Programming the ARM

Computer Design 2002, Lecture 4
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Quick Recap

- The Control Flow Model
 - □ Ordered list of instructions, fetch/execute, PC
- Instruction Set Architectures
 - □ Types of internal storage
 - Accumulator, stack and general purpose register machines
 - □ Addressing modes
 - Register indirect, displacement, memory indirect etc.
 - □ Branch conditions
 - Condition codes, condition registers, branch and compare
 - □ Encoding Instructions
 - Fixed length or variable length encodings
 - Representing immediates, opcode and operands

Quick Recap

- Making the common case fast and Amdahl's Law
 - See H&P (chapters 1 and 2)
- The Big Picture Interaction between the compiler, architecture and instruction set
 - E.g. a poorly designed instruction set may make highperformance implementation difficult and restrict the effectiveness of an optimizing compiler.

Lecture 4

- This Lecture
 - Implementing functions and procedures
 - □ The ARM Procedure Call Standard (APCS)
 - □ Development tools
 - □ Practical 4 of the ECAD labs

Functions and Procedures

- Structure program (abstraction/hierarchy)
- Package useful code so it can be reused
 - Must use a well defined interface
- Questions
 - □ How do we pass arguments to the function?
 - ☐ How do we obtain the result?

Procedure Calls

- Register values are preserved by saving and restoring them from a stack in memory
- Agree on a standard, define how registers are used (part of ABI – Application Binary Interface)
 - □ Which registers hold arguments
 - □ Which register holds the result (or pointer to)
 - Some registers may need to be saved (preserved) by caller (caller-saved)
 - Some may need to be saved by callee (callee-saved)

ARM Procedure Call Standard

- There are in fact many variants, lets look at base standard
- Four argument registers (r0-r3)
 - Not preserved by routine/function (callee)
 - May be saved by caller if necessary
 - ☐ May be used to return results to caller
- Registers r4-r11 are typically used to hold the routine's local variables
 - The value of these registers remains unchanged after a subroutine call
 - □ Registers may need to be saved and restored by callee
- Registers r12-r15 have special dedicated roles
 - □ e.g. the link register holds the return address

APCS Register Usage Convention

Register	APCS Name	APCS Role
0	a1	argument 1 / result / scratch register for function
1	a2	argument 2 / result / scratch register for function
2	а3	argument 3 / result / scratch register for function
3	a4	argument 4 / result / scratch register for function
4	v1	variable register 1
5	v2	variable register 2
6 v3 variat		variable register 3
7	v4	variable register 4
8	v5	variable register 5
9	sb or v6	static base or variable register 6
10	sl or v7	stack limit or variable register 7
11	fp	frame pointer (pointer to start of current stack frame) or variable register 8
12	12 ip the intra-procedure-call scratch register	
13	sp	lower end of current stack frame (stack pointer)
14	lr	the link register (return address)
15	рс	the program counter

Parameter Passing and result return

- a1-a4 (r0-r3) may be used to hold parameters
- Additional parameters may be passed on the stack
- Results may be returned in a1-a4 as a value or indirectly as an address

Simple function call

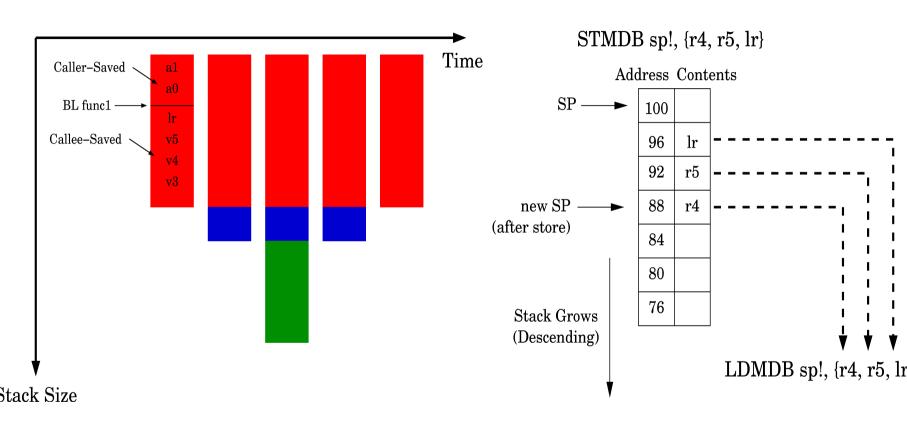
```
MOV r0, #10
            MOV r1, #5
             BL
                          max
             CMP
                          r0, r1
max
             MOVLT
                          r0, r1
                                       ; if r0<r1 r0=r1
             MOV
                          pc, Ir
```

For simple leaf functions that only use r0-r3 (a1-a4) the function overhead is small as no registers need to be saved. In real programs around half the functions may be simple leaf functions like this.

Saving and restoring registers

```
BL
                   myfunction
myfunction
            STMFD sp!, {r4-r10, lr}
                                      ; save registers
             LDMFD sp!, {r4-r10, pc}
                                      ; restore and
                                       ; return
```

The Stack



Full-descending stack - ARM stack grows down, stack pointer points to last entry (not next free entry)

ARM Assembler Examples

Assembler Directives

- Information for the assembler
- Common uses
 - □ Define and name new section, code or data
 - □ Constants, aliases
 - Allocate bytes of memory as data and initialize contents

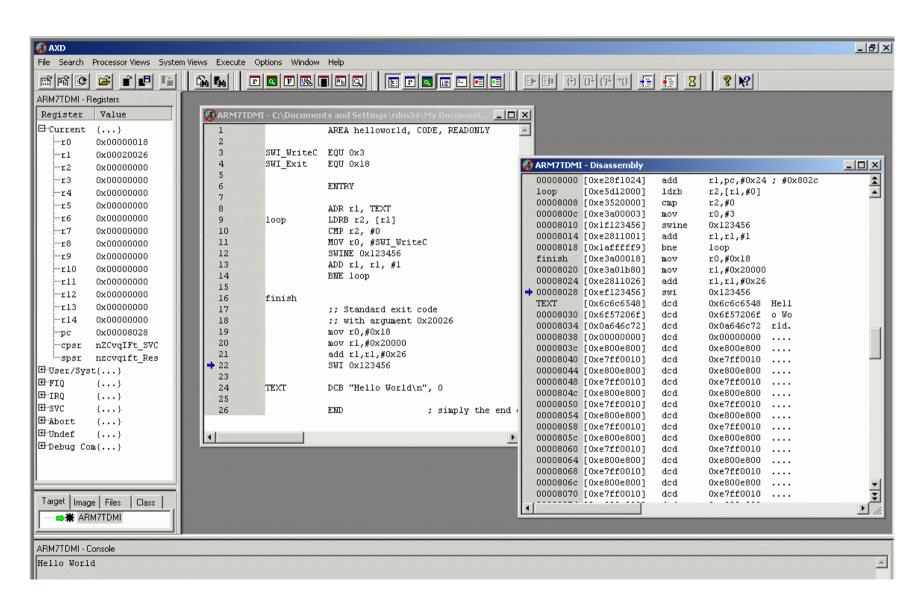
Allocating memory for data

- Allocate bytes of memory (DCB)
 - C_String DCB "MyString", 0
- Allocate words of memory (DCD)
 - Data DCD 1234, 1, 5, 20
- Reserve a 'zeroed' block of memory (SPACE or %)

```
table % 1024 ; 1024 bytes of zeroed ; memory
```

Writing Assembler, Hello World Example

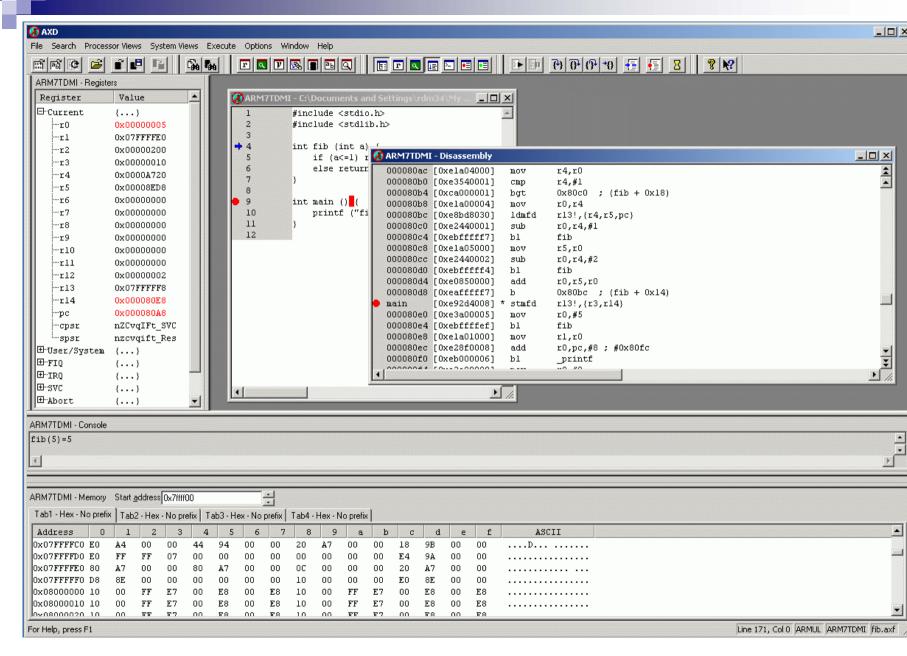
0		•	1
	AREA	hello, CODE, REAL	
SWI_WriteC SWI_Exit	EQU EQU ENTRY	0x18	; symbolic name for constant ; entry point into program
START LOOP	ADR LDRB CMP MOV SWINE ADD BNE	r2, [r1] r2, #0 r0, #SWI_WriteC 0x123456 r1, r1, #1	; ADD r1, pc, offset of TEXT ; load character ; check we are not at end of string ; call putchar (r1 points to char) ; do system call ; increase our pointer ; loop if we are not finished
	MOV MOV ADD SWI	r0, #SWI_Exit r1, #0x20000 r1,r1,#0x26 0x123456	; standard exit system call
TEXT	DCB END	"Hello World\n", 0	; end of source file!



The program in memory....

- 33					
	🚱 ARM7TDMI	- Disassembly			<u>=</u>
	0008000	[0xe28f1024]	add	rl,pc,#0x24 ; #0x802c	
	loop	[0xe5d12000]	ldrb	r2,[r1,#0]	
	00008008	[0xe3520000]	\mathtt{cmp}	r2,#0	
	0000800c	[0xe3a00003]	Mov	r0,#3	
	00008010	[0x1f123456]	swine	0x123456	
	00008014	[0xe2811001]	add	rl,rl,#l	
	00008018	[Oxlafffff9]	bne	loop	
	finish	[0xe3a00018]	MOA	r0,#0x18	
	00008020	[0xe3a01b80]	MOA	rl,#0x20000	
	00008024	[0xe2811026]	add	rl,rl,#0x26	
	→ 00008028	[0xef123456]	swi	0x123456	
	TEXT	[0x6c6c6548]	dcd	0x6c6c6548 Hell	
	00008030	[0x6f57206f]	dcd	0x6f57206f o Wo	
	00008034	[0x0a646c72]	dcd	0x0a646c72 rld.	
	00008038	rnvnnnnnnn1	ded	0	

```
int fib (int a) {
           if (a<=1) return a;
           else return fib (a-1)+fib(a-2);
fib:
                 r13!, {r4, r5, r14}
     stmfd
                                             ; preserve registers
                 r4, r0
                                             : r4=a
     mov
                 r0, #1
     cmp
                                             ; compare a with 1
                                             ; (if a<=1) result=a
     movle
              r0, r4
                                             ; (if a<=1) return
     Idmlefd
                 r13!, {r4, r5, pc}
     sub
                 r0, r4, #1
                                             ; fib (a-1)
     bl
                  fib
                 r5, r0
                                             ; r5=fib(a-1)
     mov
     sub
                 r0, r4, #2
                                             ; fib (a-2)
     bl
                  fib
                 r0, r5, r0
                                             ; result=fib(a-1)+fib(a-2)
     add
                 r13!, {r4, r5, pc}
     Idmfd
                                             ; return
```



R13 = 0x07FFFFF8 (stack pointer)

R4 = A720, R5 = 8ED8

R14 = 80E8 (address after BL fib instruction in main)

STMFD r13!, {r4, r5, r14} FF / 07 nn 00 7 00 FF nn. nn. 00 00 OxO7FFFFDO EO nn. 20 A7 0x07FFFFE0 🔑 00 00 0x07FFFFF¶ D8 00 00 8E 80 00 00 10 00 00 00 E7)x08000000 10 10 00 E8 00 PF 000080E8 00008ED8 0000A720SP at start Adding data to stack

```
Fib(5)
      R5
      R4
      Return Address to Main()
             R4=5
             Fib(R4-1)
                   R5
                   R4 (a=5)
                    Return address
                          Fib(3).... etc.
             Fib(R4-2)
             Return sum
```

ECAD Workshop Four

Sieve of Eratosthenes

```
2 3 4 5 6 7 8 9 10 11 12 13 14 15
0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
2 3 4 5 6 7 8 9 10 11 12 13 14 15
0 0 1 0 1 0 1 0 1 0 1 0 1 0
2 3 4 5 6 7 8 9 10 11 12 13 14 15
0 0 1 0 1 0 1 1 1 10 1 0 1 1
```

Later Workshops

- Make sure you use procedure call standard correctly
 - Use branch-and-link (BL) instruction to call procedure
 - □ Pass parameters in correct registers
 - □ Save registers when necessary
 - □ Restore registers on exit

Next Lecture

- OS Support and Memory Management
 - □ Virtual Memory
 - □ Interrupts/Exceptions
 - □ ARM specifics (operating modes, page table organisation etc.)

Self Study/Supervision Work

- Write a simple C program to sum the numbers 1 to n, compile it and examine the assembler produced
 - □ You might want to try executing the program in the debugging environment, this will allow you to single step through the program. You can look at the contents of the registers and memory as you go.

Challenge

- □ Write a program (in ARM assembly language) to reverse the bytes of a 32-bit register
- It's possible to do this using only one additional register to hold temporary results and 4 ARM instructions!