Computer Design

Example Tagged Token and Instruction Formats

example token format:

tag data frame pointer statement pointer port type value where frame pointer = address of the start of the activation frame statement pointer = address of the target statement port = indicates left or right operand integer, floating point etc. = type value = typically 64 bits of data example instruction format:

op-code (r) dest1 (dest2)

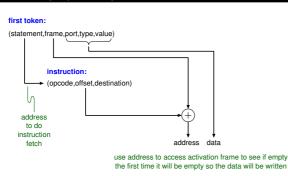
where () indicates optional parameters

op-code is the instruction identifier
r is the activation frame offset number for dyadic operations
dest1 and dest2 are the destination offsets (dest2 being optional)

Matching Algorithm for Dyadic Operations

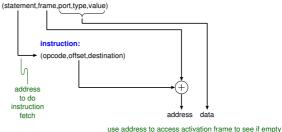
- incoming token's statement pointer is used to look up the instruction
- the instruction's activation frame offset is added to the token's activation frame number to give an effective address
- the *effective address* is then used to look up the the *presence* bit in the activation frame
- if the presence = empty then the token's value and port are written to the location
- if the presence = full then the stored value and token value should make up the two operands for the dyadic instruction (assuming their ports are different)
- the operation, its operands and the destination(s) are executed note:
- these stages correspond to the stages in the pipeline

Matching Dyadic Operations cont...



Matching Dyadic Operations cont...

second token:



this time it will be full so the data pair will be sent for execution

Example Tagged-token Data-flow Program

address	instruction		
(e.g.)	op-code	offset	destinations
0x30	mul	0,	0x31ℓ,nil
0x31	add	2,	0x33ℓ,nil
0x32	div	1,	0x31r,nil
0x33	ret	0,	(dest)ℓ,nil

note:

ret accepts a $\langle \mathit{destination instruction, port, frame} \rangle$ triplet as its left parameter

advantages:

- simple matching algorithm which may be implemented using a pipeline
- garbage collecting unmatched tokens is easy
- problems:
- pipeline bubble every time the first operand of an instruction is matched
- token explosion problem can still occur (careful code generation required)

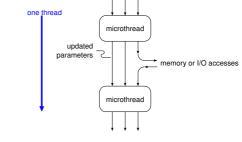
Evaluation of Data-flow

advantages:

- inherently concurrent and latency tolerant (no need for caches)
- multiprocessor applications are easy to write
- disadvantages:
- assignment a problem because there is too much concurrency, thus functional languages tend to be used. Furthermore, this makes I/O difficult
- ineffective use of very local storage (a register file or stack)
- scheduling policies have to be simple because of the instruction level concurrency

Multithreaded Processors — Combining Controlflow and Data-flow

example machine: Anaconda (Cambridge)



- unit of execution is larger so matching time does not dominate
- concurrency allows memory latency to be tolerated