

The Process Model (1)

L41 Lecture 3, Part 1: The Process Model

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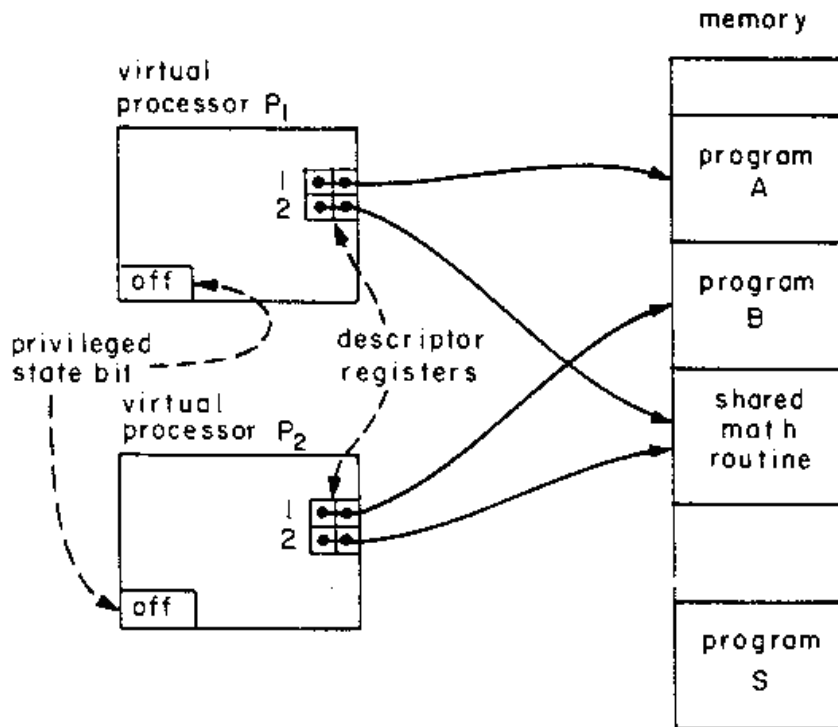
This time: The process model

- The process model and its evolution
- Brutal (re, pre)-introduction to VM
- Where do programs come from?

} Lecture 3, Part 1

} Lecture 3, Part 2

The Process Model: 1970s foundations



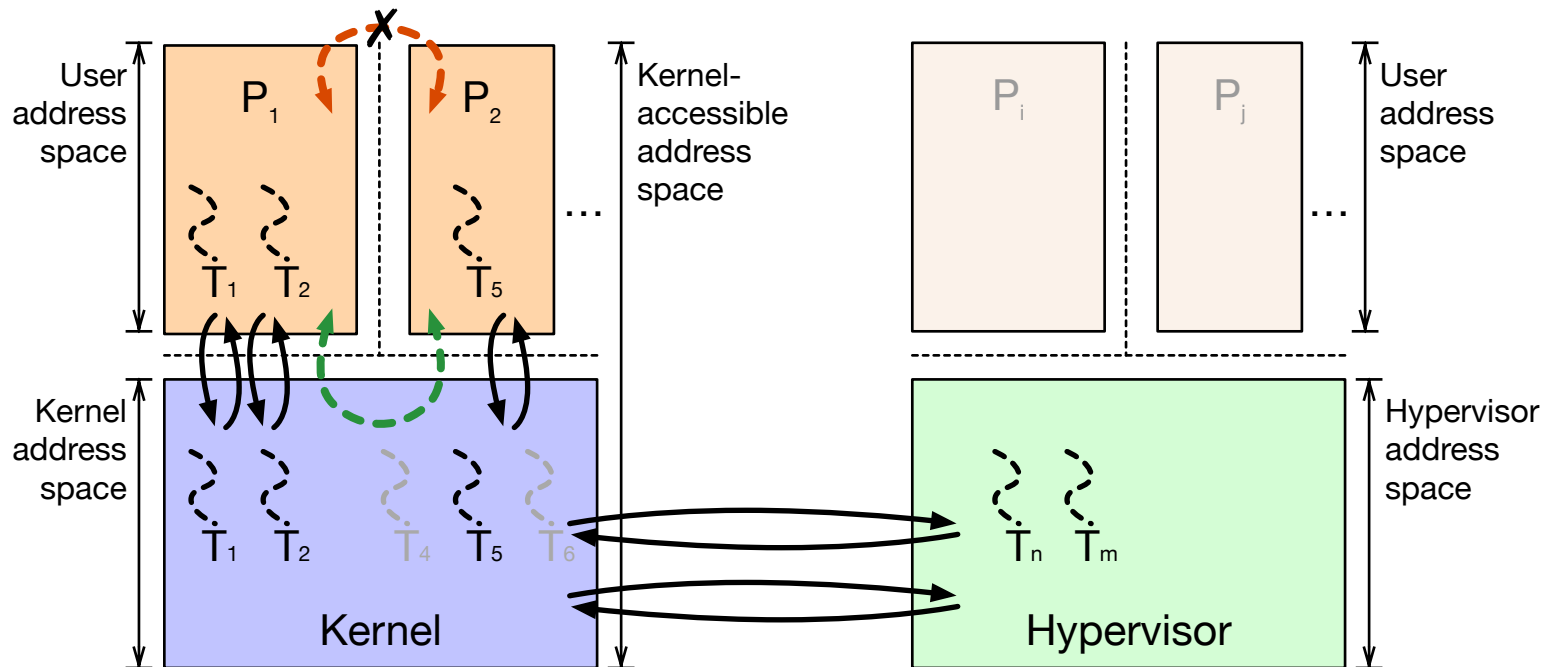
- Saltzer and Schroeder, **The Protection of Information in Computer Systems**, SOSP'73, October 1973. (CACM 1974)
- **Multics process model**
 - 'Program in execution'
 - **Process isolation** bridged by **controlled communication** via **supervisor** (kernel)
- Hardware foundations
 - Supervisor mode
 - Memory segmentation
 - Trap mechanism
- Hardware protection rings (Schroeder and Saltzer, 1972)

The process model: today - concept

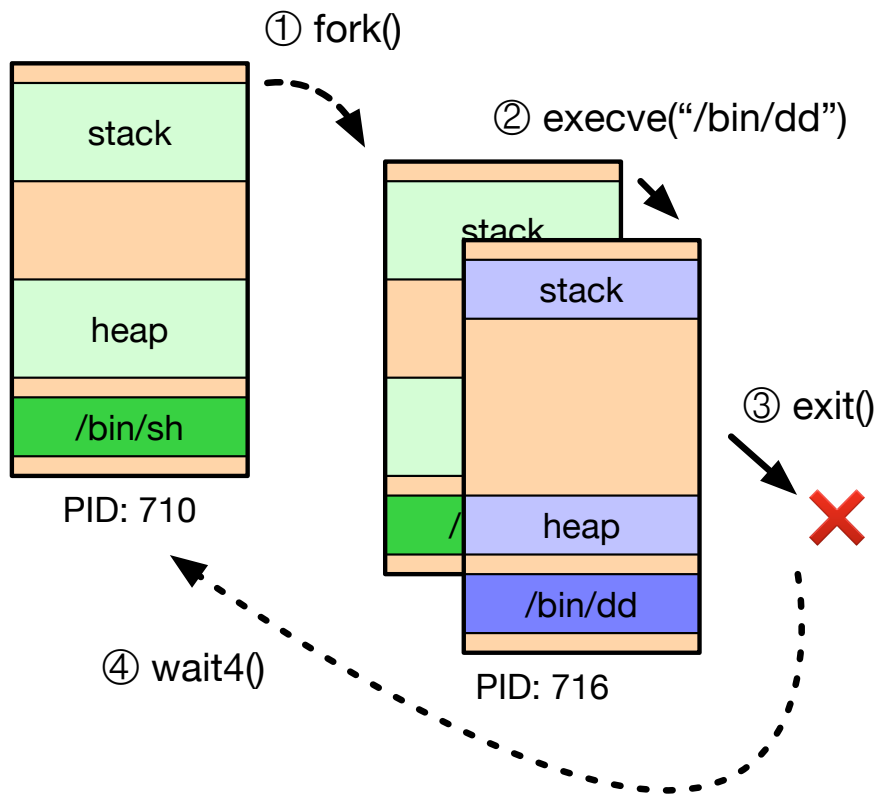
- ‘Program in execution’
 - **Process** \approx address space
 - **Threads** execute code
 - Unique instance of global variables, etc.
 - Isolated failure domain
- Unit of **resource accounting**
 - Open files, memory, ...
- Unit of privilege
 - Process credentials – UID, OS privileges, MAC, RBAC, ...
 - NB: Increasing support for per-thread credentials
- Recently: Inverted App-OS trust model
 - Third-party applications cannot trust the OS ...
 - E.g., Trustzone, SGX, ...

The process model today: isolation and controlled communication

- Hardware foundations for isolation
 - Rings control MMU, I/O, etc.
 - MMU to construct mutually exclusive **virtual address spaces**
 - Context switched **threads of control**
- Hardware foundations for controlled communication
 - Interaction via **traps**: system calls, page faults, ...
 - MMU to construct **shared memory**

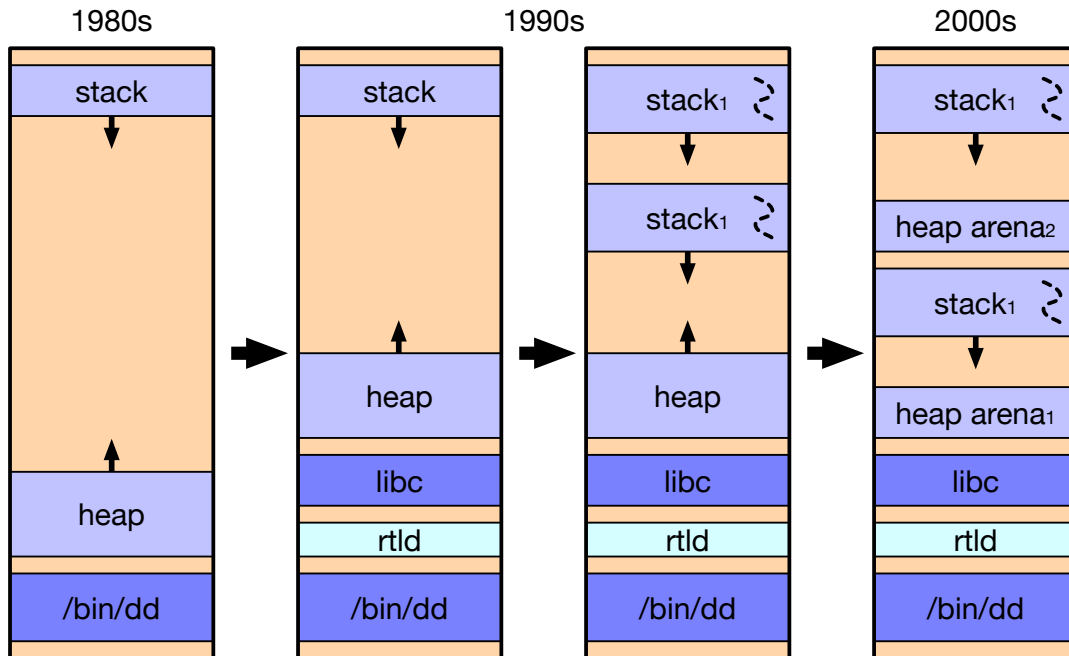


The UNIX process life cycle



- **fork()**
 - Child inherits address space and other properties
 - Program prepares process for new binary (e.g., `stdio`)
 - Copy-on-Write (COW)
- **execve()**
 - Kernel replaces address space, loads new binary, starts execution
- **exit()**
 - Process can terminate self (or be terminated)
- **wait4()** (et al)
 - Parent can await exit status
- NB: **posix_spawn()**

Evolution of the process model



- **1980s:** Code, heap, and stack
- **1990s:** Dynamic linking, threading
- **2000s:** Scalable memory allocators implement multiple **arenas** (e.g., as in `jemalloc`)
- Co-evolution with virtual memory (VM) research
 - Acetta, et al: *Mach* microkernel (1986)
 - Navarro, et al: *Superpages* (2002)