

## Probability and Computation: Problem Sheet 5

You are encouraged to submit your solutions at student reception or by emailing them to [luca.zanetti@cl.cam.ac.uk](mailto:luca.zanetti@cl.cam.ac.uk) by 2pm Friday 28th of February

### Question 1.

- (i) Prove that for every  $n \geq 2$  there is an unweighted, undirected  $n$ -vertex graph with conductance 1.
- (ii) (Open-Ended Bonus Question): Can you characterise all graphs with that property?

**Question 2.** Recall the 8-vertex graph from Lecture 11, slide 11, which has  $1 - \lambda_2(P) \approx 0.13$ . To get some idea of how small (or large) this value is, prove the following bounds on the conductance of any unweighted, connected, 3-regular graph with 8 vertices:

- (i) Show that for any such graph  $G$ ,  $\phi(G) \leq \frac{1}{2}$
- (ii) Show that for any such graph  $G$ ,  $\phi(G) \geq 1/12$
- (iii) Which lower and upper bounds on  $1 - \lambda_2$ , where  $\lambda_2$  is associated with the transition matrix of the lazy walk on  $G$ , can you deduce from (i) and (ii) using Cheeger's inequality?

**Question 3.** Find the conductance of the following graphs

1. The  $n$ -vertex path.
2. The 2-dimensional  $n \times m$  grid.
3. The complete binary tree of height  $h$ .

**Question 4.** Prove the following: A finite irreducible, aperiodic Markov chain with transition matrix  $P$  is reversible if and only if its transition probabilities satisfy

$$p_{j_1 j_2} p_{j_2 j_3} \cdots p_{j_{n-1} j_n} p_{j_n j_1} = p_{j_1 j_n} p_{j_n j_{n-1}} \cdots p_{j_3 j_2} p_{j_2 j_1},$$

for any sequence of states  $j_1, \dots, j_n$ .

**Question 5.** Let  $G$  be a connected graph and  $P$  be the transition matrix of the simple random walk on  $G$ .

1. Show that if  $-1$  is an eigenvalue of  $P$  then the walk is periodic.
2. Show that if  $G$  is bipartite and  $\mu$  is an eigenvalue of  $P$  then  $-\mu$  is also an eigenvalue of  $G$ , and that  $\mu$  and  $-\mu$  have the same multiplicity.

**Question 6.** Recall the definition of the  $\ell_1$ -mixing time  $\tau$  and the  $\ell_2$ -mixing time  $\tau_2$ . Prove that  $\tau_2(\epsilon) \geq \tau(2\epsilon)$  for any  $\epsilon \in (0, 1/2]$ .

## Hints

**Hint** (Question 2(i)). *Use a randomised algorithm similar to the one for MAX-CUT to find a subset with conductance at most  $1/2$ .*

**Hint** (Question 4). *Recall the Convergence Theorem for finite Markov chains.*