The Network Stack (2)

Lecture 6, Part 1: TCP

Dr Robert N. M. Watson

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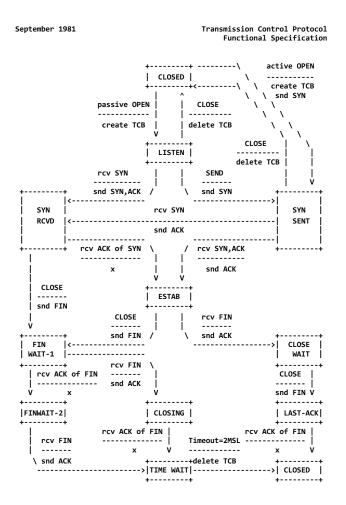
The Network Stack (2)

- The Transmission Control Protocol (TCP)
 - The TCP state machine
 - TCP congestion control
 - TCP implementations and performance
 - The evolving TCP stack
 - Lab 3 on TCP
- Wrapping up the Advanced Operating Systems lecture series

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Lecture 6, Part 1
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Lecture 6, Part 3
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The Transmission Control Protocol (TCP)



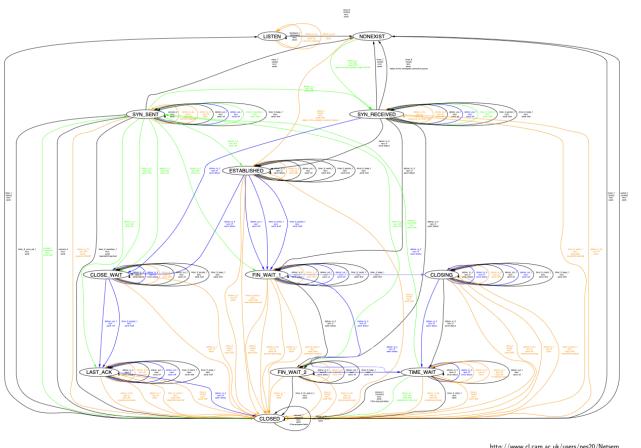
TCP Connection State Diagram Figure 6.

- V. Cerf, K. Dalal, and C. Sunshine, *Transmission Control Protocol (version 1)*,
 INWG General Note #72,
 December 1974.
- In practice: J. Postel, Ed.,
 Transmission Control Protocol: Protocol Specification, RFC
 793, September, 1981.



Compare to Bishop, et al (2005)

TCP: an approximation to the real state diagram



http://www.cl.cam.ac.uk/users/pes20/Netsem

What Is This?

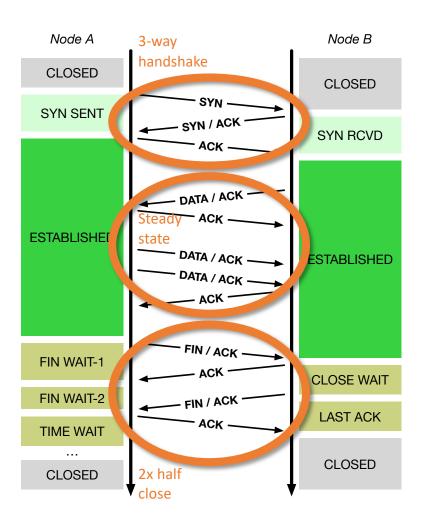
hat RST is set and r indicating RST is clear. Tra

Transition Rules		
	olane. F	Successful abortion class of a synchronised socket
	olore, T	Secondally close the last file descriptor for a socket in the CLOSED.
		SYNLSENT or SYNLRECEMED Males.
	class.#	Successfully close the last file descriptor for a listening TCP sucket.
	oresect_f	Bugin connection establishment by creating a SYN and trying to measure it on half's catories
	connect(Fall: socket has pending ever
	deliver_in_I	Passive open: receive SYN, send SYN,ACK
	ddiner-in-D	For a littering socket, receive and drap a had datagram and other generate a RST segment or ignore it. Ones the incoming segment if the sociat's passes of incomplete connections in full.
	deliver in 2	Completion of active open (in SYN-SENT receive SYN-ACK and send ACH) or simultaneous open (in SYN-SENT receive SYN and send SYN-ACK)
	deliver.in.24	Receive bad or boring datagram and RST or ignore for SYN_SENT suches
	deliver.in.2	Receive data, FINs, and AOKs in a connected state
	deliver_in_25	Receive data after process has gove away
	4dioer.in.2c	Receive stupid ACK or LAND DoS in SYNLRSCERVED state
	deliver-in-6	Receive and drop (cliently) a same segment that matches a CLOSED nucleit.
	deliver-in-7	Receive RST and cap non-(CLOSED; USTER; SYN SENT; SYN RECEIVED; TIME WWY) socket
	deliner.in.7s	Receive RST and sap SYN RECEIVED socket
	deliver.in.75	Receive RST and igners for USTEN socket
	delicer.in.7c	Receive RST and ignore for SYN.SENT(unacceptable ach) or TME WART model.
	deliver in 72	Receive RST and zap SYN.SENT (acceptable ack) secket
	4diver.in.#	Receive SYN is son-ICLOSED: LISTEN: SYN-SENT:
		TIME WAITH GAR
	deliver_in_9	Receive SYN in TIME WAIT state if there is no matching USTEN
		seclet or sequence number has not increased
	deliver_out_1	Common case TCP colput
	histenf	Successfully put sucket in LISTEN state
	listen-fr	Successfully put social in the LISTEN state from any nen- (CLOSED LISTEN) state on FreeESD
	skatdron, J	Shut down read or write half of TCP connection
	sucket. J	Successfully nature a new file descriptor for a fresh socket.
	timer At Prod. J	2°MSL time supires
	Sinur_M.com.est./	connection establishment times expires
		/FW:WAIT.2 timer eights
	Sinuer_M.keep. J	keegafive timer expires
	Newer-M-persist-f	persist timer eques
	biner_N_reput_f	



Steve Bishop, Matthew Fairbairn, Michael Norrish, Peter Sewell, Michael Smith, and Keith Wansbrough. Rigorou Specification and Conformance Testing Techniques for Network Protocols, as Applied to TCP, UDP, and Sockets. Proceedings of SIGCOMM 2005, ACM, 2005.

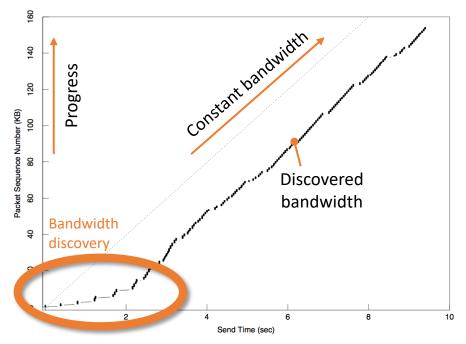
TCP principles and properties



- Assumptions: Network may delay, (reorder), drop, corrupt IP packets
- TCP implements reliable, ordered, stream transport protocol over IP
- Three-way handshake: SYN / SYN-ACK / ACK (mostly!)
- Steady state
 - Sequence numbers ACK'd
 - Round-Trip Time (RTT) measured to time out loss
 - Data retransmitted on loss
 - Flow control via advertised window size in ACKs
 - Congestion control ('fairness') detects congestion via loss (and, recently, via delay: BBR)
- NB: "Half close" allows communications in one direction to end while the other continues

TCP congestion control and avoidance

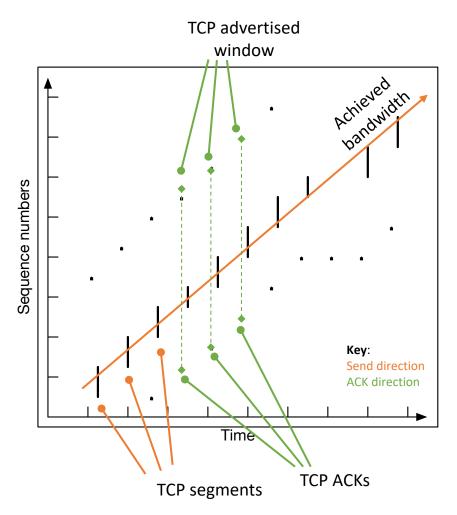
Figure 4: Startup behavior of TCP with Slow-start



Same conditions as the previous figure (same time of day, same Suns, same network path, same buffer and window sizes), except the machines were running the $4.3^+\mathrm{TCP}$ with slow-start. No bandwidth is wasted on retransmits but two seconds is spent on the slow-start so the effective bandwidth of this part of the trace is $16~\mathrm{KBps}$ — two times better than figure 3. (This is slightly misleading: Unlike the previous figure, the slope of the trace is $20~\mathrm{KBps}$ and the effect of the 2 second offset decreases as the trace lengthens. E.g., if this trace had run a minute, the effective bandwidth would have been $19~\mathrm{KBps}$. The effective bandwidth without slow-start stays at $7~\mathrm{KBps}$ no matter how long the trace.)

- 1986 Internet CC collapse
 - 32Kbps → **40bps**
- Van Jacobson, SIGCOMM 1988
 - Don't send more data than the network can handle!
 - Conservation of packets via ACK clocking
 - Exponential retransmit timer, slow start, aggressive receiver ACK, dynamic window sizing on congestion, and (later) ABC
- ECN (RFC 3168), ABC (RFC 3465), Compound (Tan, et al, INFOCOM 2006), Cubic (Rhee and Xu, ACM OSR 2008), BBR (Cardwell, ACM Queue 2016)

TCP time/sequence graphs (Van Jacobson)



- Extracted from TCP packet traces (e.g., via tcpdump)
- Visualize windows, congestion response, buffering, RTT, etc:
 - X: Time
 - Y: Sequence number
- We can extract this data from the network stack directly using DTrace
 - Allows correlation/plotting with respect to other variables / events
 - E.g., TCP and socket-buffer state
- TCP time/sequence diagrams have since been extended to represent additional information
 - E.g., SACK (selective acknowledgement) blocks

