

UNIVERSITY OF
CAMBRIDGE
COMPUTER LABORATORY



Advanced Graphics & Image Processing

Parallel programming in OpenCL

Part 1/3 – OpenCL framework

Rafał Mantiuk

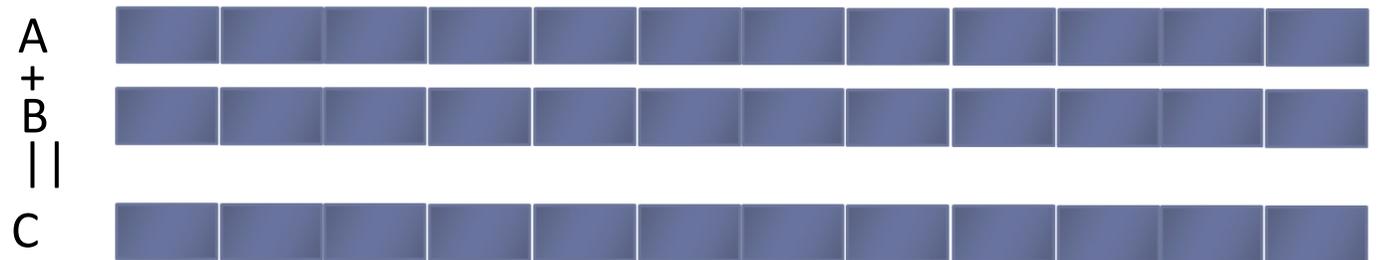
Computer Laboratory, University of Cambridge

Single Program Multiple Data (SPMD)

- ▶ Consider the following vector addition example

Serial program:
one program completes
the entire task

```
for( i = 0:11 ) {  
    C[ i ] = A[ i ] + B[ i ]  
}
```



Multiple copies of the same program execute on different data in parallel

SPMD program:
multiple copies of the
same program run on
different chunks of the
data

```
for( i = 0:3 ) {  
    C[ i ] = A[ i ] + B[ i ]  
}  
for( i = 4:7 ) {  
    C[ i ] = A[ i ] + B[ i ]  
}  
for( i = 8:11 ) {  
    C[ i ] = A[ i ] + B[ i ]  
}
```



Parallel Software – SPMD

- ▶ In the vector addition example, each chunk of data could be executed as an independent thread
- ▶ On modern CPUs, the overhead of creating threads is so high that the chunks need to be large
 - ▶ In practice, usually a few threads (about as many as the number of CPU cores) and each is given a large amount of work to do
- ▶ For GPU programming, there is low overhead for thread creation, so we can create one thread per loop iteration

Parallel Software – SPMD

 = loop iteration

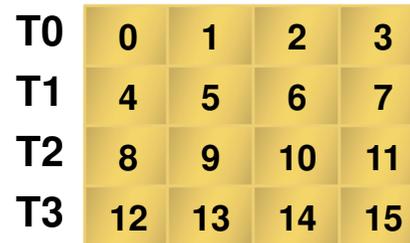
Single-threaded (CPU)

```
// there are N elements
for(i = 0; i < N; i++)
    C[i] = A[i] + B[i]
```



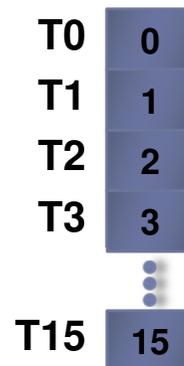
Multi-threaded (CPU)

```
// tid is the thread id
// P is the number of cores
for(i = 0; i < tid*N/P; i++)
    C[i] = A[i] + B[i]
```



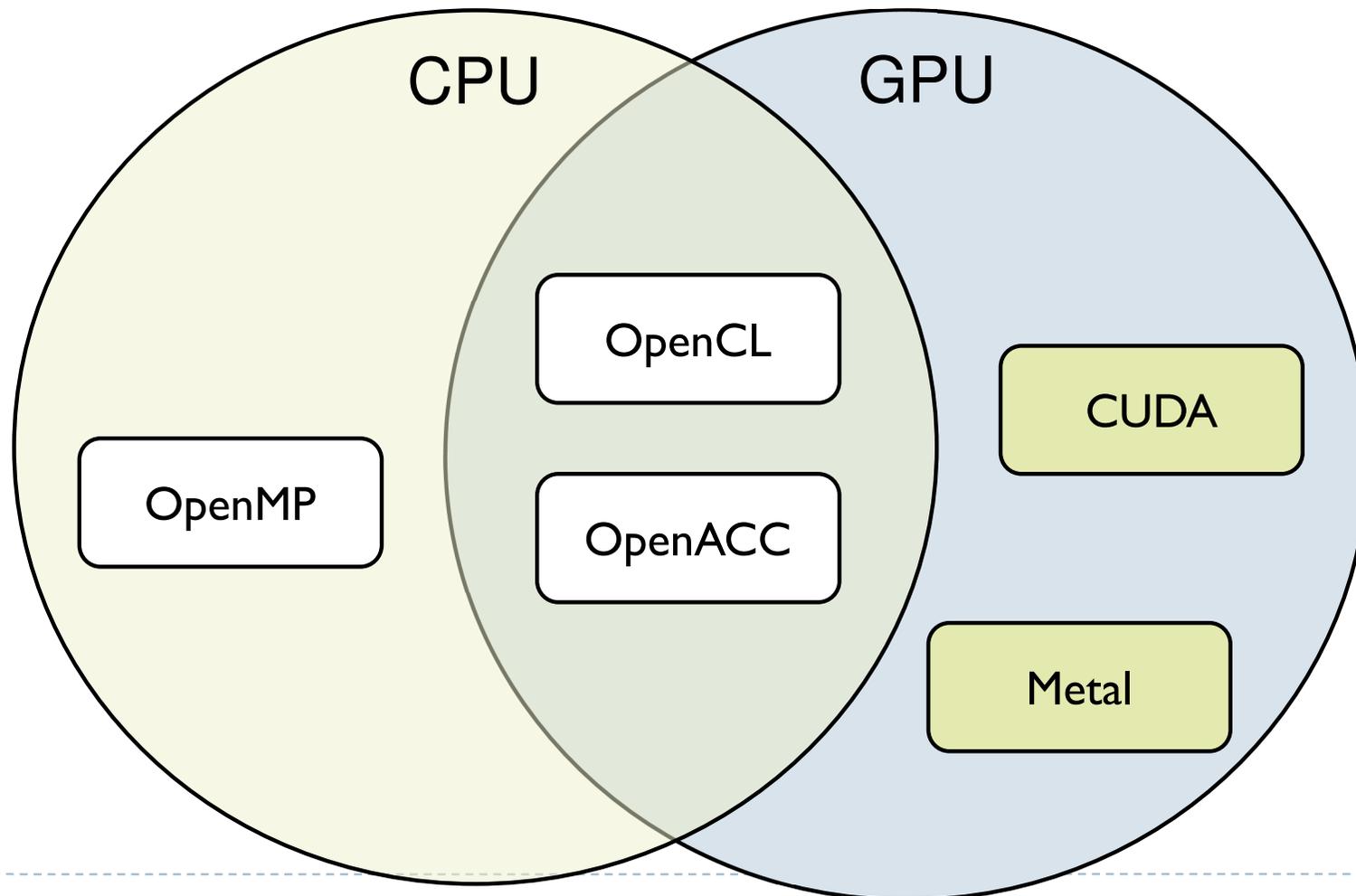
Massively Multi-threaded (GPU)

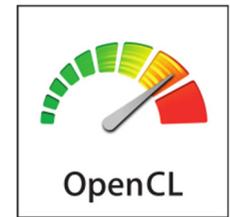
```
// tid is the thread id
C[tid] = A[tid] + B[tid]
```



Parallel programming frameworks

- ▶ These are some of more relevant frameworks for creating parallelized code





OpenCL

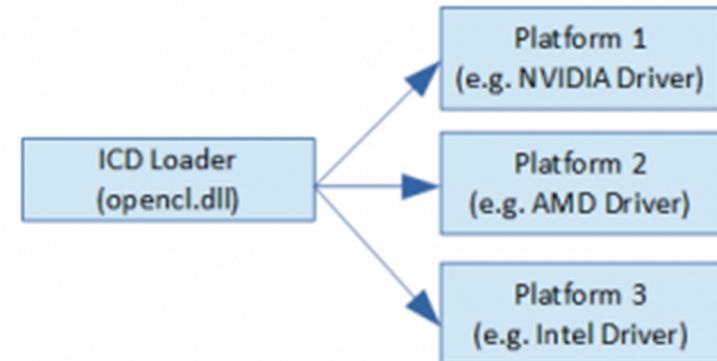
- ▶ OpenCL is a framework for writing parallelized code for CPUs, GPUs, DSPs, FPGAs and other processors
- ▶ Initially developed by Apple, now supported by AMD, IBM, Qualcomm, Intel and Nvidia (reluctantly)
- ▶ Versions
 - ▶ Latest: OpenCL 2.2
 - ▶ OpenCL C++ kernel language
 - ▶ SPIR-V as intermediate representation for kernels
 - Vulkan uses the same Standard Portable Intermediate Representation
 - ▶ AMD, Intel
 - ▶ Mostly supported: OpenCL 1.2
 - ▶ Nvidia, OSX



OpenCL platforms and drivers

- ▶ To run OpenCL code you need:

- ▶ Generic ICD loader
 - ▶ Included in the OS
- ▶ Installable Client Driver
 - ▶ From Nvidia, Intel, etc.



- ▶ This applies to Windows and Linux, only one platform on Mac

- ▶ To develop OpenCL code you need:

- ▶ OpenCL headers/libraries
 - ▶ Included in the SDKs
 - Nvidia – CUDA Toolkit
 - Intel OpenCL SDK
 - ▶ But lightweight options are also available

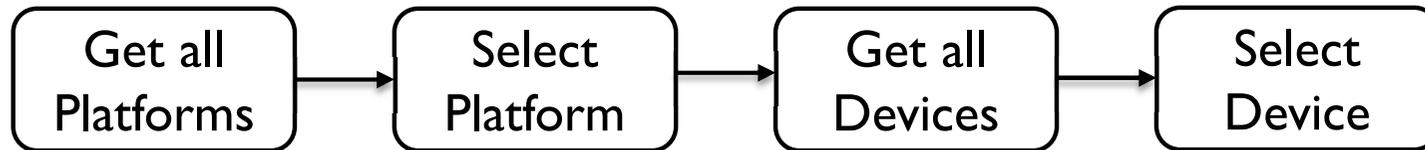


Programming OpenCL

- ▶ OpenCL natively offers C99 API
- ▶ But there is also a standard OpenCL C++ API wrapper
 - ▶ Strongly recommended – reduces the amount of code
- ▶ Programming OpenCL is similar to programming shaders in OpenGL
 - ▶ Host code runs on CPU and invokes **kernels**
 - ▶ Kernels are written in C-like programming language
 - ▶ In many respects similar to GLSL
 - ▶ Kernels are passed to API as strings and compiled at runtime
 - ▶ Kernels are usually stored in text files
 - ▶ Kernels can be precompiled into SPIR from OpenCL 2.1



Example: Step 1 - Select device

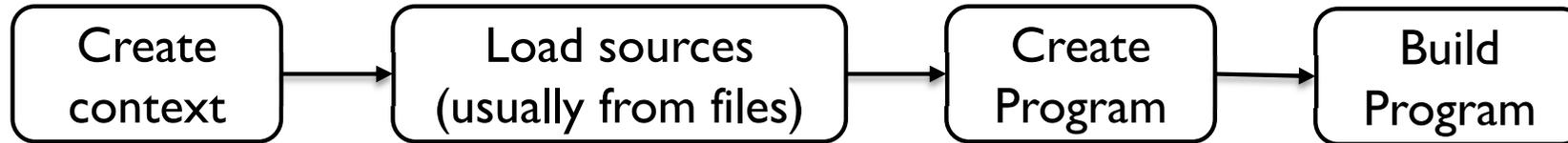


```
//get all platforms (drivers)
std::vector<cl::Platform> all_platforms;
cl::Platform::get(&all_platforms);
if (all_platforms.size() == 0){
    std::cout << " No platforms found. Check OpenCL installation!\n";
    exit(1);
}
cl::Platform default_platform = all_platforms[0];
std::cout << "Using platform: " << default_platform.getInfo<CL_PLATFORM_NAME>() << "\n";

//get default device of the default platform
std::vector<cl::Device> all_devices;
default_platform.getDevices(CL_DEVICE_TYPE_ALL, &all_devices);
if (all_devices.size() == 0){
    std::cout << " No devices found. Check OpenCL installation!\n";
    exit(1);
}
cl::Device default_device = all_devices[0];
std::cout << "Using device: " << default_device.getInfo<CL_DEVICE_NAME>() << "\n";
```



Example: Step 2 - Build program



```
cl::Context context({ default_device });
```

```
cl::Program::Sources sources;
```

```
// kernel calculates for each element C=A+B
```

```
std::string kernel_code =
```

```
    "__kernel void simple_add(__global const int* A, __global const int* B, __global int* C) {"  
    "    int index = get_global_id(0);"  
    "    C[index] = A[index] + B[index];"  
    "};";
```

```
sources.push_back({ kernel_code.c_str(), kernel_code.length() });
```

```
cl::Program program(context, sources);
```

```
try {
```

```
    program.build({ default_device });
```

```
}
```

```
catch (cl::Error err) {
```

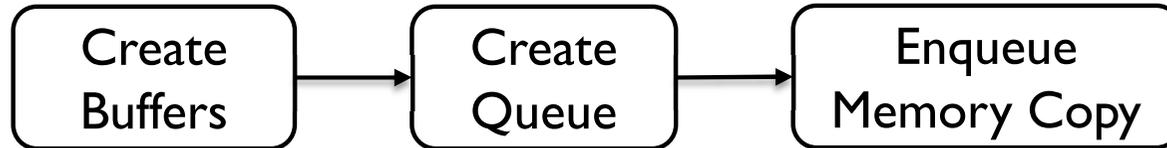
```
    std::cout << " Error building: " <<
```

```
        program.getBuildInfo<CL_PROGRAM_BUILD_LOG>(default_device) << "\n";
```

```
    exit(1);
```

```
}
```

Example: Step 3 - Create Buffers and copy memory



```
// create buffers on the device
cl::Buffer buffer_A(context, CL_MEM_READ_WRITE, sizeof(int) * 10);
cl::Buffer buffer_B(context, CL_MEM_READ_WRITE, sizeof(int) * 10);
cl::Buffer buffer_C(context, CL_MEM_READ_WRITE, sizeof(int) * 10);

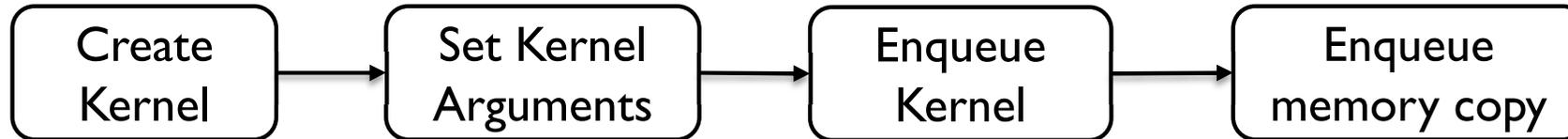
int A[] = { 0, 1, 2, 3, 4, 5, 6, 7, 8, 9 };
int B[] = { 0, 1, 2, 0, 1, 2, 0, 1, 2, 0 };

//create queue to which we will push commands for the device.
cl::CommandQueue queue(context, default_device);

//write arrays A and B to the device
queue.enqueueWriteBuffer(buffer_A, CL_TRUE, 0, sizeof(int) * 10, A);
queue.enqueueWriteBuffer(buffer_B, CL_TRUE, 0, sizeof(int) * 10, B);
```



Example: Step 4 - Execute Kernel and retrieve the results



```
cl::Kernel kernel(program, "simple_add");

kernel.setArg(0, buffer_A);
kernel.setArg(1, buffer_B);
kernel.setArg(2, buffer_C);
queue.enqueueNDRangeKernel(kernel, cl::NullRange, cl::NDRange(10), cl::NullRange);

int C[10];
//read result C from the device to array C
queue.enqueueReadBuffer(buffer_C, CL_TRUE, 0, sizeof(int) * 10, C);
queue.finish();

std::cout << " result: \n";
for (int i = 0; i < 10; i++){
    std::cout << C[i] << " ";
}
std::cout << std::endl;
```

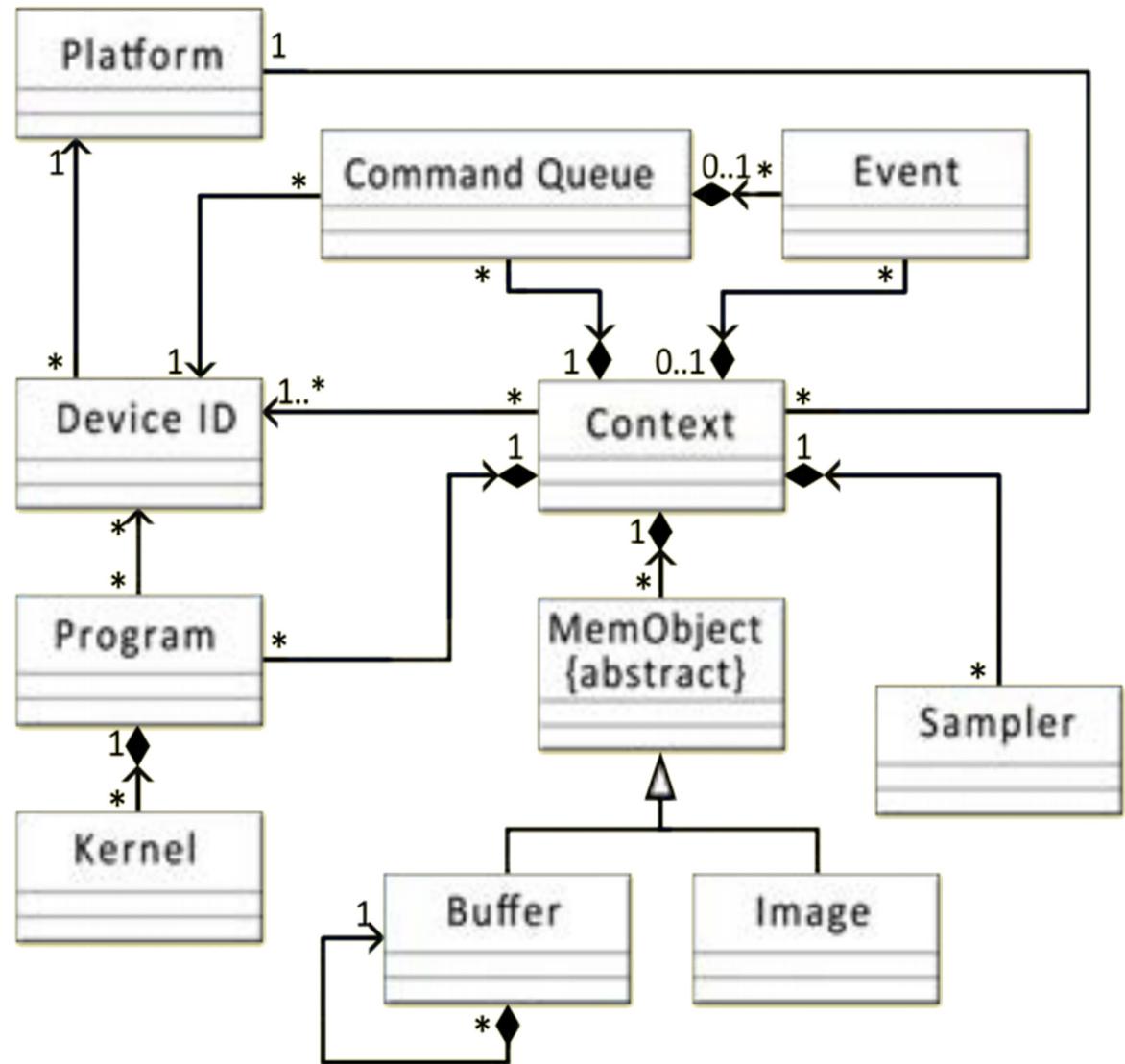
Our Kernel was

```
__kernel void simple_add(__read_only const int* A,
                        __read_only const int* B,
                        __write_only int* C) {

    int index = get_global_id(0);
    C[index]=A[index]+B[index];
};
```

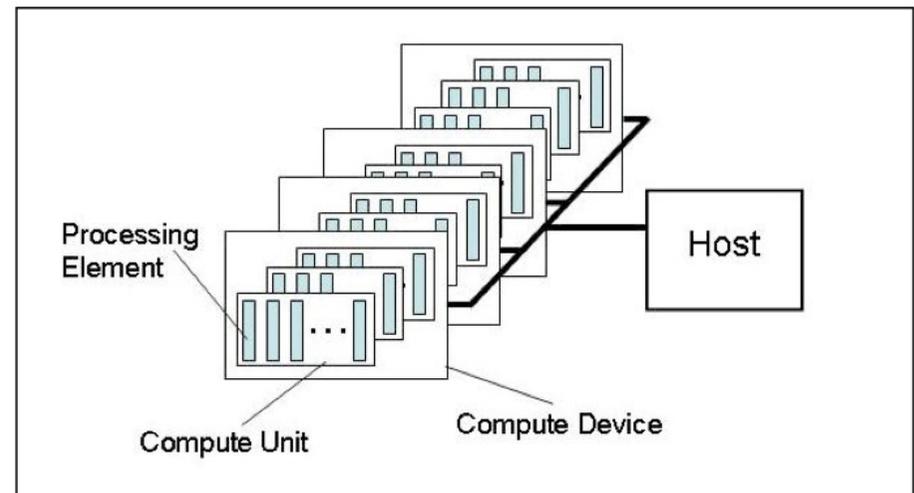
OpenCL API Class Diagram

- ▶ *Platform* – Nvidia CUDA
- ▶ *Device* – GeForce 780
- ▶ *Program* – collection of kernels
- ▶ *Buffer / Image* – device memory
- ▶ *Sampler* – how to interpolate values for *Image*
- ▶ *Command Queue* – put a sequence of operations there
- ▶ *Event* – to notify that something has been done



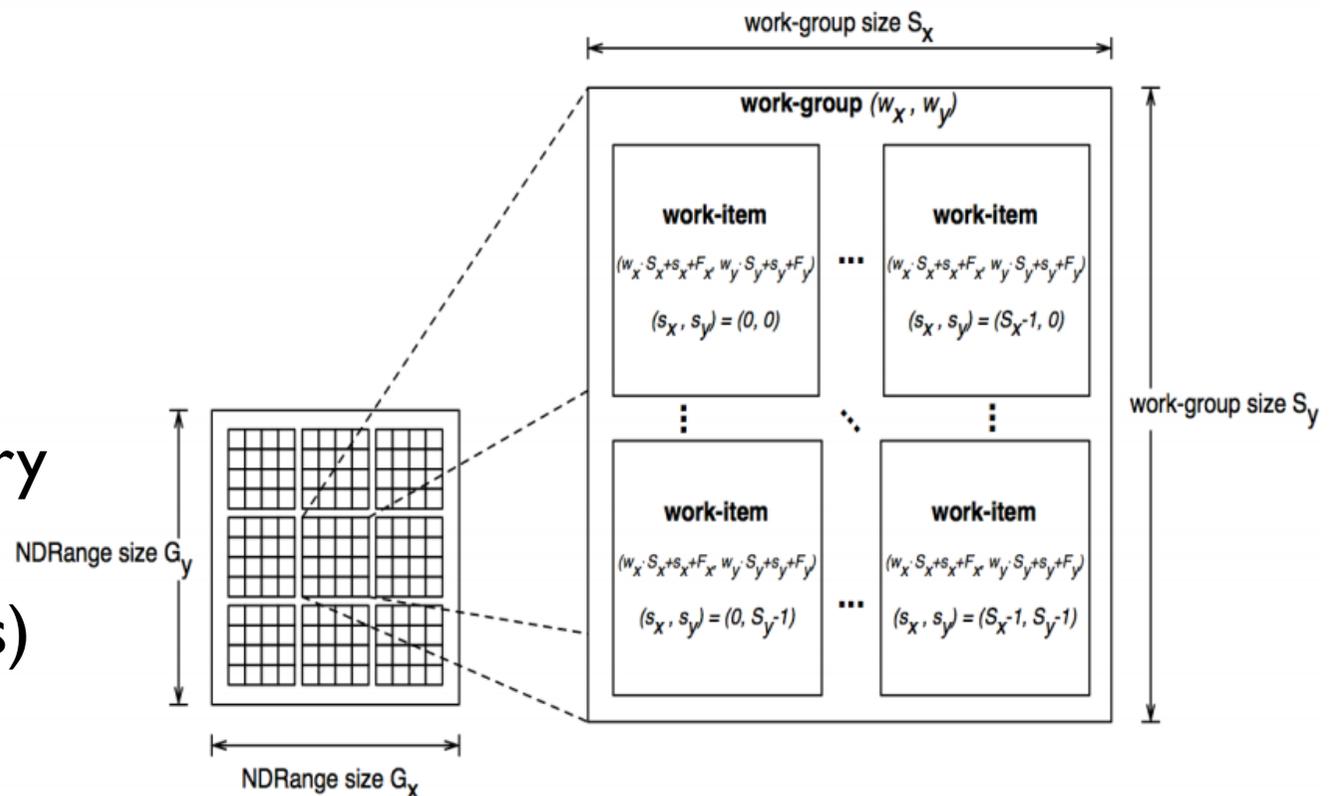
Platform model

- ▶ The host is whatever the OpenCL library runs on
 - ▶ Usually x86 CPUs for both NVIDIA and AMD
- ▶ Devices are processors that the library can talk to
 - ▶ CPUs, GPUs, DSP,s and generic accelerators
- ▶ For AMD
 - ▶ All CPUs are combined into a single device (each core is a compute unit and processing element)
 - ▶ Each GPU is a separate device



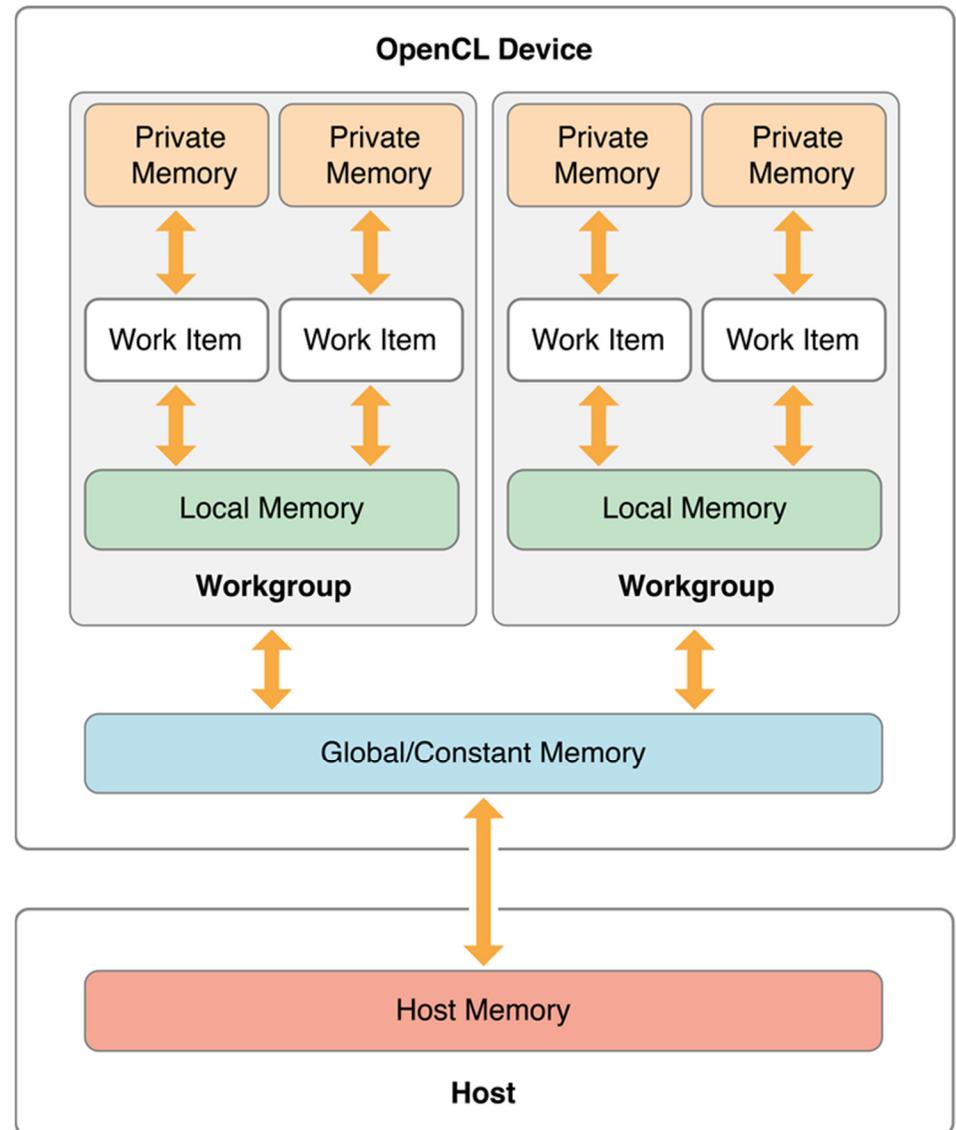
Execution model

- ▶ Each kernel executes on 1D, 2D or 3D array (NDRange)
- ▶ The array is split into work-groups
- ▶ Work items (threads) in each work-group share some local memory
- ▶ Kernel can query
 - ▶ `get_global_id(dim)`
 - ▶ `get_group_id(dim)`
 - ▶ `get_local_id(dim)`
- ▶ Work items are not bound to any memory entity (unlike GLSL shaders)

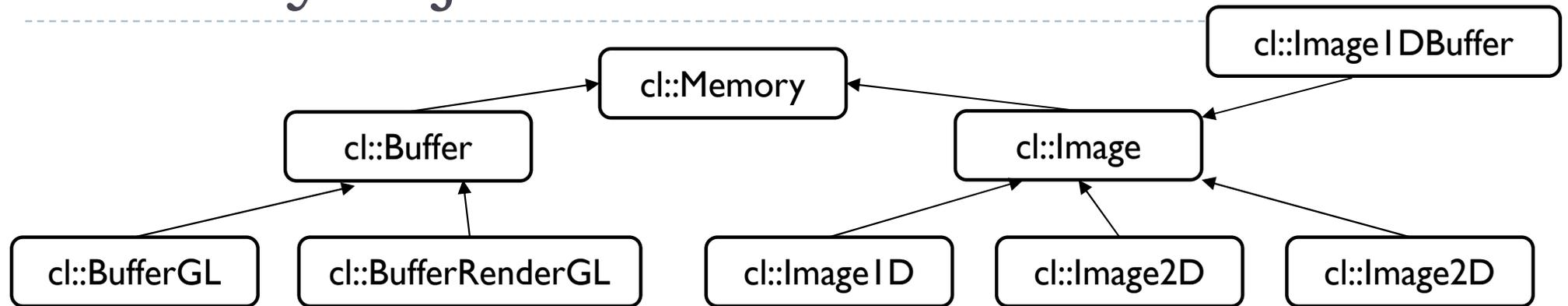


Memory model

- ▶ **Host memory**
 - ▶ Usually CPU memory, device does not have access to that memory
- ▶ **Global memory [`__global`]**
 - ▶ Device memory, for storing large data
- ▶ **Constant memory [`__constant`]**
- ▶ **Local memory [`__local`]**
 - ▶ Fast, accessible to all work-items (threads) within a workgroup
- ▶ **Private memory [`__private`]**
 - ▶ Accessible to a single work-item (thread)



Memory objects



This diagram is incomplete – there are more memory objects

▶ Buffer

- ▶ ArrayBuffer in OpenGL
- ▶ Accessed directly via C pointers

▶ Image

- ▶ Texture in OpenGL
- ▶ Access via texture look-up function
- ▶ Can interpolate values, clamp, etc.



Programming model

- ▶ Data parallel programming
 - ▶ Each NDRange element is assigned to a work-item (thread)
 - ▶ Each kernel can use vector-types of the device (float4, etc.)
- ▶ Task-parallel programming
 - ▶ Multiple different kernels can be executed in parallel
- ▶ Command queue

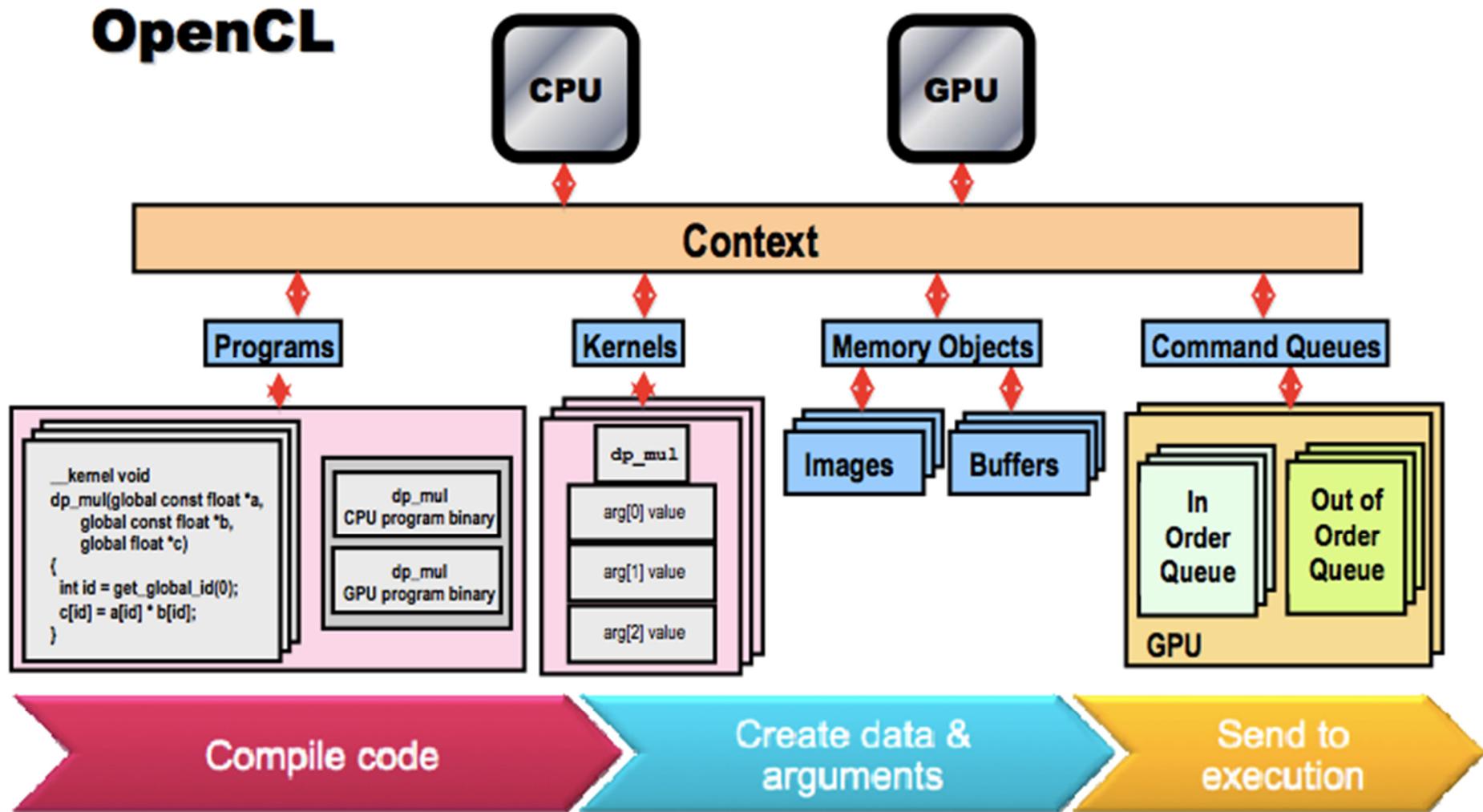
```
clCreateCommandQueue(  
    cl_context context,  
    cl_device_id device,  
    cl_command_queue_properties properties,  
    cl_int* errcode_ret)
```

CL_QUEUE_OUT_OF_ORDER_EXEC_MODE_ENABLE
Execute out-of-order if specified, in order otherwise

- ▶ Provides means to both synchronize kernels and execute them in parallel
-



Big Picture



UNIVERSITY OF
CAMBRIDGE
COMPUTER LABORATORY



Advanced Graphics & Image Processing

Parallel programming in OpenCL

Part 2/3 – Thread mapping

Rafał Mantiuk

Computer Laboratory, University of Cambridge

Thread Mapping

- ▶ By using different mappings, the same thread can be assigned to access different data elements
 - ▶ The examples below show three different possible mappings of threads to data (assuming the thread id is used to access an element)

Mapping

```
int tid =
get_global_id(1) *
get_global_size(0) +
get_global_id(0);
```

```
int tid =
get_global_id(0) *
get_global_size(1) +
get_global_id(1);
```

```
int group_size =
get_local_size(0) *
get_local_size(1);
```

```
int tid =
get_group_id(1) *
get_num_groups(0) *
group_size +
get_group_id(0) *
group_size +
get_local_id(1) *
get_local_size(0) +
get_local_id(0)
```

Thread IDs

0	1	2	3
4	5	6	7
8	9	10	11
12	13	14	15

0	4	8	12
1	5	9	13
2	6	10	14
3	7	11	15

0	1	4	5
2	3	6	7
8	9	12	13
10	11	14	15

*assuming 2x2 groups

Thread Mapping

- ▶ Consider a serial matrix multiplication algorithm

```
for(i1=0; i1 < M; i1++)  
  for(i2=0; i2 < N; i2++)  
    for(i3=0; i3 < P; i3++)  
      C[i1][i2] += A[i1][i3]*B[i3][i2];
```

- ▶ This algorithm is suited for output data decomposition
 - ▶ We will create $N \times M$ threads
 - ▶ Effectively removing the outer two loops
 - ▶ Each thread will perform P calculations
 - ▶ The inner loop will remain as part of the kernel
- ▶ Should the index space be $M \times N$ or $N \times M$?

Thread Mapping

- ▶ Thread mapping 1: with an $M \times N$ index space, the kernel would be:

```
int tx = get_global_id(0);
int ty = get_global_id(1);
for(i3=0; i3<P; i3++)
    C[tx][ty] += A[tx][i3]*B[i3][ty];
```

Mapping for C

0	4	8	12
1	5	9	13
2	6	10	14
3	7	11	15

- ▶ Thread mapping 2: with an $N \times M$ index space, the kernel would be:

```
int tx = get_global_id (0);
int ty = get_global_id (1);
for(i3=0; i3<P; i3++)
    C[ty][tx] += A[ty][i3]*B[i3][tx];
```

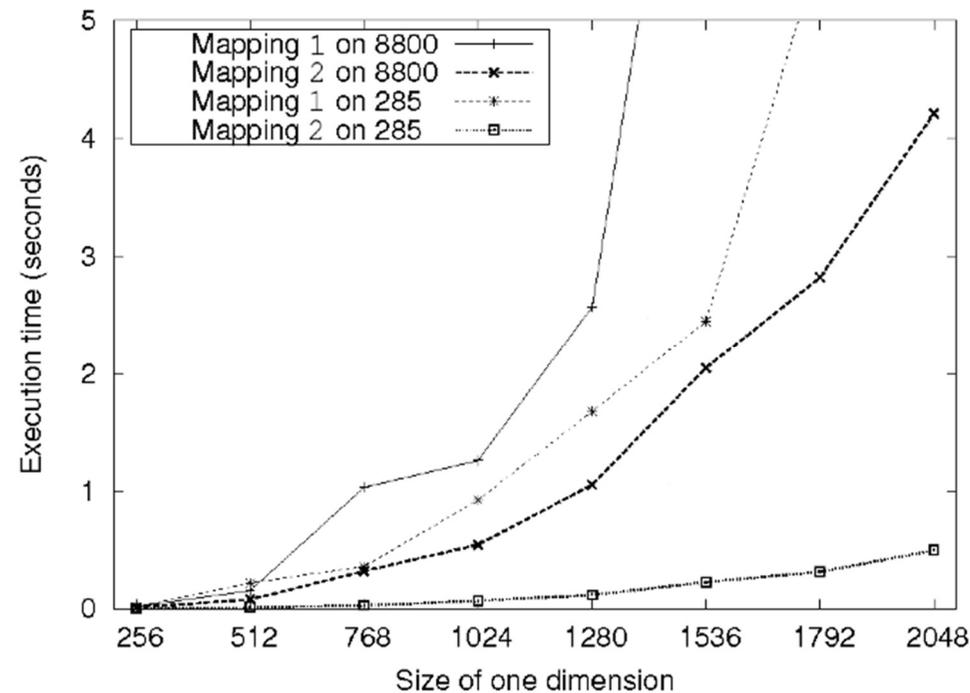
Mapping for C

0	1	2	3
4	5	6	7
8	9	10	11
12	13	14	15

- ▶ Both mappings produce functionally equivalent versions of the program

Thread Mapping

- ▶ This figure shows the execution of the two thread mappings on NVIDIA GeForce 285 and 8800 GPUs



- ▶ Notice that mapping 2 is far superior in performance for both GPUs

Thread Mapping

- ▶ The discrepancy in execution times between the mappings is due to data accesses on the global memory bus
 - ▶ Assuming row-major data, data in a row (i.e., elements in adjacent columns) are stored sequentially in memory
 - ▶ To ensure coalesced accesses, consecutive threads in the same wavefront should be mapped to columns (the second dimension) of the matrices
 - ▶ This will give coalesced accesses in Matrices B and C
 - ▶ For Matrix A, the iterator $i3$ determines the access pattern for row-major data, so thread mapping does not affect it

UNIVERSITY OF
CAMBRIDGE
COMPUTER LABORATORY



Advanced Graphics & Image Processing

Parallel programming in OpenCL

Part 3/3 – Reduction

Rafał Mantiuk

Computer Laboratory, University of Cambridge

Reduction

- ▶ GPU offers very good performance for tasks in which the results are stored independently
 - ▶ Process N data items and store in N memory location

```
float reduce_sum(float* input, int length)
{
    float accumulator = input[0];
    for(int i = 1; i < length; i++)
        accumulator += input[i];
    return accumulator;
}
```

- ▶ But many common operations require reducing N values into 1 or few values
 - ▶ sum, min, max, prod, min, histogram, ...
- ▶ Those operations require an efficient implementation of reduction
- ▶ The following slides are based on AMD's OpenCL™ Optimization Case Study: Simple Reductions
 - ▶ <http://developer.amd.com/resources/articles-whitepapers/opencl-optimization-case-study-simple-reductions/>



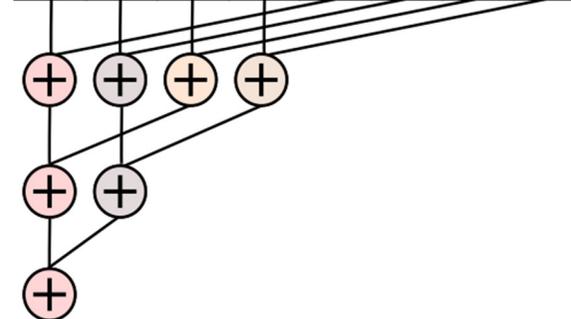
Reduction tree for the min operation

```
__kernel
void reduce_min(__global float* buffer,
               __local float* scratch,
               __const int length,
               __global float* result) {

    int global_index = get_global_id(0);
    int local_index = get_local_id(0);
    // Load data into local memory
    if (global_index < length) {
        scratch[local_index] = buffer[global_index];
    } else {
        scratch[local_index] = INFINITY;
    }
    barrier(CLK_LOCAL_MEM_FENCE);
    for(int offset = get_local_size(0) / 2;
        offset > 0; offset >>= 1) {
        if (local_index < offset) {
            float other = scratch[local_index + offset];
            float mine = scratch[local_index];
            scratch[local_index] = (mine < other) ? mine :
other;
        }
        barrier(CLK_LOCAL_MEM_FENCE);
    }
    if (local_index == 0) {
        result[get_group_id(0)] = scratch[0];
    }
}
```

- ▶ barrier ensures that all threads (work units) in the local group reach that point before execution continue
- ▶ Each iteration of the for loop computes next level of the reduction pyramid

Local memory



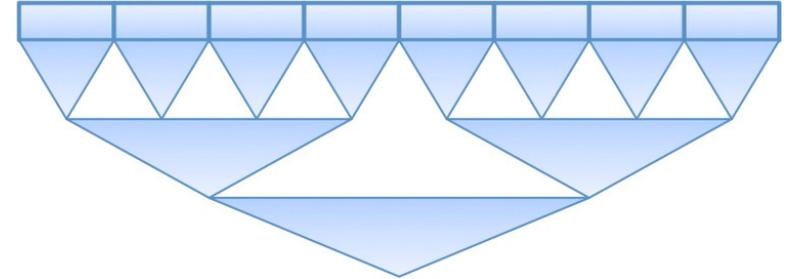
Parallel Reduction Tree for Commutative Operator



SIMD Utilization for Reduction Tree

Multistage reduction

- ▶ The local memory is usually limited (e.g. 50kB), which restricts the maximum size of the array that can be processed
- ▶ Therefore, for large arrays need to be processed in multiple stages
 - ▶ The result of a local memory reduction is stored in the array and then this array is reduced



Two-stage reduction

Stage 1

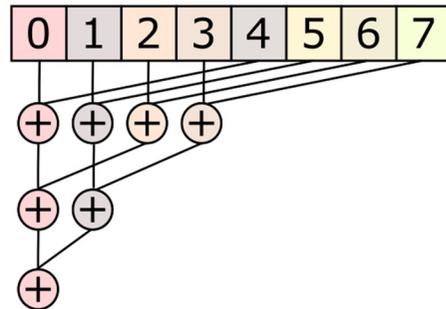
Different colours denote different threads

Global memory



Local memory

Stage 2



```
__kernel
void reduce(__global float* buffer,
            __local float* scratch,
            __const int length,
            __global float* result) {

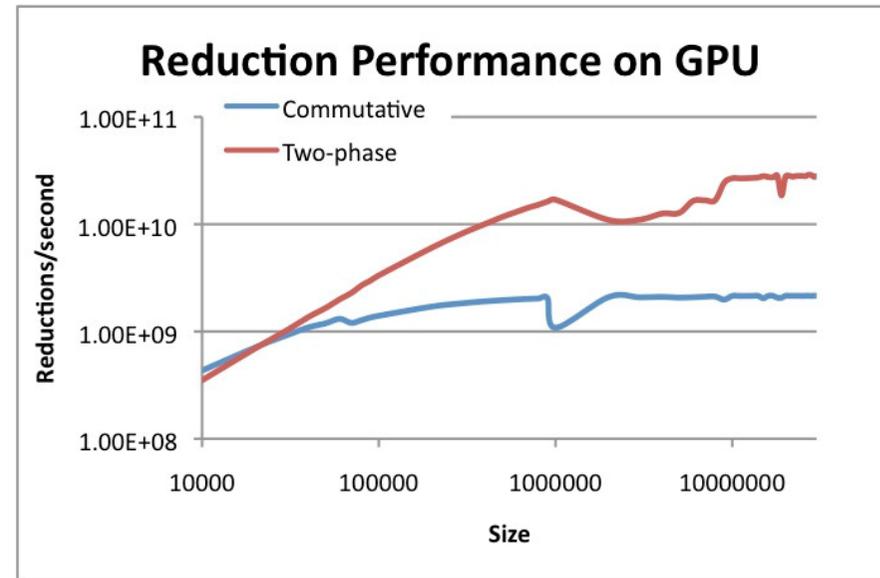
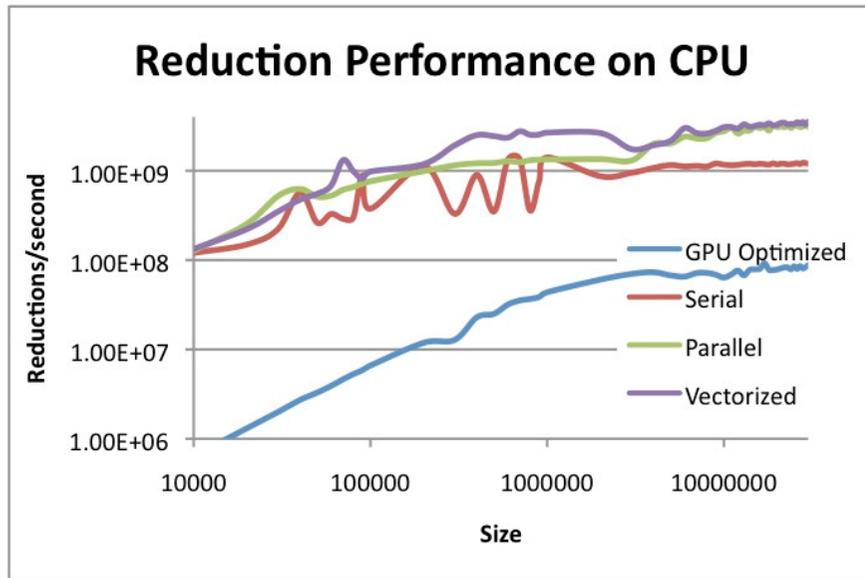
    int global_index = get_global_id(0);
    float accumulator = INFINITY;
    // Loop sequentially over chunks of input
    vector
    while (global_index < length) {
        float element = buffer[global_index];
        accumulator = (accumulator < element) ?
accumulator : element;
        global_index += get_global_size(0);
    }

    // Perform parallel reduction
    [The same code as in the previous example]
}
```

- ▶ First stage: serial reduction by N concurrent threads
 - ▶ Number of threads $<$ data items
- ▶ Second stage: parallel reduction in local memory



Reduction performance CPU/GPU

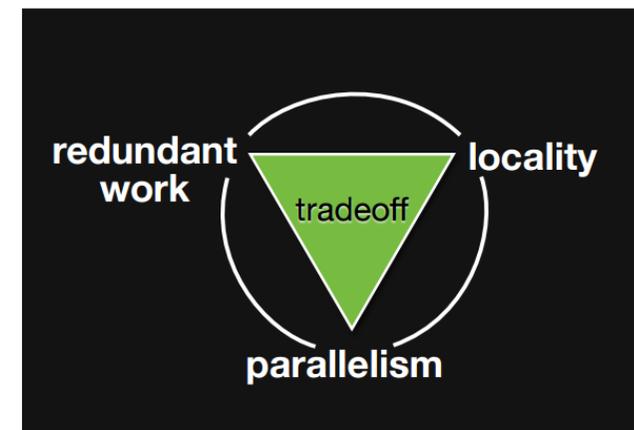


- ▶ Different reduction algorithm may be optimal for CPU and GPU
- ▶ This can also vary from one GPU to another
- ▶ The results from: <http://developer.amd.com/resources/articles-whitepapers/opencl-optimization-case-study-simple-reductions/>



Better way?

- ▶ **Halide** - a language for image processing and computational photography
 - ▶ <http://halide-lang.org/>
 - ▶ Code written in a high-level language, then translated to x86/SSE, ARM, CUDA, OpenCL
 - ▶ The optimization strategy defined separately as a *schedule*
 - ▶ Auto-tune software can test thousands of schedules and choose the one that is the best for a particular platform
 - ▶ (Semi-)automatically find the best trade-offs for a particular platform
 - ▶ Designed for image processing but similar languages created for other purposes



OpenCL resources

- ▶ <https://www.khronos.org/registry/OpenCL/>
- ▶ Reference cards
 - ▶ Google: “OpenCL API Reference Card”
- ▶ **AMD OpenCL Programming Guide**
 - ▶ http://developer.amd.com/wordpress/media/2013/07/AMD_Accelerated_Parallel_Processing_OC_L_Programming_Guide-2013-06-21.pdf

