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## COMPUTER SCIENCE TRIPOS Part IA – 2015 – Paper 1

## 1 Foundations of Computer Science (LCP)

This question has been translated from Standard ML to OCaml

*O*-notation, lists, binary trees, functional arrays (a) Write brief notes about a tree representation of functional arrays, subscripted by positive integers according to their representation in binary notation. How efficient are the lookup and update operations? [6 marks]

Answer: The underlying data structure is the binary tree. A location in the tree is found by starting at the root, testing whether the subscript is even or odd, and descending into the left or right subtree, respectively; this process terminates when 1 is reached. Here is the code for lookup:

```
exception Subscript
let rec sub k = function
| Lf -> raise Subscript
| Br (v, t1, t2) ->
    if k = 1 then v
    else if k mod 2 = 0 then
        sub (k / 2) t1
    else
        sub (k / 2) t2
```

Lookup and update both take  $O(\log n)$  time, where n is the size of the array, because the representation guarantees balancing. The update operation naturally copies only the path from the root to the updated node, rather than the entire tree.

(b) Write an OCaml function arrayoflist to convert the list  $[x_1; \ldots; x_n]$  to the corresponding functional array having  $x_i$  at subscript position *i* for  $i = 1, \ldots, n$ . Your function should not call the update operation. [6 marks]

*Answer:* The point is to realise the tree structure directly, rather than repeatedly updating. Here is a straightforward solution:

```
let rec revalts ys zs = function
| [] -> (List.rev ys, List.rev zs)
| [x] -> (List.rev (x::ys), List.rev zs)
| x1::x2::xs -> revalts (x1::ys) (x2::zs) xs
let alts = revalts [] []
let rec arrayoflist = function
| [] -> Lf
| x::xs ->
    let (evens, odds) = alts xs in
    Br (x, arrayoflist evens, arrayoflist odds)
There is an elegant solution based on the following "cons" operation for Braun trees:
```

let rec tcons v = function
| Lf -> Br (v, Lf, Lf)
| Br (w, t1, t2) -> Br (v, tcons w t2, t1)

(c) Consider the task of finding out which elements of an array satisfy the predicate  $\mathbf{p}$ , returning the corresponding subscript positions as a list. For

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example, the list [2; 3; 6] indicates that these three designated array elements, and no others, satisfy **p**. Write an OCaml functional to do this for a given array and predicate, returning the subscripts in increasing order. [8 marks]

```
let rec merge xs (ys : int list) =
  match xs, ys with
  | [], ys -> ys
  | xs, [] -> xs
  | x::xs, y::ys ->
      if x<=y then
        x::(merge xs (y::ys))
      else
        y::(merge (x::xs) ys)
let rec mfilter p = function
| Lf -> []
| Br (x, t1, t2) ->
    let ks = merge (List.map (fun k -> 2 * k) (mfilter p t1))
                    (List.map (fun k \rightarrow 2 * k + 1) (mfilter p t2))
    in
      if p x then
        1 :: ks
      else
        ks
```

All OCaml code must be explained clearly and should be free of needless complexity.

Answer: The algorithm is a straightforward recursion. Using **merge** delivers a sorted result. A solution that returns an unsorted result, combined with a sorting function, is likely to lose marks due to inelegance and inefficiency.