Kernels and Tracing

Lecture 2, Part 2: The Probe Effect
Prof. Robert N. M. Watson
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The **Probe Effect**

- The **probe effect** is the unintended alteration of system behaviour that arises from measurement
  - Software instrumentation is **active**: execution is changed
- DTrace minimises probe effect when unused...
  - ... but has a very significant impact when it is used
  - Disproportionate effect on probed events
- Potential perturbations:
  - Speed relative to other cores (e.g., lock hold times)
  - Speed relative to external events (e.g., timer ticks)
  - Microarchitectural effects (e.g., cache, branch predictor)
Probe effect example: dd(1) execution time

- Simple (naïve) microbenchmark – dd(1)
  - dd copies blocks from input to output
  - Copy 10M buffer from /dev/zero to /dev/null
  - ("Do nothing .. But do it slowly")
  - Execution time measured with /usr/bin/time
  - **Workload chosen to illustrate high overhead**

```
# dd if=/dev/zero of=/dev/null bs=10m count=1 status=none
```

- Simultaneously, run various DTrace scripts
  - Compare resulting execution times using ministat
  - Difference is probe effect (+/- measurement error)
Probe effect 1: memory allocation

- Using the dtmalloc provider, count kernel memory allocations:
  
  ```
  dtmalloc:::
  { @count = count(); }
  ```

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<tr>
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<td>0.22</td>
<td>0.21</td>
<td>0.20818182</td>
<td>0.0060302269</td>
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<td>0.21</td>
<td>0.2127272727</td>
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No difference proven at 95.0% confidence

- No statistically significant overhead at 95% confidence level
Probe effect 2: locking

- Using the lockstat provider, track kernel lock acquire, release:

  ```
  lockstat:::
  { @count = count(); }
  ```

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<td>0.21</td>
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Difference at 95.0% confidence
0.226364 +/- 0.00575196
108.734% +/- 2.76295%
(Student's t, pooled s = 0.0064667)

- **109% overhead** – 170K locking operations vs. 6 malloc() calls!
Probe effect 3: limiting to dd(1)?

• Limit the action to processes with the name dd:

```c
lockstat::: /execname == "dd"/
{ @count = count(); }
```

• Well, crumbs. Now **168% overhead**!
Probe effect 4: stack traces

• Gather more locking information in action – capture call stacks:

```c
lockstat::: { @stacks[stack()] = count(); }
lockstat::: /execname == "dd"/ { @stacks[stack()] = count(); }
```

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<td>1.30455 +/− 0.00970671</td>
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1.30455 +/− 0.00970671 | 626.638% +/− 4.66261%
What does this mean for us?

• Always think about the potential role of the probe effect when instrumenting a workload
  • E.g., avoid benchmarking while running DTrace ...
  • ... unless measuring or accounting for the probe effect

• Traced applications may behave (very) differently
  • E.g., more timer ticks will fire, affecting thread inverleaving
  • E.g., I/O will “seem faster” relative to computation, as latter may slow down due to probe effect

• Performance overheads may be disproportionate
  • E.g., if you instrument one way of doing things, but not another, and workloads have a different functional footprint

• Consider ways to decide if an analysis is representative
  • E.g., are the performance inflection points consistent even if absolute performance is lower?