

DIVISION  
THEOREM  
AND  
ALGORITHM

## The division theorem and algorithm

**Theorem 43 (Division Theorem)** *For every natural number  $m$  and positive natural number  $n$ , there exists a unique pair of integers  $q$  and  $r$  such that  $q \geq 0$ ,  $0 \leq r < n$ , and  $m = q \cdot n + r$ .*

## The division theorem and algorithm

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**Definition 44** *The natural numbers  $q$  and  $r$  associated to a given pair of a natural number  $m$  and a positive integer  $n$  determined by the Division Theorem are respectively denoted  $\text{quo}(m, n)$  and  $\text{rem}(m, n)$ .*

PROOF OF Theorem 43:

Uniqueness      Let

$$\begin{aligned} m &= q_1 \cdot n + r_1 & 0 \leq r_1 < n \\ m &= q_2 \cdot n + r_2 & 0 \leq r_2 < n \end{aligned} \quad (*)$$

RTP  $r_1 = r_2$  and  $q_1 = q_2$

From  $(*)$ , we have  $m \equiv r_1 \pmod{n}$      $0 \leq r_1 < n$   
and  $m \equiv r_2 \pmod{n}$      $0 \leq r_2 < n$

Therefore, by a previous proposition,  $r_1 = r_2$ .  
Moreover,  $q_1 n = q_2 n$  and, by cancellation,  $q_1 = q_2$ .

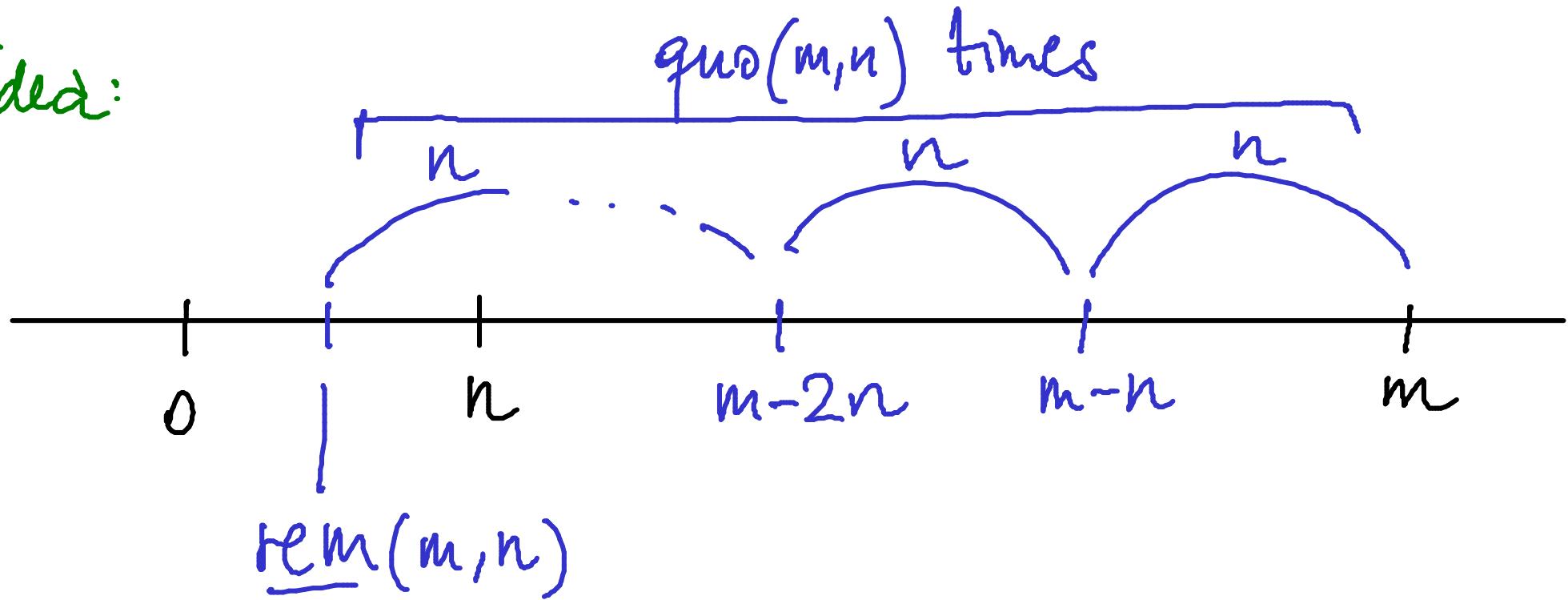
Given a natural number  $m$  and a positive integer  $n$ , it remains to show that there are natural numbers  $\underline{\text{quo}}(m, n)$  and  $\underline{\text{rem}}(m, n)$ , the latter below  $n$ , such that

$$m = \underline{\text{quo}}(m, n) \cdot n + \underline{\text{rem}}(m, n)$$

We will in fact compute them by means of the

DIVISION ALGORITHM

Idea:



That is:  $\underline{\text{quo}}(m,n) = \begin{cases} 0 & \text{if } m < n \\ 1 + \underline{\text{quo}}(m-n, n) & \text{else} \end{cases}$

$\underline{\text{rem}}(m,n) = \begin{cases} m & \text{if } m < n \\ \underline{\text{rem}}(m-n, n) & \text{else} \end{cases}$

## The Division Algorithm in ML:

```
fun divalg( m , n )
= let
  fun diviter( q , r )
  = if r < n then ( q , r )
    else diviter( q+1 , r-n )
in
  diviter( 0 , m )
end
```

```
fun quo( m , n ) = #1( divalg( m , n ) )
```

```
fun rem( m , n ) = #2( divalg( m , n ) )
```

# Computation Tree

$$\underline{\text{divalg}}(m, n) = \underline{\text{diviter}}(0, m)$$

$$m < n \quad / \quad \backslash \quad m \geq n$$

$$(0, m)$$

$$\underline{\text{diviter}}(1, m-n)$$

$$m-n < n \quad / \quad \backslash \quad m-n \geq n$$

$$(1, m-n)$$

$$\dots$$

$$\underline{\text{diviter}}(q, r)$$

$$r < n \quad / \quad \backslash \quad r \geq n$$

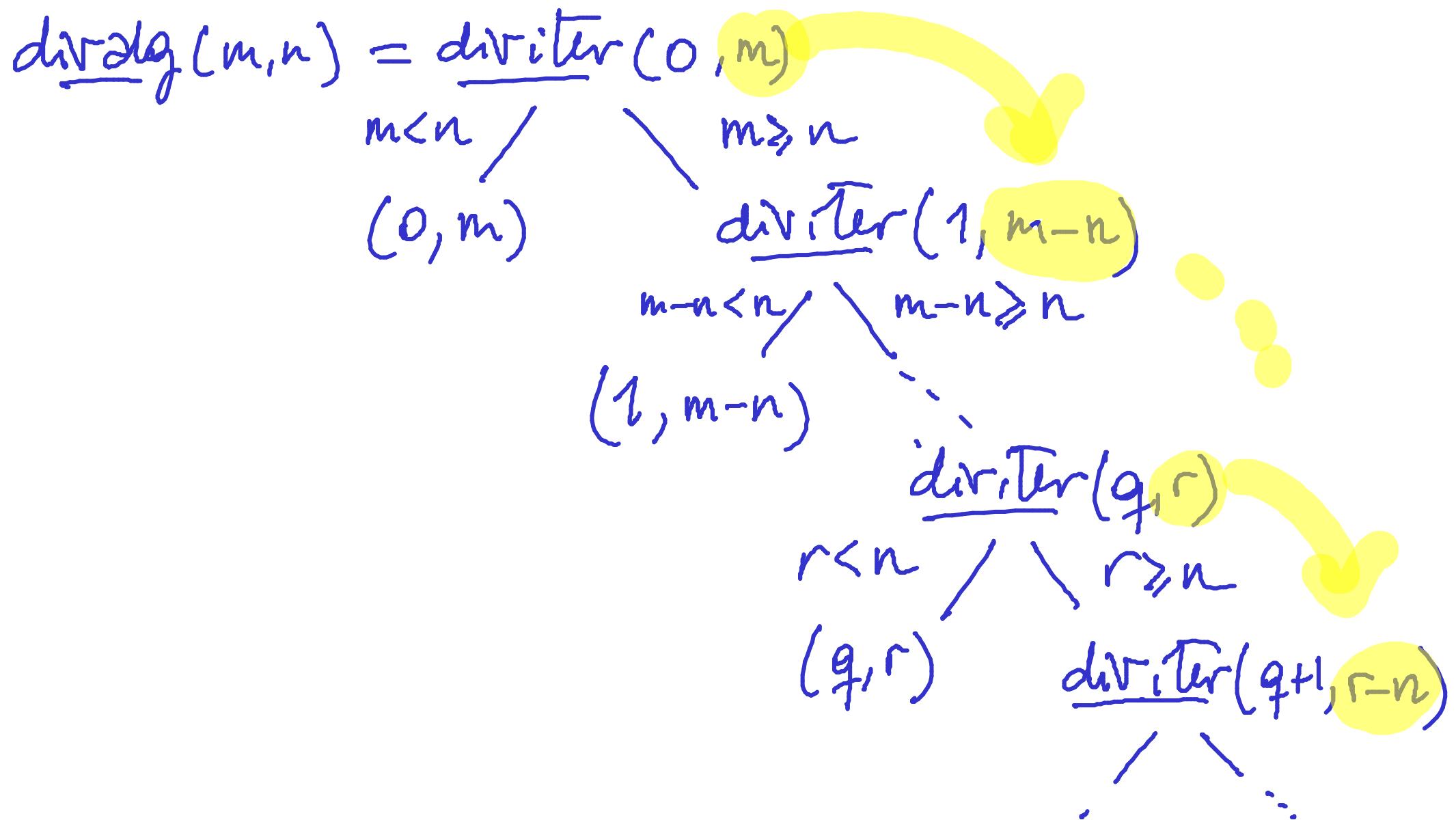
$$(q, r)$$

$$\underline{\text{diviter}}(q+1, r-n)$$

$$\dots$$

**Theorem 45** *For every natural number  $m$  and positive natural number  $n$ , the evaluation of  $\text{divalg}(m, n)$  terminates, outputting a pair of natural numbers  $(q_0, r_0)$  such that  $r_0 < n$  and  $m = q_0 \cdot n + r_0$ .*

PROOF:



As for partial correctness; i.e. That

$$m = \underline{\text{quo}}(m, n) \cdot n + \underline{\text{rem}}(m, n)$$

$$0 \leq \underline{\text{rem}}(m, n) < n$$

(\*)

We show the invariant property That  
on all calls of diviter( $q, r$ ) one has

$$m = q \cdot n + r$$

The last call will therefore yield (\*).

$$m = 0 \cdot n + m, m \geq 0$$

$$\underline{\text{divAlg}}(m, n) = \underline{\text{divIter}}(0, m)$$

$$\begin{array}{c} m < n \\ \swarrow \quad \searrow \\ (0, m) \end{array}$$

Suppose

$$m = qn + r, r \geq 0$$

If  $r < n$  then

$$\underline{\text{quo}}(m, n) = q$$

$$\underline{\text{rem}}(m, n) = r$$

satisfy the required properties

Otherwise  $r \geq n$  and we note that

The invariant is maintained as

$$m = (q+1)n + (r-n), r-n \geq 0.$$

$$\underline{\text{divIter}}(1, m-n)$$

$$\begin{array}{c} m-n < n \\ \swarrow \quad \searrow \\ (1, m-n) \end{array}$$

$$\underline{\text{divIter}}(q, r)$$

$$\begin{array}{c} r < n \\ \swarrow \quad \searrow \\ (q, r) \end{array}$$

$$\underline{\text{divIter}}(q+1, r-n)$$

$$\begin{array}{c} \vdots \\ \vdots \end{array}$$

**Proposition 46** Let  $m$  be a positive integer. For all natural numbers  $k$  and  $l$ ,

$$k \equiv l \pmod{m} \iff \text{rem}(k, m) = \text{rem}(l, m)$$

PROOF: Let  $m$  be a positive integer.

Let  $k$  and  $l$  be natural numbers.

$\Rightarrow$  Assume  $k \equiv l \pmod{m}$

Then,  $k = l + im$  for some integer  $i$

$$= [\underline{\text{quo}}(l, m) + i] \cdot m + \underline{\text{rem}}(l, m)$$

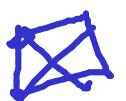
Therefore  $\underline{\text{rem}}(k, m) = \underline{\text{rem}}(l, m)$

Let  $m$  be a positive integer.

Let  $k$  and  $l$  be natural numbers.

( $\Leftarrow$ ) Assume:  $\underline{\text{rem}}(k, m) = \underline{\text{rem}}(l, m)$

$$\begin{aligned} \text{Then, } k - l &= [\underline{\text{quo}}(R, m) - \underline{\text{quo}}(l, m)].m \\ &\quad + [\underline{\text{rem}}(R, m) - \underline{\text{rem}}(l, m)] \\ &= [\underline{\text{quo}}(R, m) - \underline{\text{quo}}(l, m)].m \end{aligned}$$



**Corollary 47** Let  $m$  be a positive integer.

1. For every natural number  $n$ ,

$$n \equiv \text{rem}(n, m) \pmod{m} .$$

PROOF:

**Corollary 47** Let  $m$  be a positive integer.

1. For every natural number  $n$ ,

$$n \equiv \text{rem}(n, m) \pmod{m}$$

2. For every integer  $k$  there exists a unique integer  $[k]_m$  such that

$$0 \leq [k]_m < m \text{ and } k \equiv [k]_m \pmod{m}$$

PROOF:

$$[k]_m = \begin{cases} \underline{\text{if }} \underline{\text{rem}}(k, m) = 0 \text{ Then } 0 \\ \underline{\text{else if }} k > 0 \text{ Then } \underline{\text{rem}}(k, m) \\ \underline{\text{else }} m - \underline{\text{rem}}(-k, m) \end{cases}$$