

Theorem 118 For $R \subseteq A \times A$, let

$$\mathcal{F}_R = \{ Q \subseteq A \times A \mid R \subseteq Q \wedge Q \text{ is a preorder} \}.$$

Then, (i) $R^{\circ*} \in \mathcal{F}_R$ and (ii) $R^{\circ*} \subseteq \bigcap \mathcal{F}_R$. Hence, $R^{\circ*} = \bigcap \mathcal{F}_R$.

PROOF:

$$(i) \quad R^{\circ*} \subseteq A \times A$$

$$R \subseteq R^{\circ*} = \bigcup_{n \in \mathbb{N}} R^{\circ n}$$

$$R^{\circ 1} = R$$

$R^{\circ*}$ is a preorder

- reflexivity: $\text{id}_A \subseteq R^{\circ*}$

$$R^{\circ 0} = \text{id}_A$$

- transitivity: $R^{\circ*} \circ R^{\circ*} \subseteq R^{\circ*}$

$$\forall x, y, z. x R^{\circ*} y \wedge y R^{\circ*} z \Rightarrow x R^{\circ*} z.$$

$$R^{on} \subseteq \bigcap F_R$$

$$\begin{aligned} & X \subseteq \bigcap \{A_i \mid i \in I\} \\ \text{iff} \quad & \forall i \in I, X \subseteq A_i \end{aligned}$$

$$\Leftrightarrow R^{on} \subseteq Q \quad \forall Q \subseteq A \times A, R \subseteq Q \wedge Q \text{ a preorder.}$$

$$\begin{aligned} & \parallel \\ & \bigcup_{n \in \mathbb{N}} R^{on} \end{aligned}$$

$$\begin{aligned} & \bigcup \{A_i \mid i \in I\} \subseteq X \\ \text{iff} \quad & \forall i \in I, A_i \subseteq X \end{aligned}$$

$$\Leftrightarrow \left(\begin{aligned} & \forall Q \subseteq A \times A, R \subseteq Q \wedge Q \text{ preorder} \\ & \Rightarrow \forall n \in \mathbb{N}, R^{on} \subseteq Q \end{aligned} \right)$$

Let $Q \subseteq A \times A$ s.t. $R \subseteq Q$ & Q is a preorder

We show $R^n \subseteq Q$ for all $n \in \mathbb{N}$ by induction.

Base case: $R^0 \subseteq Q$ ✓
 \parallel
 id_A

Ind. Step: Assume $R^n \subseteq Q$ for $n \in \mathbb{N}$. (IH)

RTP: $R^{n+1} \subseteq Q$ ✓
 \parallel
 $R \circ R^n$

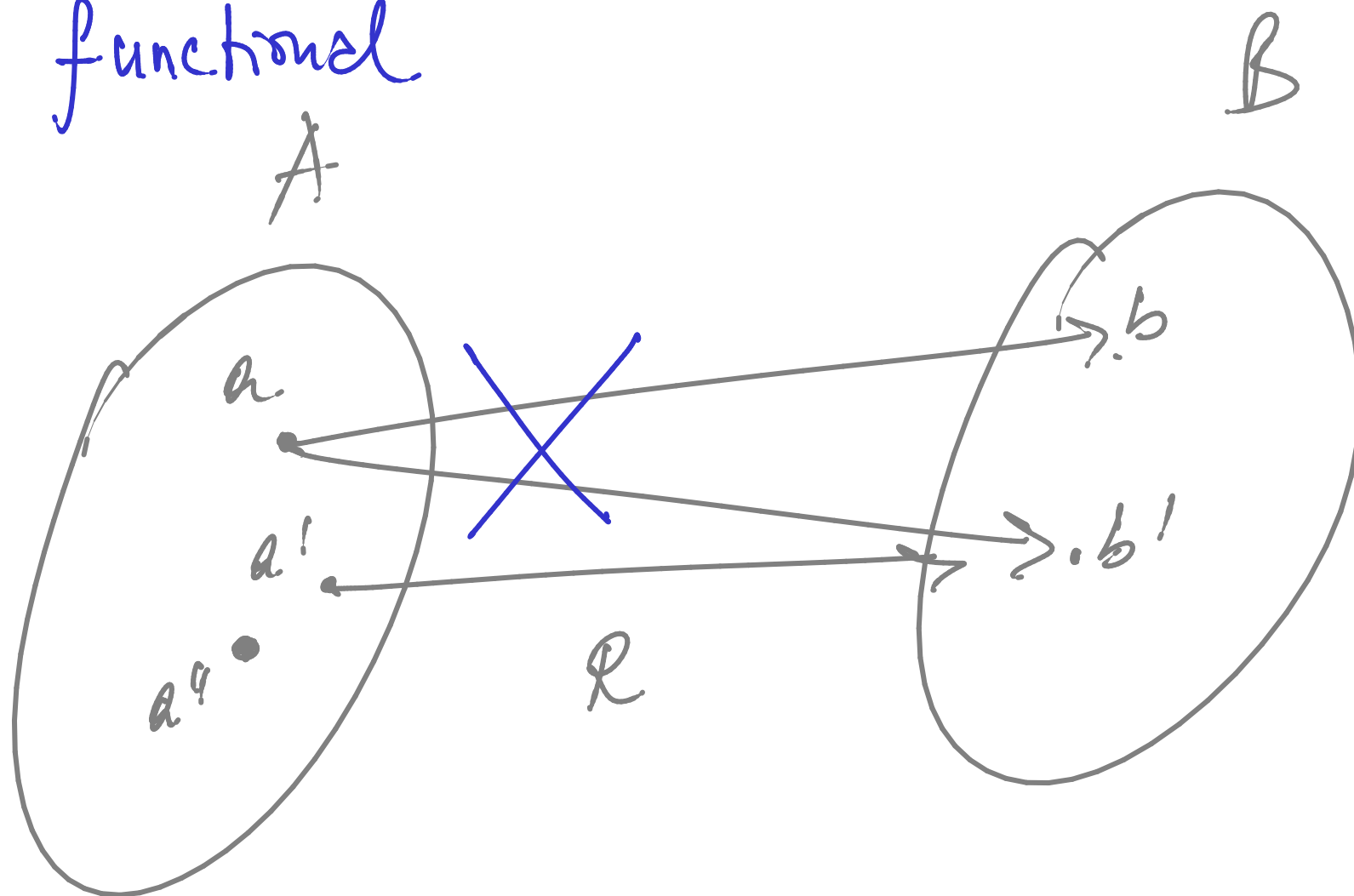
Let $x (R \circ R^n) y \Leftrightarrow \exists z. x R z \wedge z R^n y$

$\Rightarrow \exists z. x Q z \wedge z Q y \Rightarrow x Q y$ \square

$$R: A \rightarrow B$$

functional

non-deterministic
input-output behaviours.



Partial functions

Definition 119 A relation $R : A \rightarrow B$ is said to be functional, and called a partial function, whenever it is such that

$$\forall a \in A. \forall b_1, b_2 \in B. a R b_1 \wedge a R b_2 \implies b_1 = b_2 .$$

If $f: A \rightarrow B$ is a partial function
Then $\forall a \in A. \exists a f b \in B$ Then such a b is unique.

NOTATION

$$\left| \begin{array}{l} f(a) \downarrow \text{ iff } \exists b. a f b \\ f(a) \uparrow \text{ iff } \neg (f(a) \downarrow) \end{array} \right.$$

NOTATION

we typically denote it
by $f(a)$

EXAMPLE: id_A is a partial function.

NOTATION

$$f: A \rightarrow B$$

partial functions.

$$g: B \rightarrow C$$

So

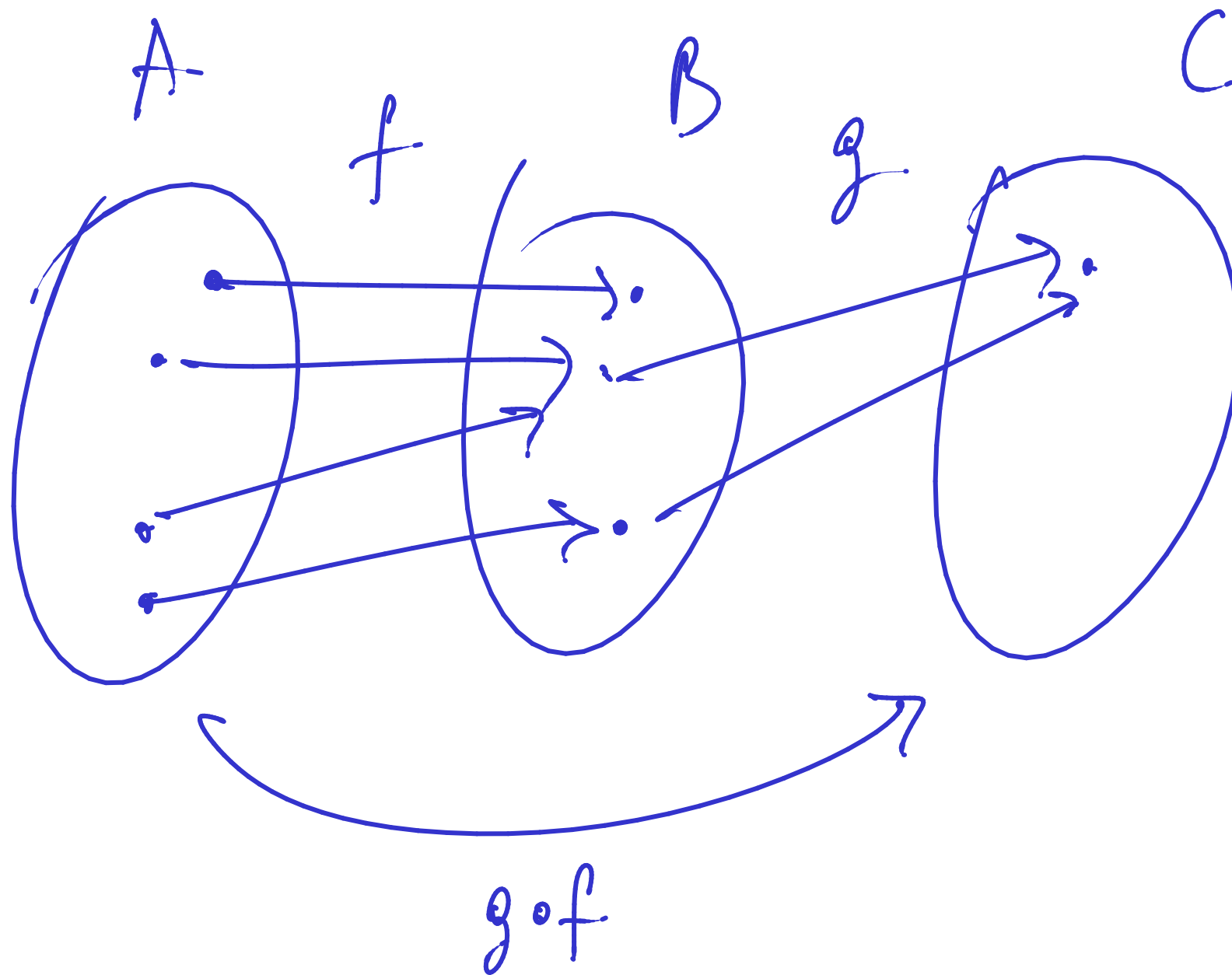
$$g \circ f: A \rightarrow C$$

is it functional?

YES!

$$a \in A$$

$$(g \circ f)(a) = \begin{cases} \uparrow & \text{if } f(a) \uparrow \\ \uparrow & \text{if } f(a) \downarrow \text{ but } g(f(a)) \uparrow \\ g(f(a)) & \text{if } f(a) \downarrow \text{ and } g(f(a)) \downarrow \end{cases}$$



Theorem 121 *The identity relation is a partial function, and the composition of partial functions yields a partial function.*

NB

$$f = g : A \multimap B$$

iff

$$\forall a \in A. (f(a) \downarrow \iff g(a) \downarrow) \wedge f(a) = g(a)$$

Convention: Typically one defines partial functions $f: A \multimap B$ by selecting $D \subseteq A$ (domain of definition) and giving a rule/mapping that associates $a \in D$ to $f(a) \in B$

$(a \xrightarrow{f} f(a))$

Recall $\mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N} \times \mathbb{N} : (n, m) \mapsto (\text{quo}(n, m), \text{rem}(n, m))$

Example: The following defines a partial function $\mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z} \times \mathbb{N}$:

- ▶ for $n \geq 0$ and $m > 0$,
 $(n, m) \mapsto (\text{quo}(n, m), \text{rem}(n, m))$
- ▶ for $n \geq 0$ and $m < 0$,
 $(n, m) \mapsto (-\text{quo}(n, -m), \text{rem}(n, -m))$
- ▶ for $n < 0$ and $m > 0$,
 $(n, m) \mapsto (-\text{quo}(-n, m) - 1, \text{rem}(m - \text{rem}(-n, m), m))$
- ▶ for $n < 0$ and $m < 0$,
 $(n, m) \mapsto (\text{quo}(-n, -m) + 1, \text{rem}(-m - \text{rem}(-n, -m), -m))$

Its domain of definition is $\{ (n, m) \in \mathbb{Z} \times \mathbb{Z} \mid m \neq 0 \}$.

$(A \Rightarrow B) =$ the set of all partial functions from A to B

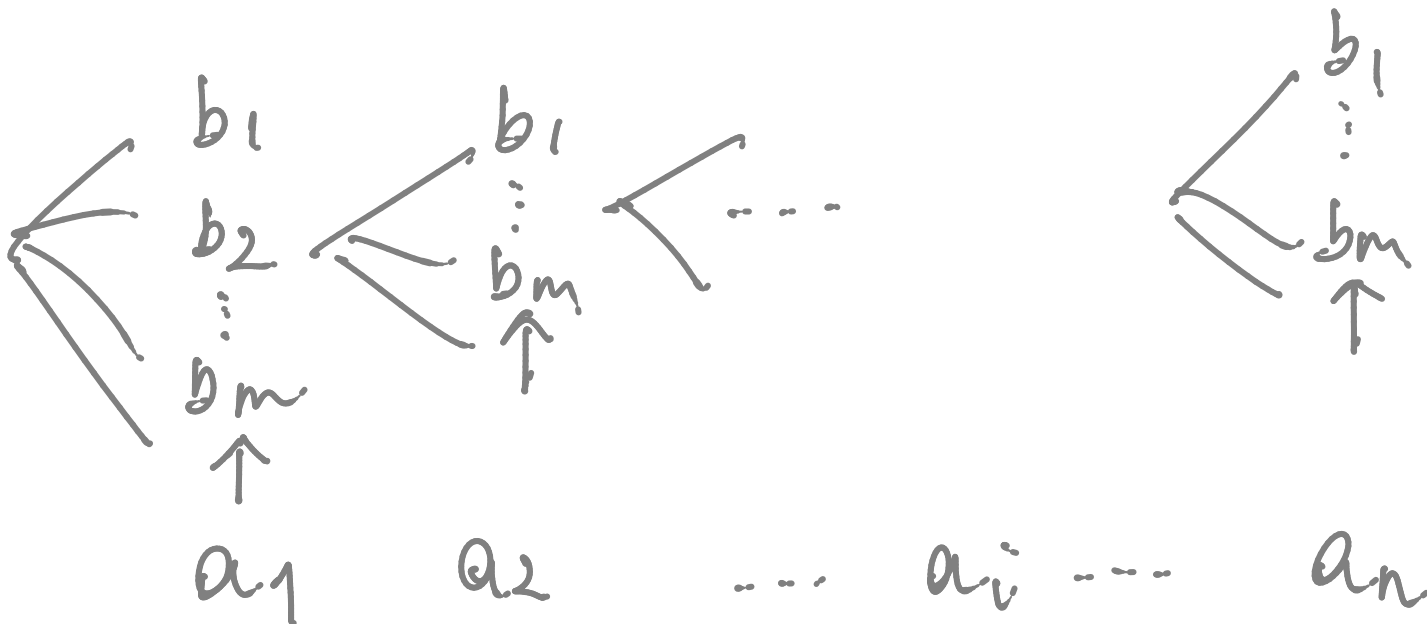
Proposition 122 For all finite sets A and B ,

$$\#(A \Rightarrow B) = (\#B + 1)^{\#A} = (m+1)^n$$

PROOF IDEA:

$$A = \{a_1, \dots, a_n\} \quad \#A = n$$

$$B = \{b_1, \dots, b_m\} \quad \#B = m$$



Functions (or maps)

Definition 123 A partial function is said to be total, and referred to as a (total) function or map, whenever its domain of definition coincides with its source.

iff $f: A \rightarrow B$ is total

$$\text{dom}(f) = \{a \in A \mid f(a) \downarrow\} = A$$

NOTATION:

For these we write $f: A \rightarrow B$.

$(A \Rightarrow B)$ = set of all functions from A to B .

Theorem 124 For all $f \in \text{Rel}(A, B)$, $\left\{ \begin{array}{l} \text{the unique such } b \text{ is} \\ \text{denoted } f(a). \end{array} \right.$

$$f \in (A \Rightarrow B) \iff \forall a \in A. \exists! b \in B. a f b.$$

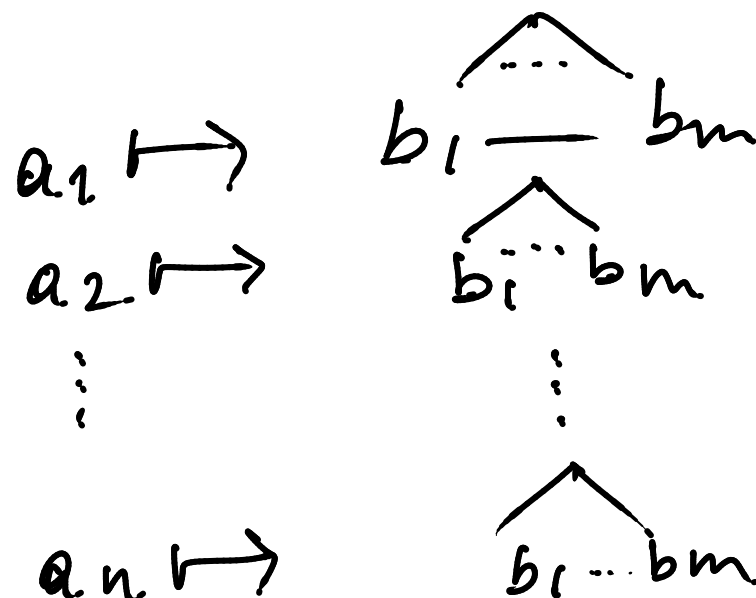
Proposition 125 For all finite sets A and B ,

$$\#(A \Rightarrow B) = \#B^{\#A} = m^n$$

PROOF IDEA:

$$A = \{a_1, \dots, a_n\}$$

$$B = \{b_1, \dots, b_m\}$$



$$\begin{aligned} & \frac{\#B}{\#(A \Rightarrow B)} \\ &= \#(A \Rightarrow (B \cup [1])) \end{aligned}$$

$$[n] = \{0, \dots, n-1\}$$

$$[1] = \{0\}$$

Theorem 126 *The identity partial function is a function, and the composition of functions yields a function.*

NB

1. $f = g : A \rightarrow B$ iff $\forall a \in A. f(a) = g(a)$.
2. For all sets A , the identity function $\text{id}_A : A \rightarrow A$ is given by the rule

$$\text{id}_A(a) = a$$

and, for all functions $f : A \rightarrow B$ and $g : B \rightarrow C$, the composition function $g \circ f : A \rightarrow C$ is given by the rule

$$(g \circ f)(a) = g(f(a)) \quad .$$