracts away how the particular CPU works ...

# Computer System Modularity (cnt' d)

- Well-defined interfaces hide information
  - Isolate **assumptions**
  - Present high-level **abstractions**
- **But can impair performance!**
- Ease of implementation vs worse performance

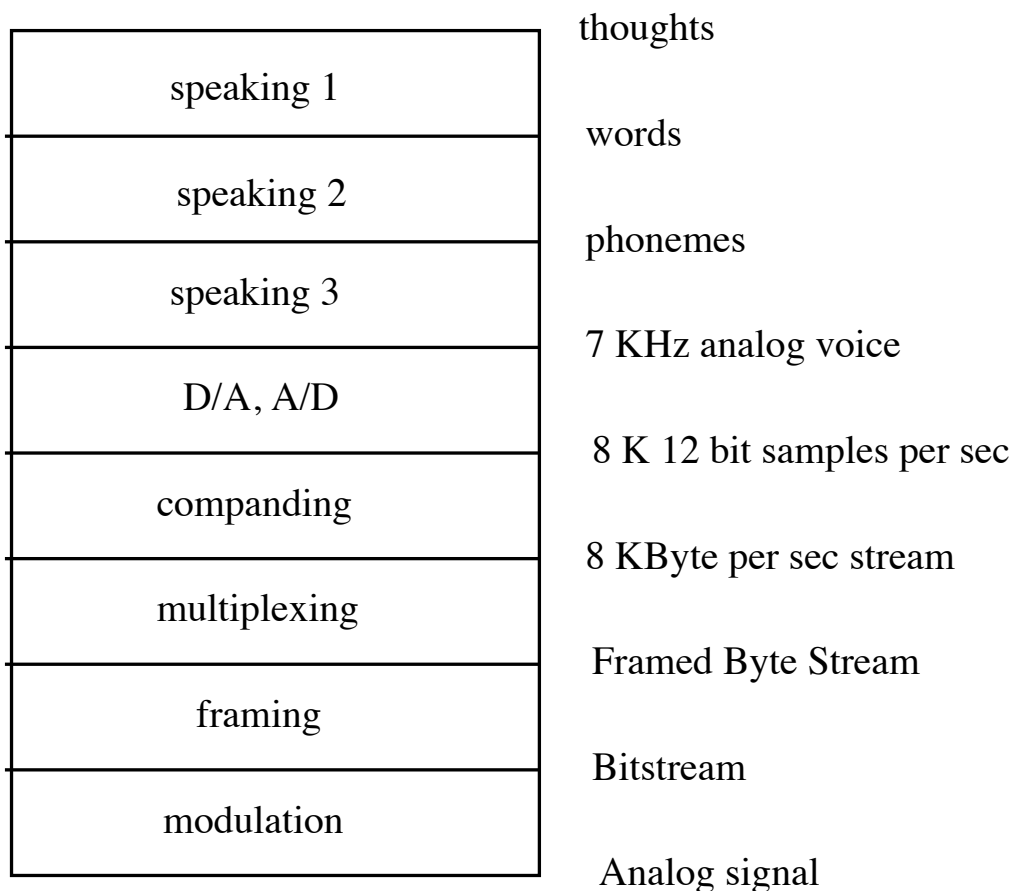
# Network System Modularity

Like software modularity, but:

- Implementation is distributed across many machines (routers and hosts)
- Must decide:
  - How to break system into modules
    - **Layering**
  - Where modules are implemented
    - **End-to-End Principle**
  - Where state is stored
    - **Fate-sharing**

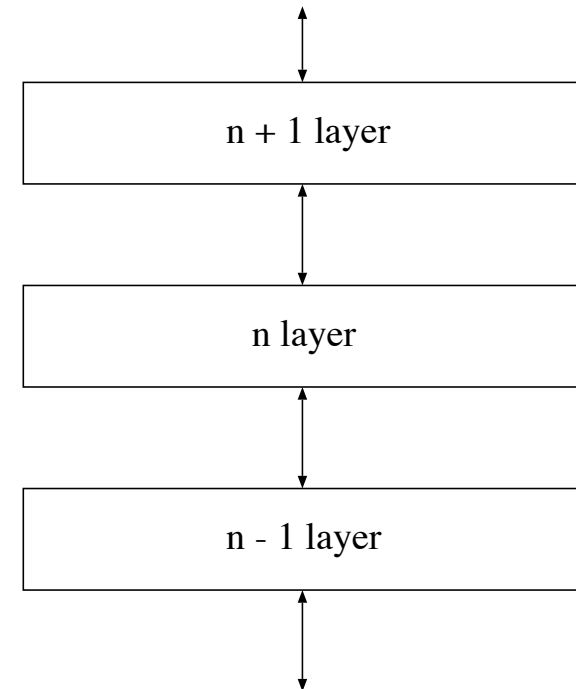
# Layering Concept

- A restricted form of abstraction: system functions are divided into layers, one built upon another
- Often called a *stack*; but **not** a data structure!



# Layers and Communications

- Interaction only between adjacent layers
- *layer n* uses services provided by *layer n-1*
- *layer n* provides service to *layer n+1*
- Bottom layer is physical media
- Top layer is application





# Entities and Peers

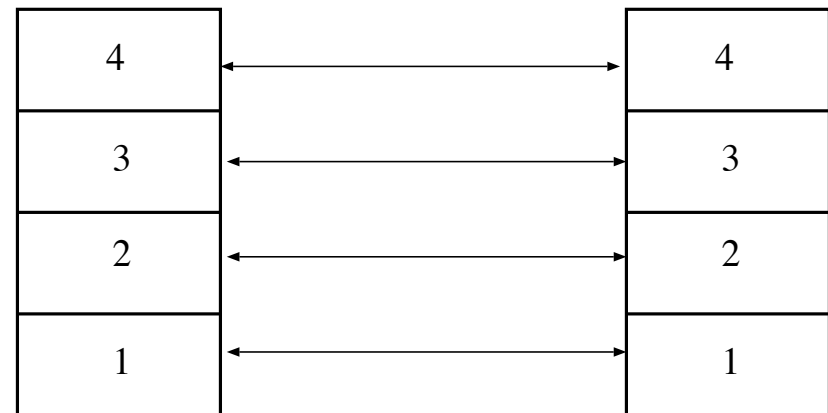
*Entity* – a *thing* (an independent existence)

Entities *interact* with the layers above and below

Entities *communicate* with *peer* entities

- same level but different place (eg different person, different box, different host)

Communications between peers is supported by entities at the lower layers



# Entities and Peers

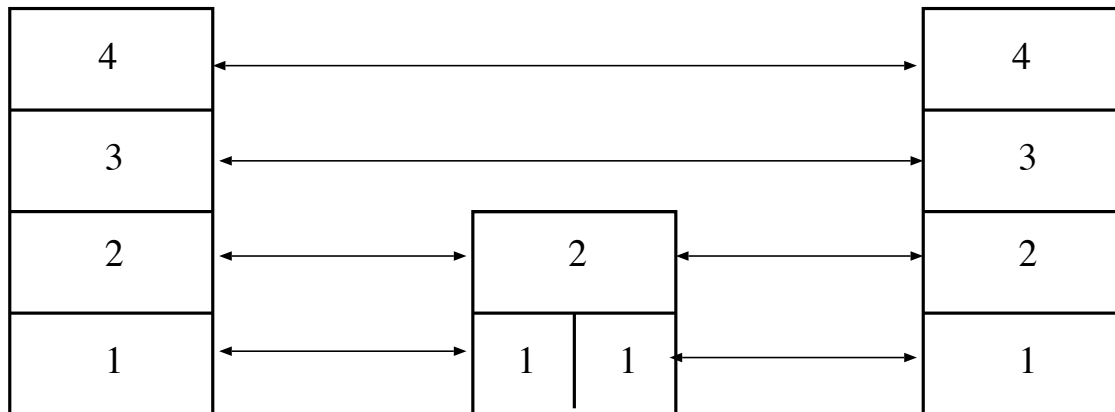
Entities usually do something useful

- Encryption – Error correction – Reliable Delivery
- Nothing at all is also reasonable

Not all communications is end-to-end

Examples for things in the middle

- IP Router – Mobile Phone Cell Tower
- Person translating French to English



# Layering and Embedding

In Computer Networks we often see higher-layer information embedded within lower-layer information

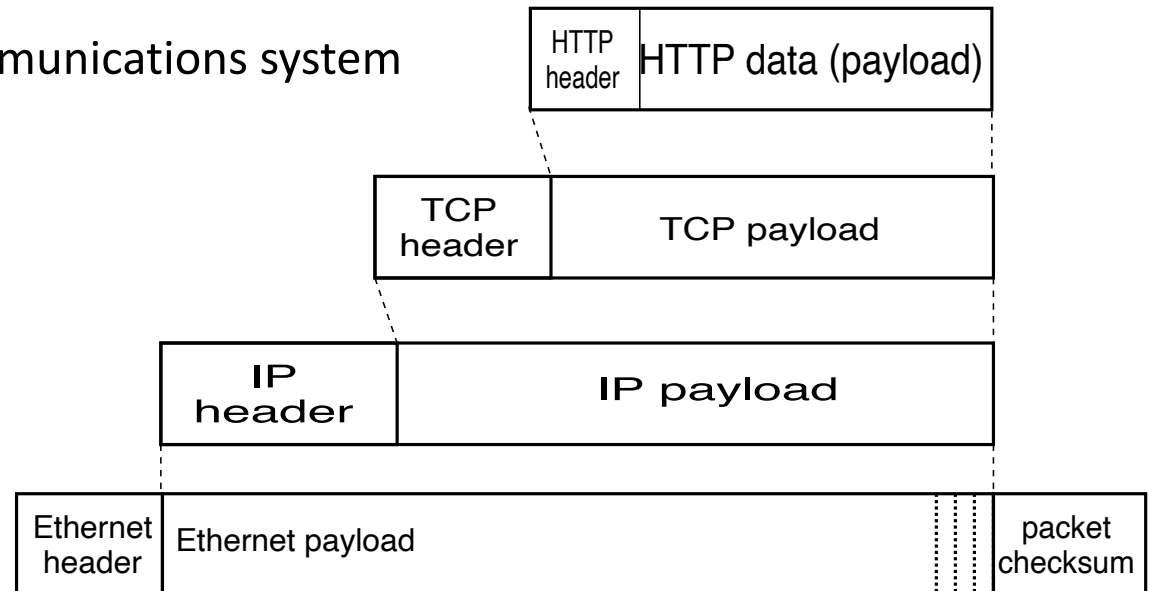
- Such embedding can be considered a form of layering
- Higher layer information is generated by stripping off headers and trailers of the current layer
- eg an IP entity only looks at the IP headers

***BUT embedding is not the only form of layering***

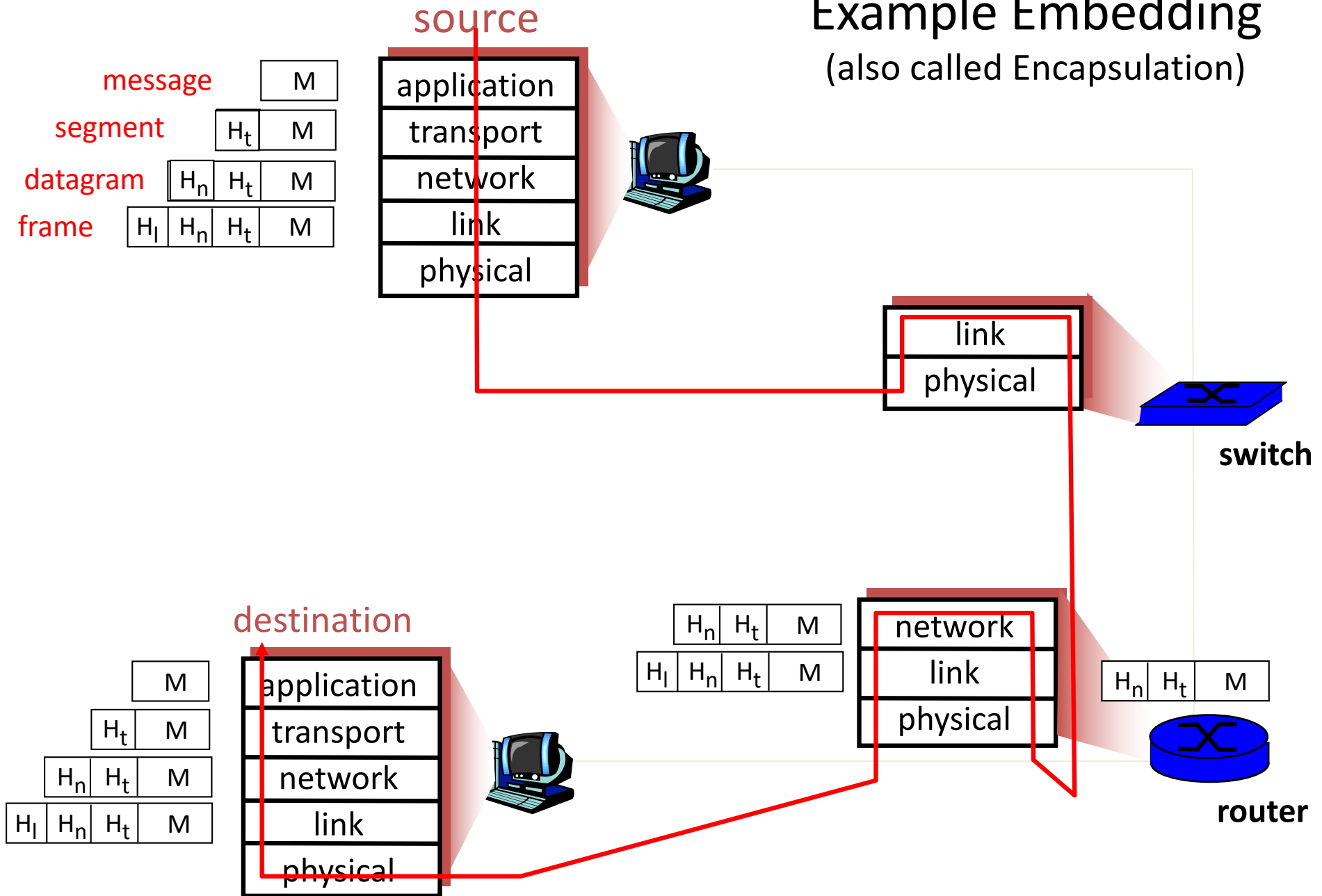
Layering is to help understand a communications system

**NOT**

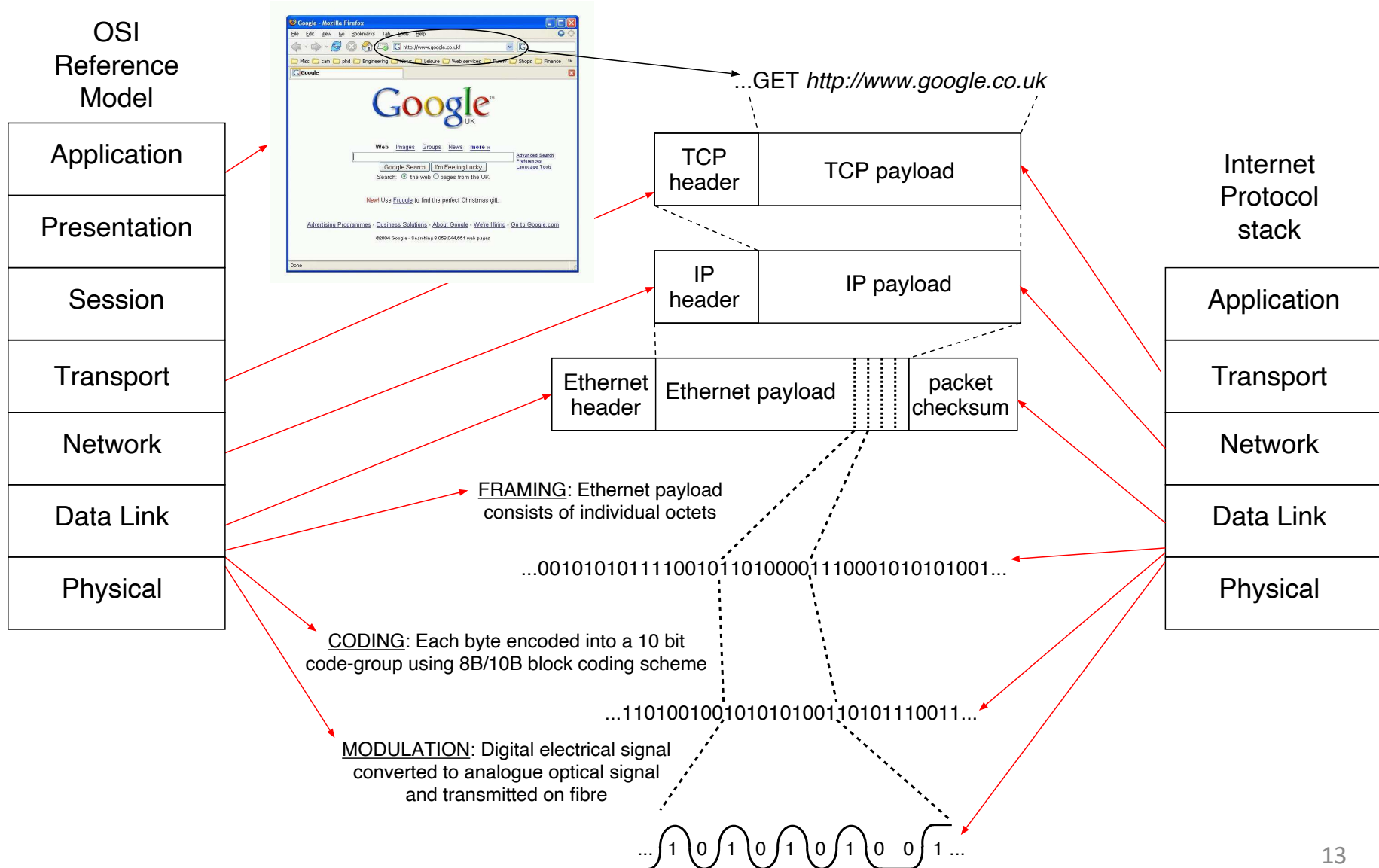
determine implementation strategy



# Example Embedding (also called Encapsulation)

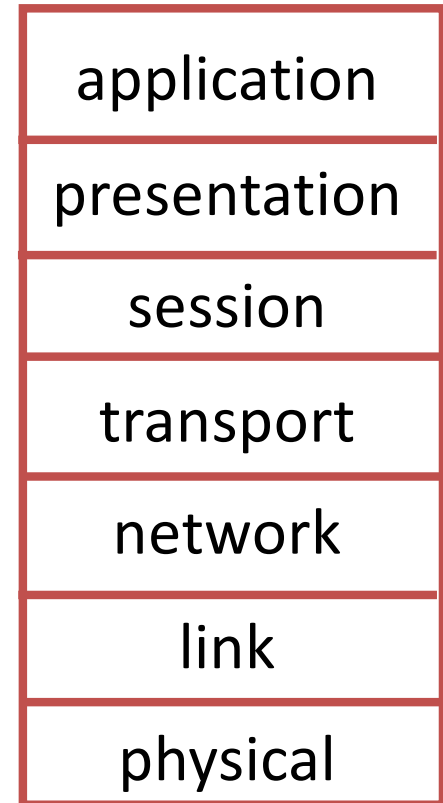


# Internet protocol stack *versus* OSI Reference Model

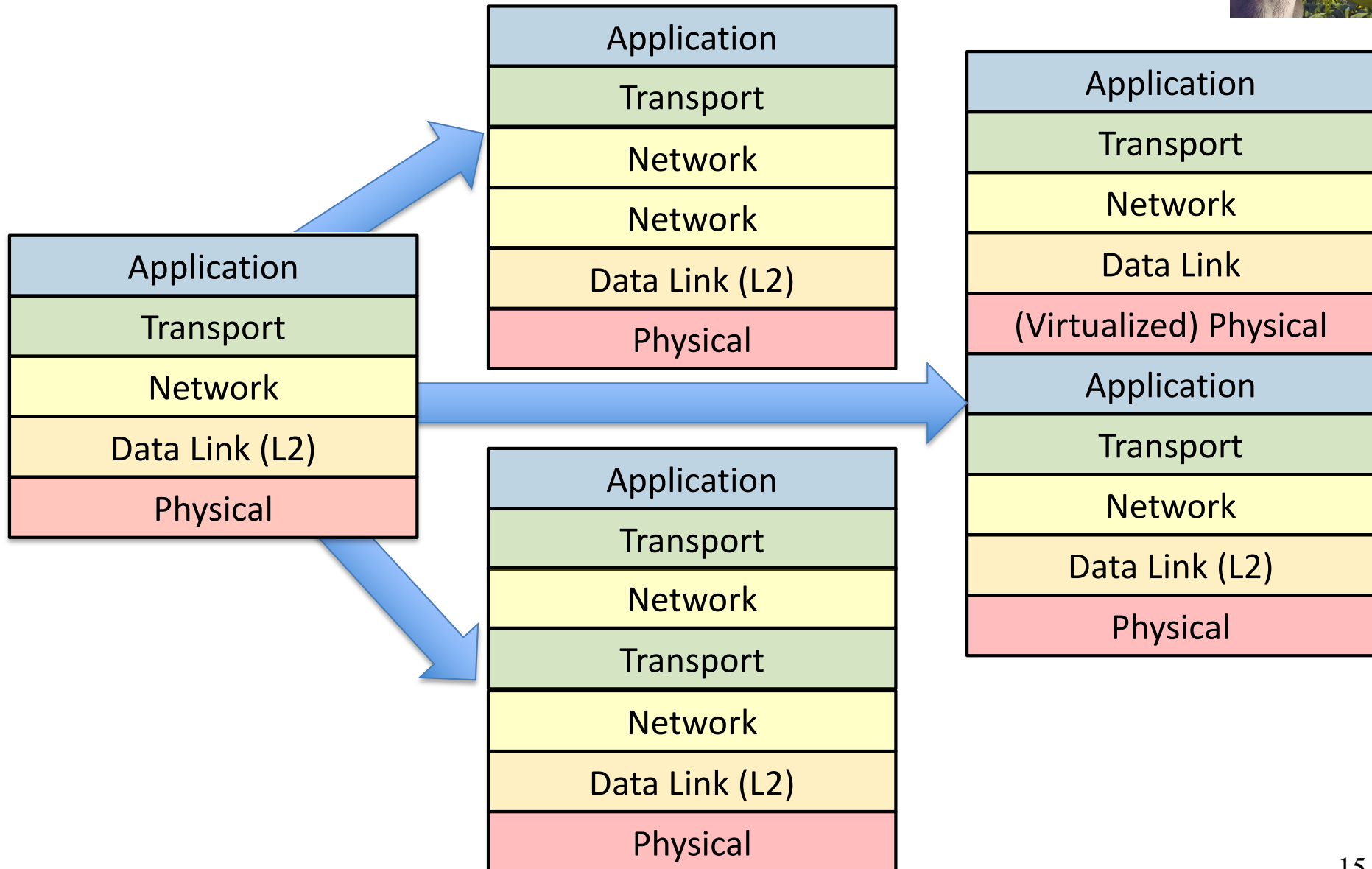


# ISO/OSI reference model

- **presentation:** allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- **session:** synchronization, checkpointing, recovery of data exchange
- Internet stack “missing” these layers!
  - these services, *if needed*, must be implemented in application



# Layers on Layers examples



# What is a protocol?

## human protocols:

- “what’s the time?”
- “I have a question”
- introductions

... specific msgs sent

... specific actions taken  
when msgs received, or  
other events

## network protocols:

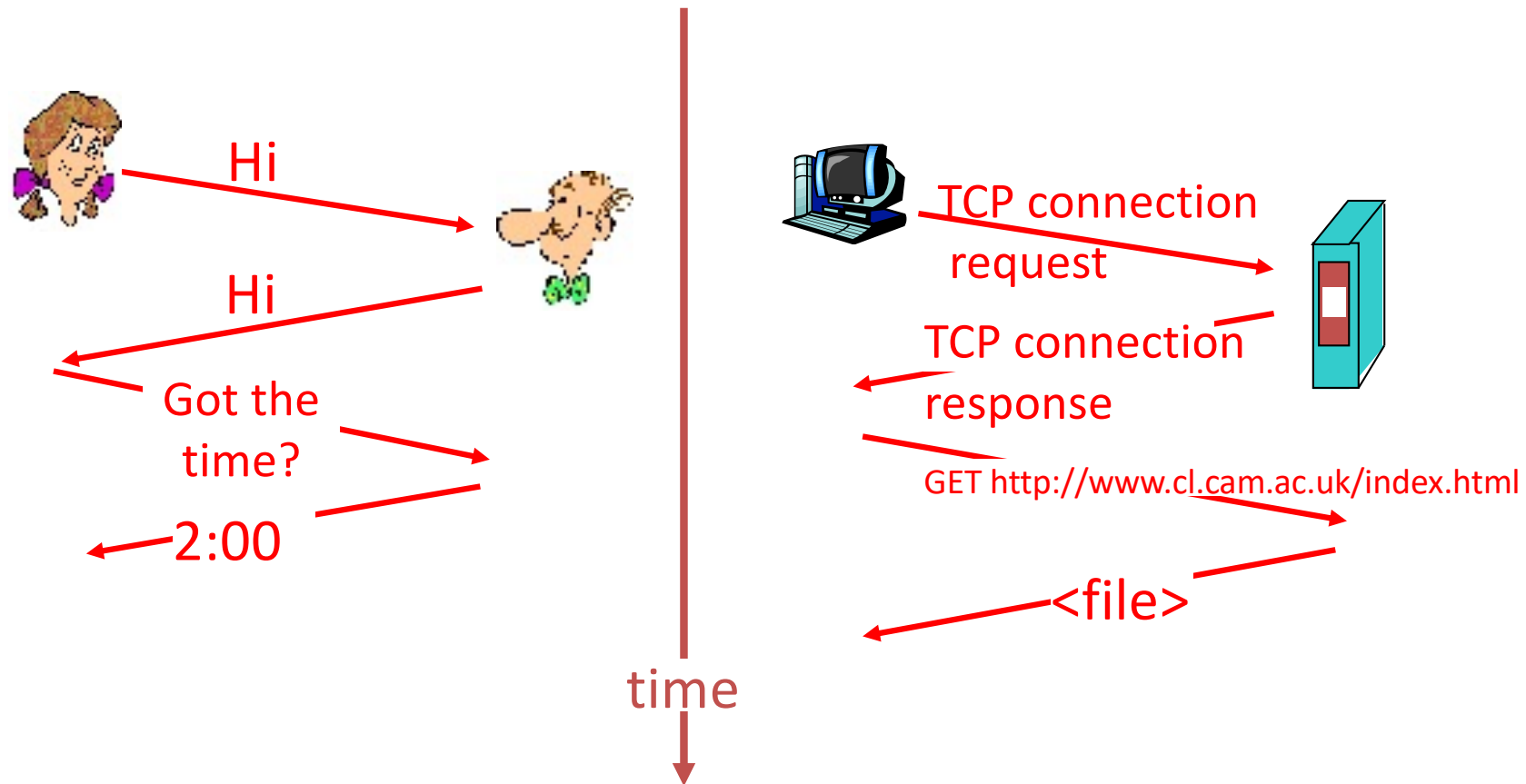
- machines rather than humans
- all communication activity in Internet governed by protocols

*protocols define format, order of msgs sent  
and received among network entities,  
and actions taken on msg transmission,  
receipt*



# What is a protocol?

a human protocol and a computer network protocol:



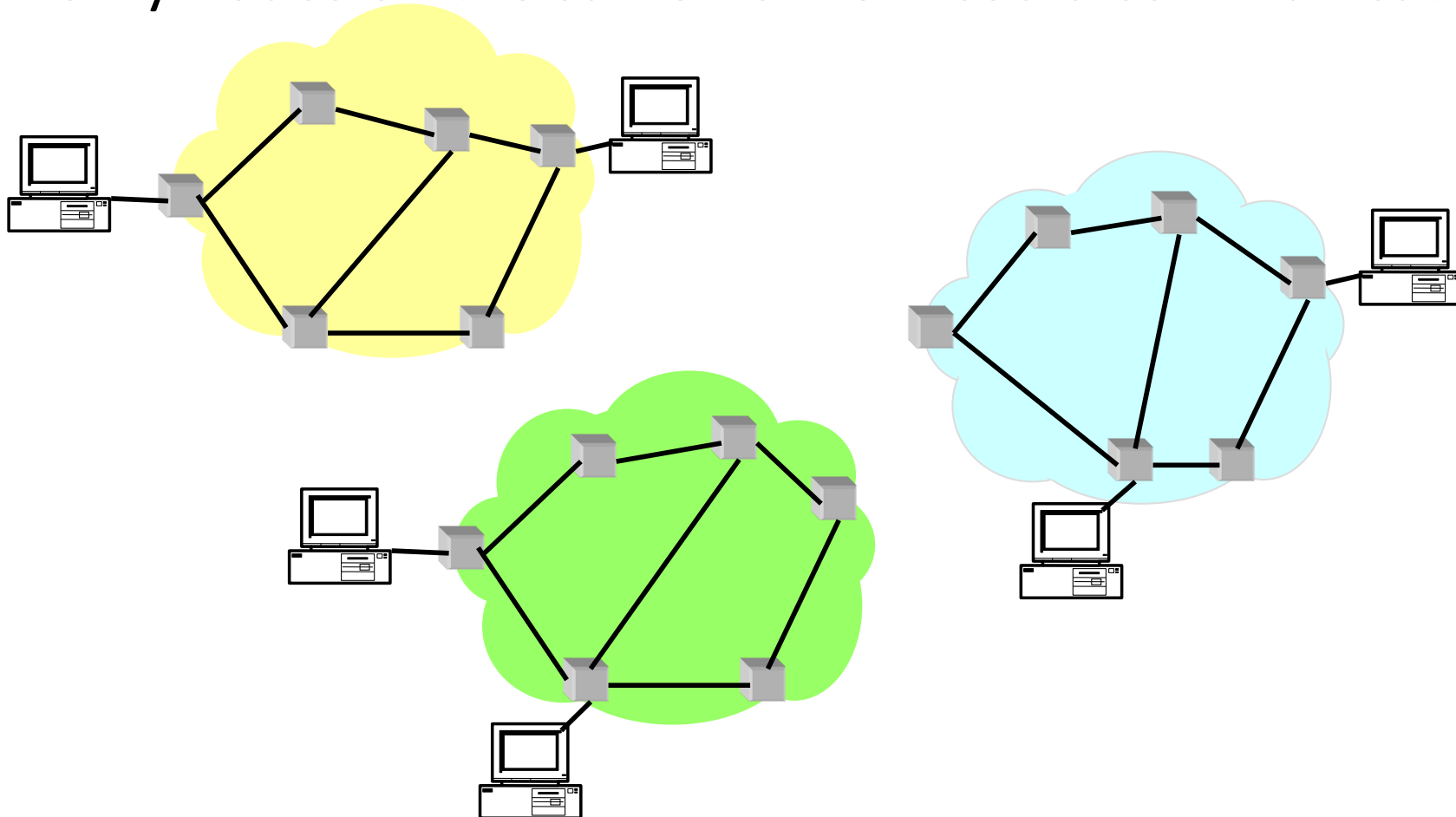
Q: Other human protocols?

# Protocol Standardization

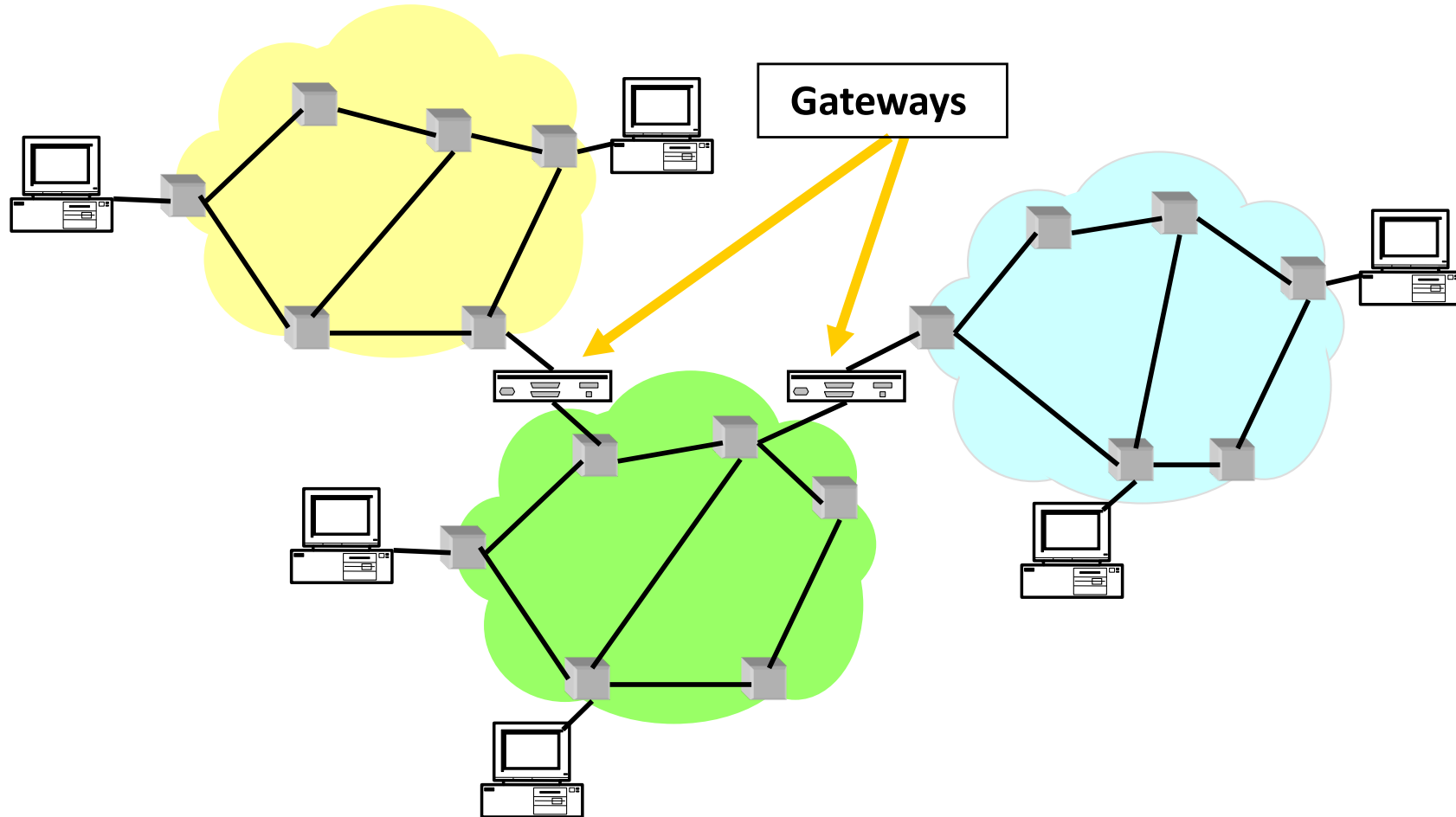
- All hosts must follow same protocol
  - Very small modifications can make a big difference
  - Or prevent it from working altogether
- This is why we have standards
  - Can have multiple implementations of protocol
- Internet Engineering Task Force (IETF)
  - Based on working groups that focus on specific issues
  - Produces “Request For Comments” (RFCs)
  - IETF Web site is ***<http://www.ietf.org>***
  - RFCs archived at ***<http://www.rfc-editor.org>***

# So many Standards Problem

- Many different packet-switching networks
- Each with its own Protocol
- Only nodes on the same network could communicate



# INTERNET Solution



# Internet Design Goals (Clark '88)

- **Connect existing networks**
- Robust in face of failures
- Support multiple types of delivery services
- Accommodate a variety of networks
- Allow distributed management
- Easy host attachment
- Cost effective
- Allow resource accountability

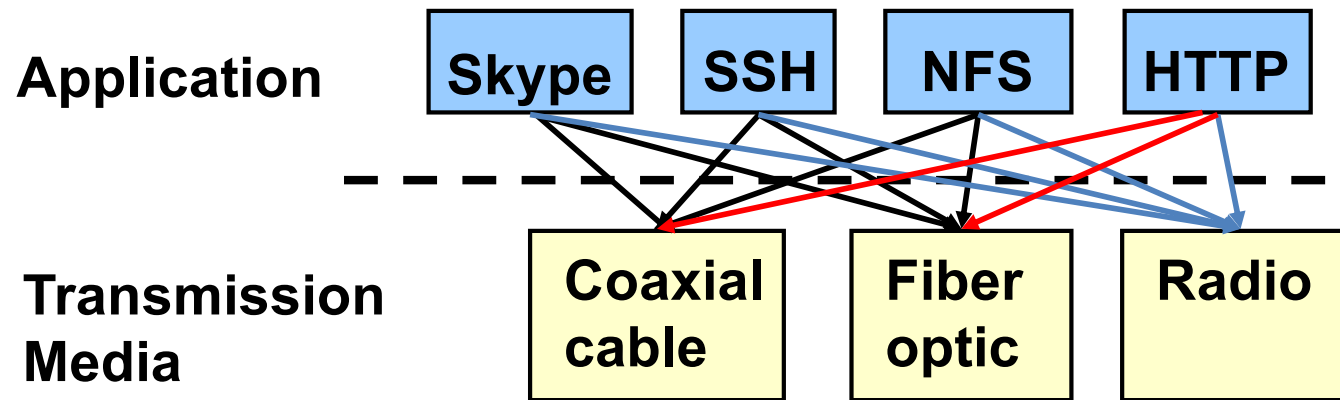
# Real Goals

Internet Motto

*We reject kings , presidents, and voting. We believe in rough consensus and running code.*“ – David Clark

- **Build something that works!**
- Connect existing networks
- Robust in face of failures
- Support multiple types of delivery services
- Accommodate a variety of networks
- Allow distributed management
- Easy host attachment
- Cost effective
- ~~Allow resource accountability~~

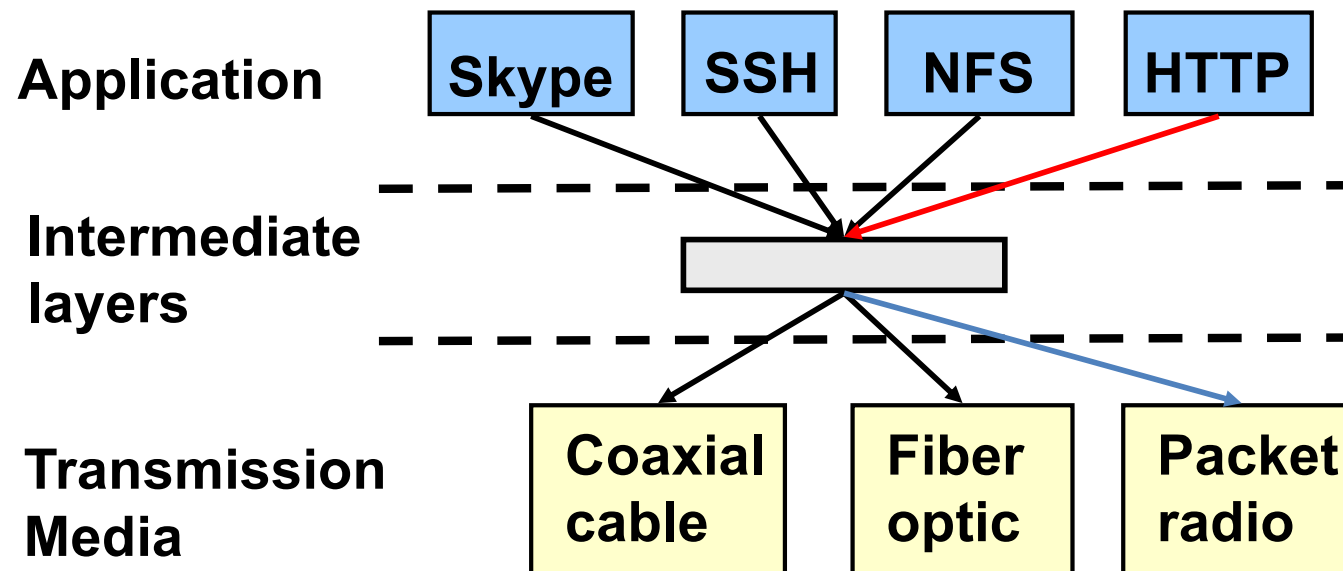
# A Multitude of Apps Problem



- Re-implement every application for every technology?
- No! But how does the Internet design avoid this?

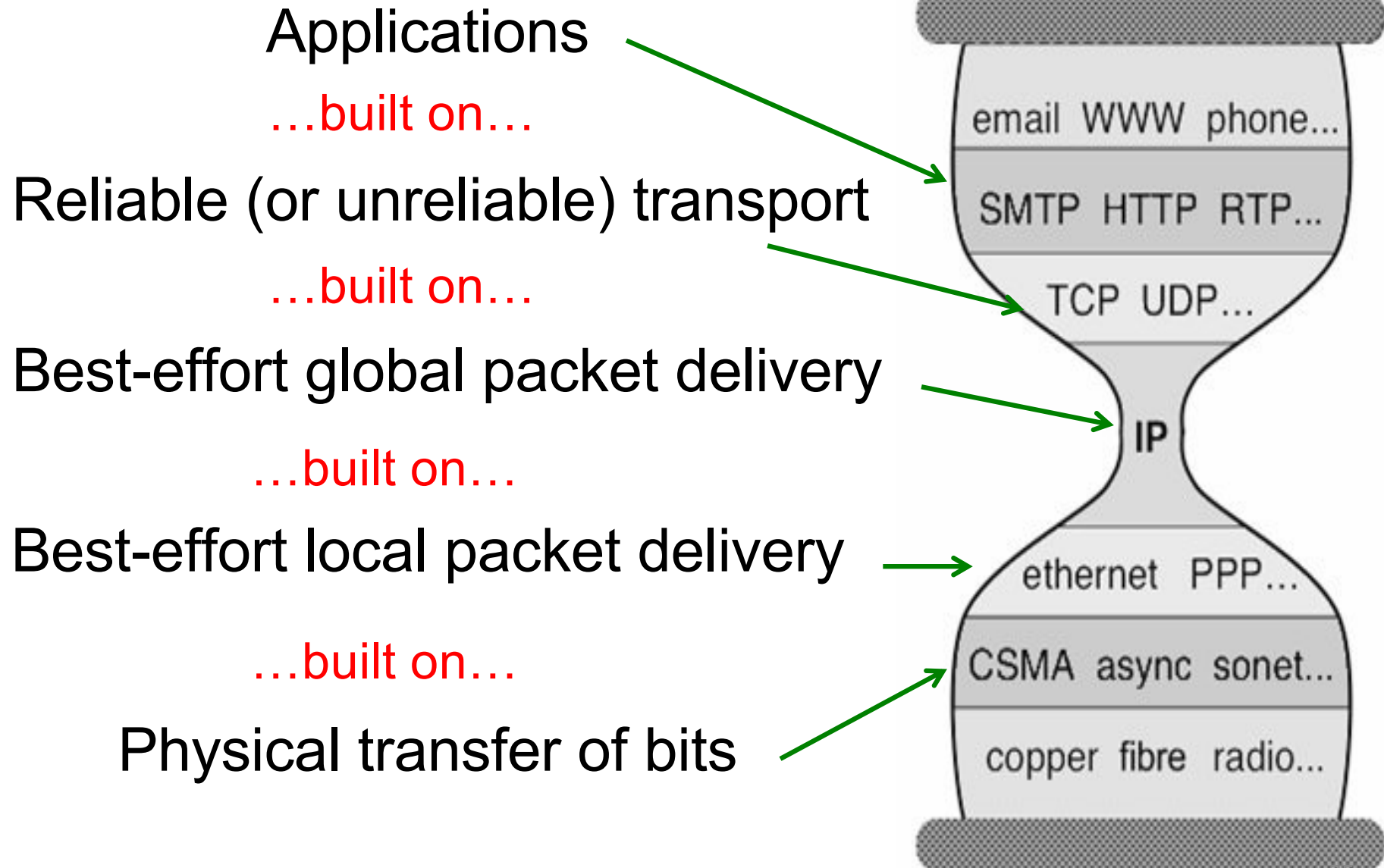
# Solution: Intermediate Layers

- Introduce intermediate layers that provide **set of abstractions** for various network functionality and technologies
  - A new app/media implemented only once
  - Variation on “add another level of indirection”



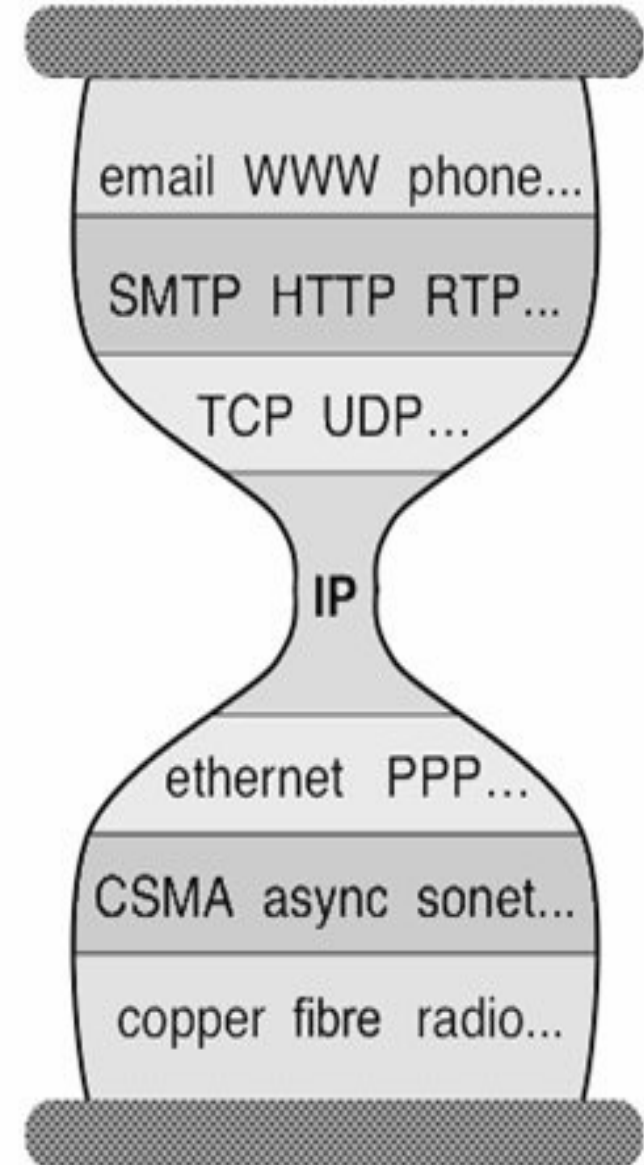


# In the context of the Internet



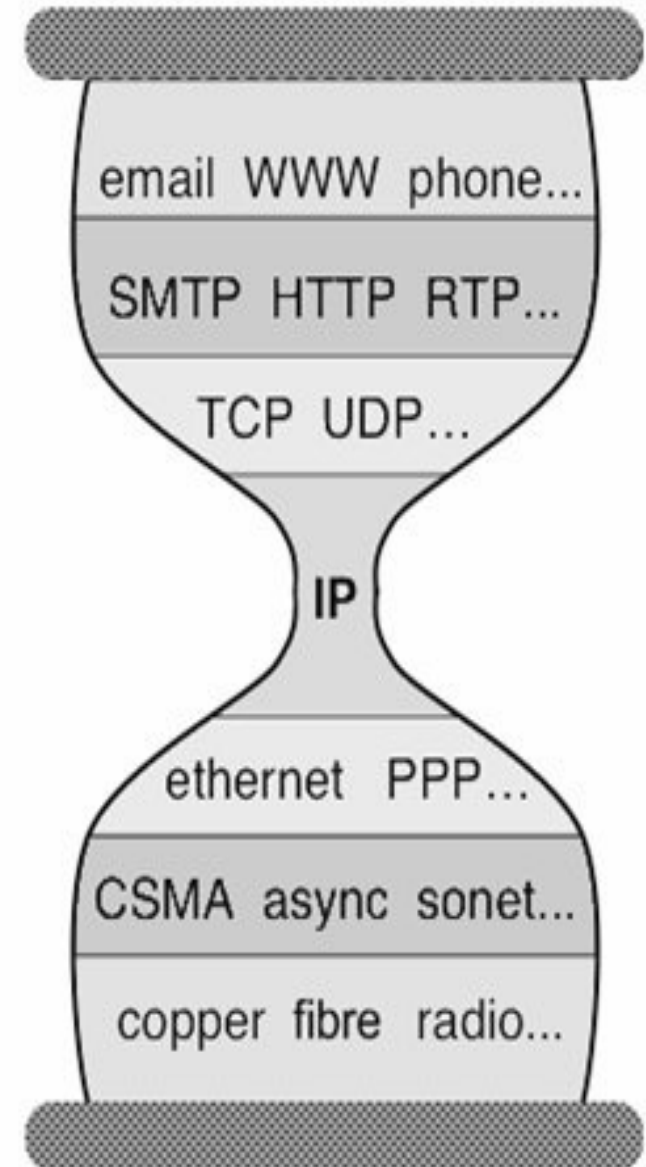
# Three Observations

- Each layer:
  - Depends on layer below
  - Supports layer above
  - Independent of others
- Multiple versions in layer
  - Interfaces differ somewhat
  - Components pick which lower-level protocol to use
- But only one IP layer
  - Unifying protocol



# Layering Crucial to Internet's Success

- Reuse
- Hides underlying detail
- Innovation at each level can proceed in parallel
- Pursued by very different communities



What are some of the drawbacks of protocols and layering?

# Drawbacks of Layering

- Layer N may duplicate lower layer functionality
  - e.g., error recovery to retransmit lost data
- Information hiding may hurt performance
  - e.g., packet loss due to corruption vs. congestion
- Headers start to get really big
  - e.g., typical TCP+IP+Ethernet is 54 bytes
- Layer violations when the gains too great to resist
  - e.g., TCP-over-wireless
- Layer violations when network doesn't trust ends
  - e.g., firewalls

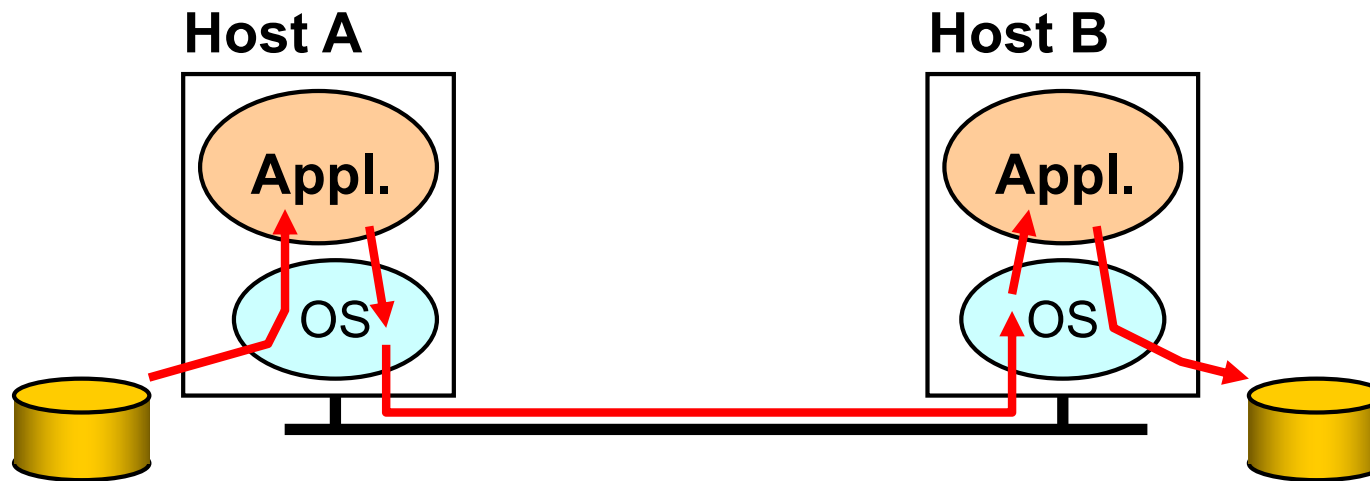
# Placing Network Functionality

- Hugely influential paper: “End-to-End Arguments in System Design” by Saltzer, Reed, and Clark ( ‘84)
  - articulated as the “End-to-End Principle” (E2E)
- Endless debate over what it means
- Everyone cites it as supporting their position  
(regardless of the position!)

# Basic Observation

- Some application requirements can only be correctly implemented **end-to-end**
  - reliability, security, *etc.*
- Implementing these in the network is hard
  - every step along the way must be fail proof
- Hosts
  - **Can** satisfy the requirement without network's help
  - **Will/must** do so, since they can't rely on the network

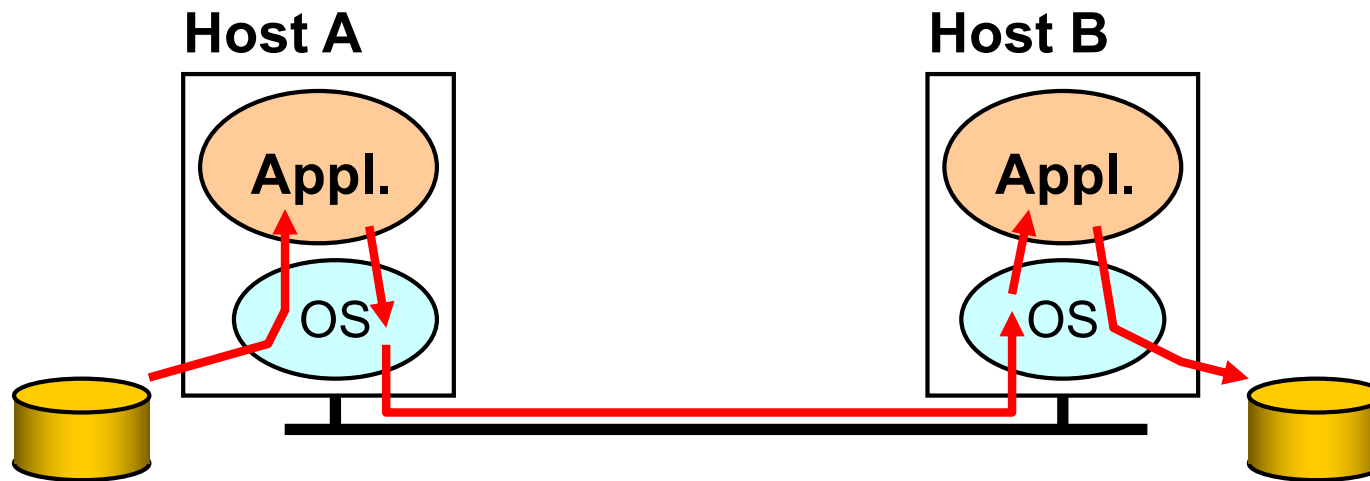
# Example: Reliable File Transfer



- Solution 1: make each step reliable, and string them together to make reliable end-to-end process



# Example: Reliable File Transfer



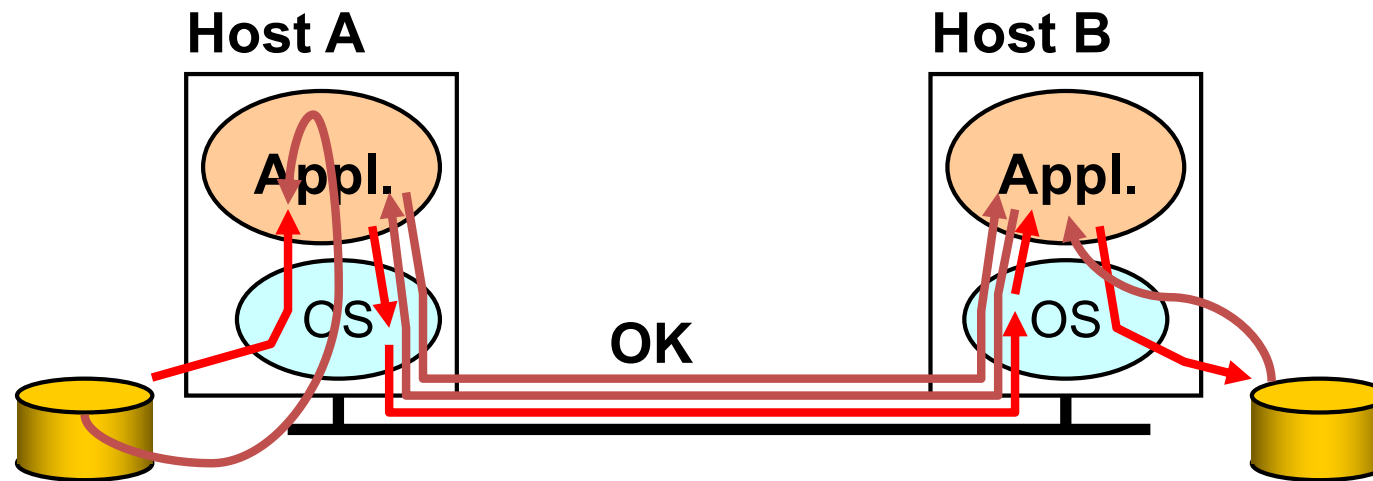
- Solution 1: make each step reliable, and string them together to make reliable end-to-end process

So what is the problem?

each component is 0.9 reliable

leads to total system failure of  $>0.4^*$

# Example: Reliable File Transfer



- Solution 1: make each step reliable, and string them together to make reliable end-to-end process
- Solution 2: end-to-end **check** and retry

# Discussion

- Solution 1 is incomplete
  - What happens if any network element misbehaves?
  - Receiver has to do the check anyway!
- Solution 2 is complete
  - Full functionality can be entirely implemented at application layer with no need for reliability from lower layers
- Is there any need to implement reliability at lower layers?

# Summary of End-to-End Principle

- Implementing functionality (e.g., reliability) in the network
  - Doesn't reduce host implementation complexity
  - Does increase network complexity
  - Probably increases delay and overhead on all applications even if they don't need the functionality (e.g. VoIP)
- However, implementing in the network can improve performance in some cases
  - e.g., consider a very lossy link

# “Only-if-Sufficient” Interpretation

- Don't implement a function at the lower levels of the system unless it can be completely implemented at this level
- *Unless you can relieve the burden from hosts, don't bother*

# “Only-if-Necessary” Interpretation

- Don't implement *anything* in the network that can be implemented correctly by the hosts
- Make network layer absolutely minimal
  - This E2E interpretation trumps performance issues
  - Increases flexibility, since lower layers stay **simple**

# “Only-if-Useful” Interpretation

- If hosts can implement functionality correctly, implement it in a lower layer **only** as a performance enhancement
- But do so only if it **does not impose burden** on applications that do not require that functionality

# We have some tools:

- Abstraction
- Layering
- Layers and Communications
- Entities and Peers
- Protocol as motivation
- Examples of the architects process
- Internet Philosophy and Tensions

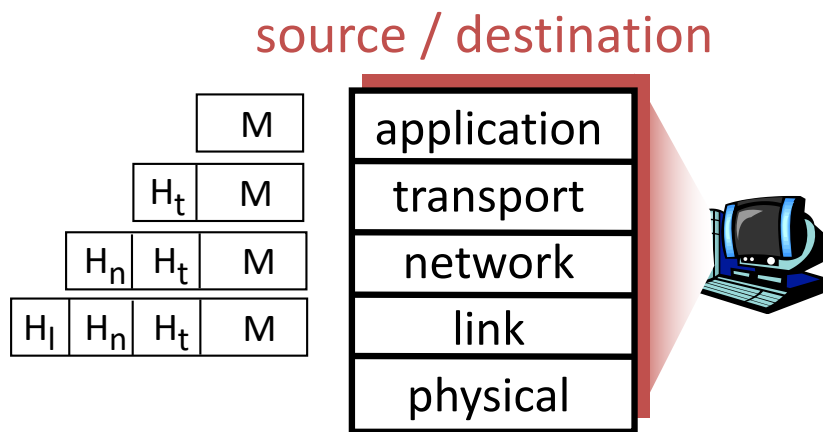


# Distributing Layers Across Network

- Layers are simple if only on a single machine
  - Just stack of modules interacting with those above/below
- But we need to implement layers across machines
  - Hosts
  - Routers (switches)
- What gets implemented where?

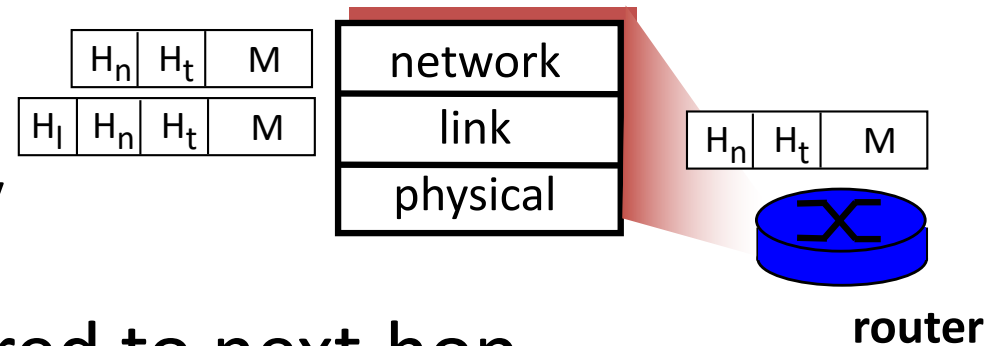
# What Gets Implemented on Host?

- Bits arrive on wire, must make it up to application
- Therefore, all layers must exist at the host



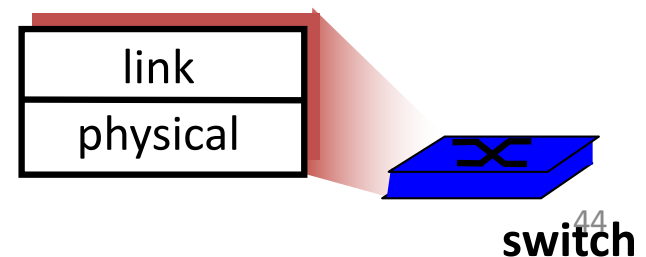
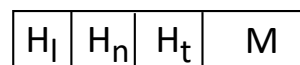
# What Gets Implemented on a Router?

- Bits arrive on wire
  - Physical layer necessary
- Packets must be delivered to next-hop
  - Datalink layer necessary
- Routers participate in global delivery
  - Network layer necessary
- Routers don't support reliable delivery
  - Transport layer (and above) **not** supported

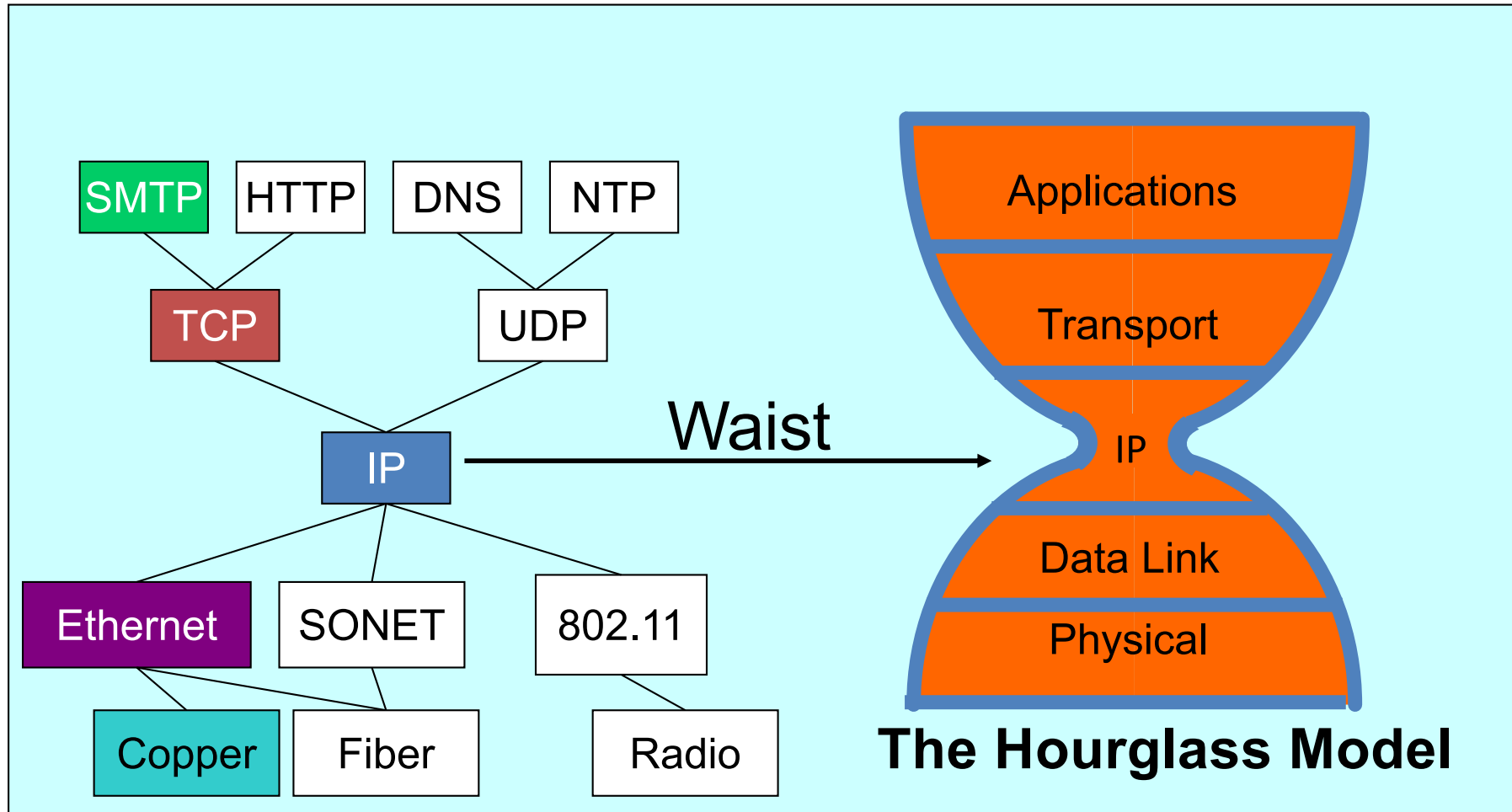


# What Gets Implemented on Switches?

- Switches do what routers do, except they don't participate in global delivery, just local delivery
- They only need to support Physical and Datalink
  - Don't need to support Network layer
- Won't focus on the router/switch distinction
  - Almost all boxes support network layer these days
  - Routers have switches but switches do not have routers

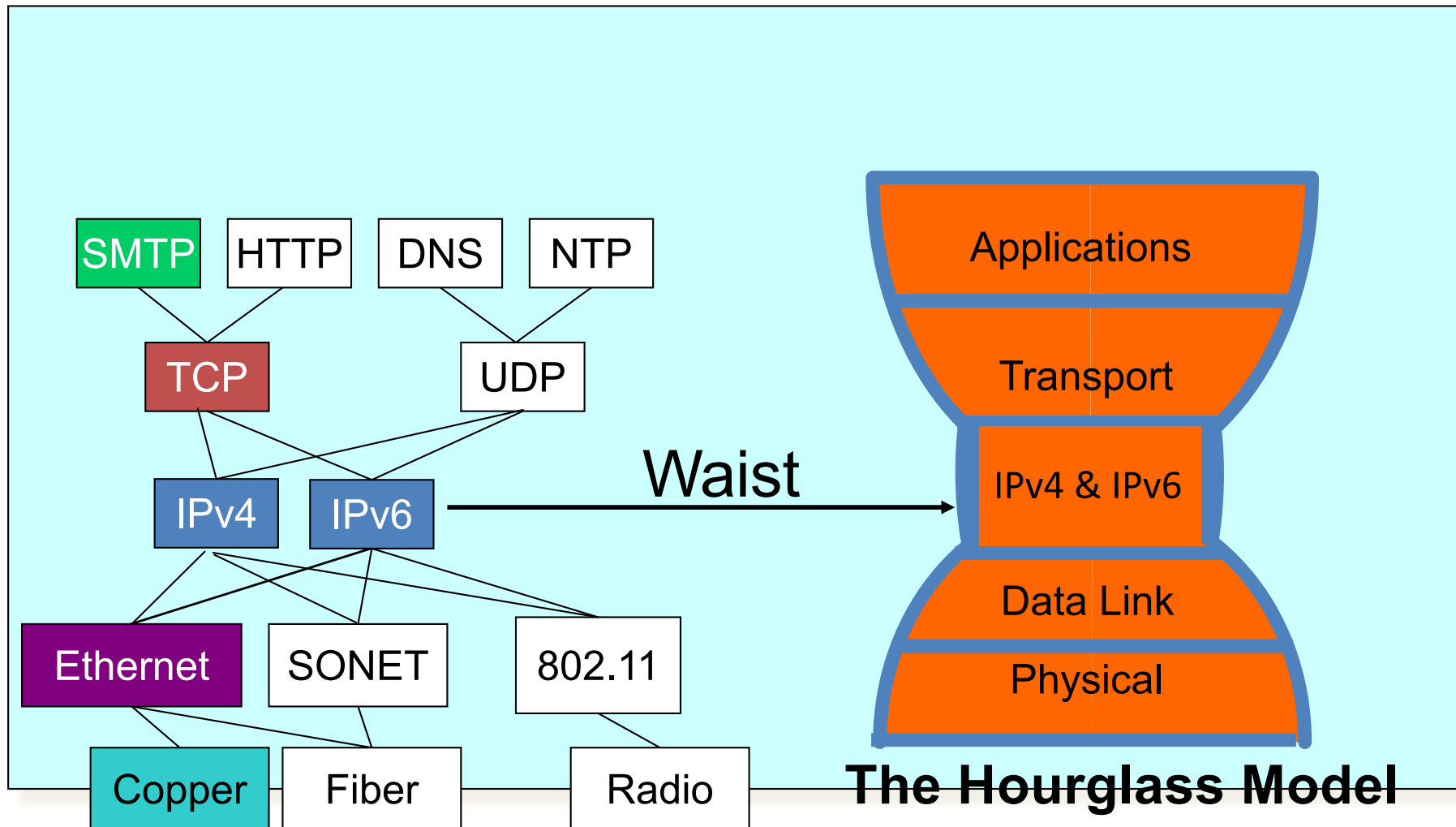


# The Internet *Hourglass*



There is just **one** network-layer protocol, **IP**.  
The “narrow waist” facilitates **interoperability**.

# The middle-age Internet *Hourglass*



There is just ~~one~~<sup>TWO</sup> network-layer protocol, IPv4 + v6  
The “narrow waist” facilitates interoperability(???)

# Alternative to Standardization?

- Have one implementation used by everyone
- Open-source projects
  - Which has had more impact, Linux or POSIX?
- Or just sole-sourced implementation
  - Skype, Signal, FaceTime, etc.