

## The division theorem and algorithm

**Theorem 43 (Division Theorem)** For every natural number  $m$  and positive natural number  $n$ , there exists a unique pair of integers  $q$  and  $r$  such that  $q \geq 0$ ,  $0 \leq r < n$ , and  $m = q \cdot n + r$ .

$\forall \text{ nat. } m, \text{ pos. nat. } n.$

$(\exists \text{ int } q, r. q \geq 0, 0 \leq r < n, m = q \cdot n + r)$  Existence.

$\wedge \left[ \begin{array}{l} \forall q, q', r, r'. q, q' \geq 0, 0 \leq r, r' < n, \\ m = q \cdot n + r \wedge m = q' \cdot n + r' \\ \Rightarrow (q = q' \wedge r = r') \end{array} \right]$  Uniqueness.

Uniqueness

Assume nat.  $m$ , pos. nat.  $n$ .

Assume  $q, q', r, r' : q, q' \geq 0, \underline{0 \leq r, r' < n}$ , (3)  
 $\underline{m = q \cdot n + r}$  (1)  $\wedge$   $\underline{m = q' \cdot n + r'}$  (2)

RTP :  $q = q' \wedge r = r'$

From (1) and (2)  $\underline{(q - q') \cdot n = r' - r}$  (4)

Case (i) : Suppose  $r' \geq r$  . so that  $r' - r$  is a natural...

that by (3) is  $< n$  . and by (4) is a multiple of  $n$  .

Therefore  $q - q' = 0$  and hence  $q = q'$ .

$$q \cdot n + r = m = q' \cdot n + r'$$

$$\text{As } q = q'$$

By cancellation,  $r = r'$ .

Case  $r \geq r'$ : Consider  $(q' - q) \cdot n = r - r'$  and proceed analogously.

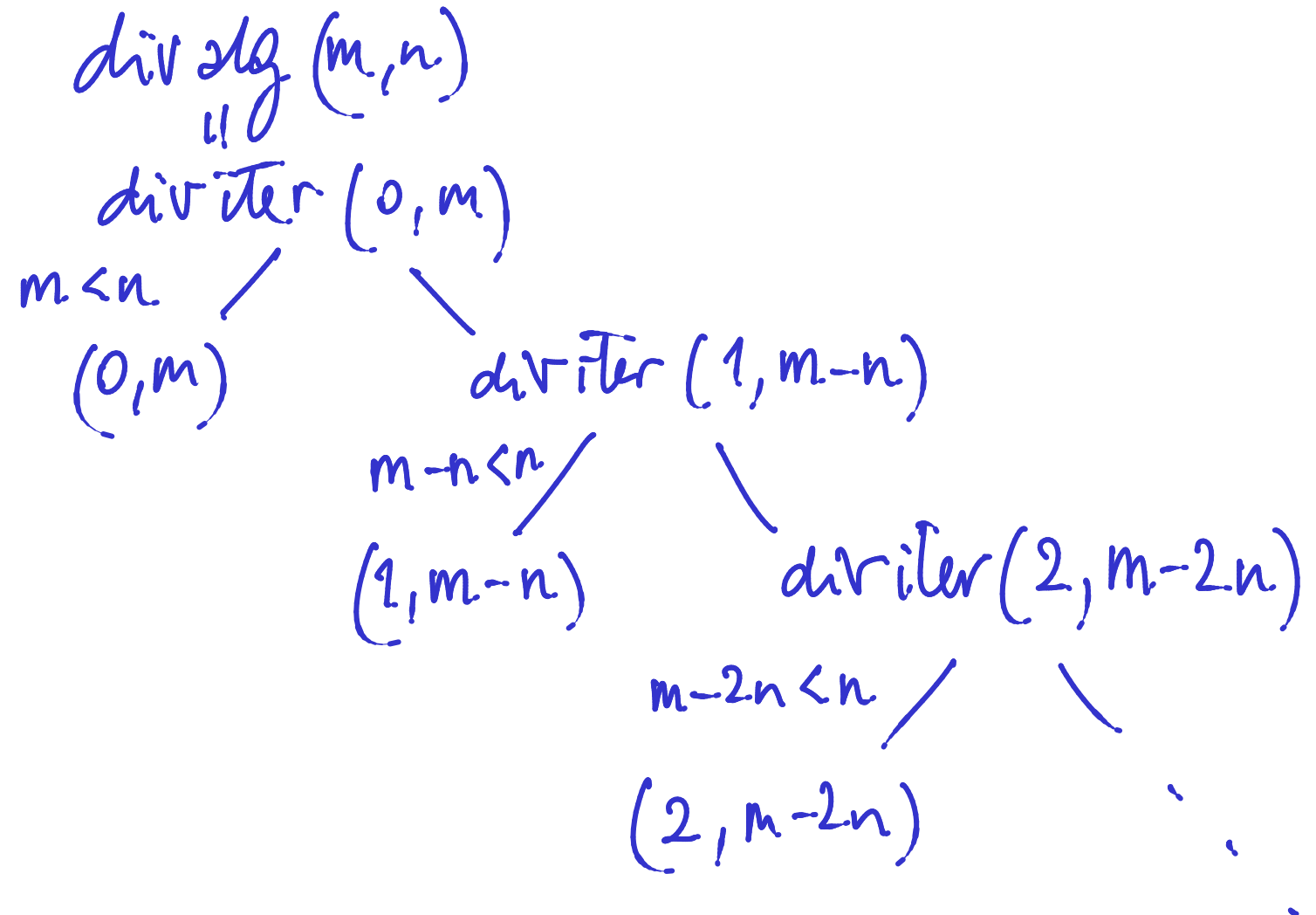


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**Theorem 43 (Division Theorem)** *For every natural number  $m$  and positive natural number  $n$ , there exists a unique pair of integers  $q$  and  $r$  such that  $q \geq 0$ ,  $0 \leq r < n$ , and  $m = q \cdot n + r$ .*

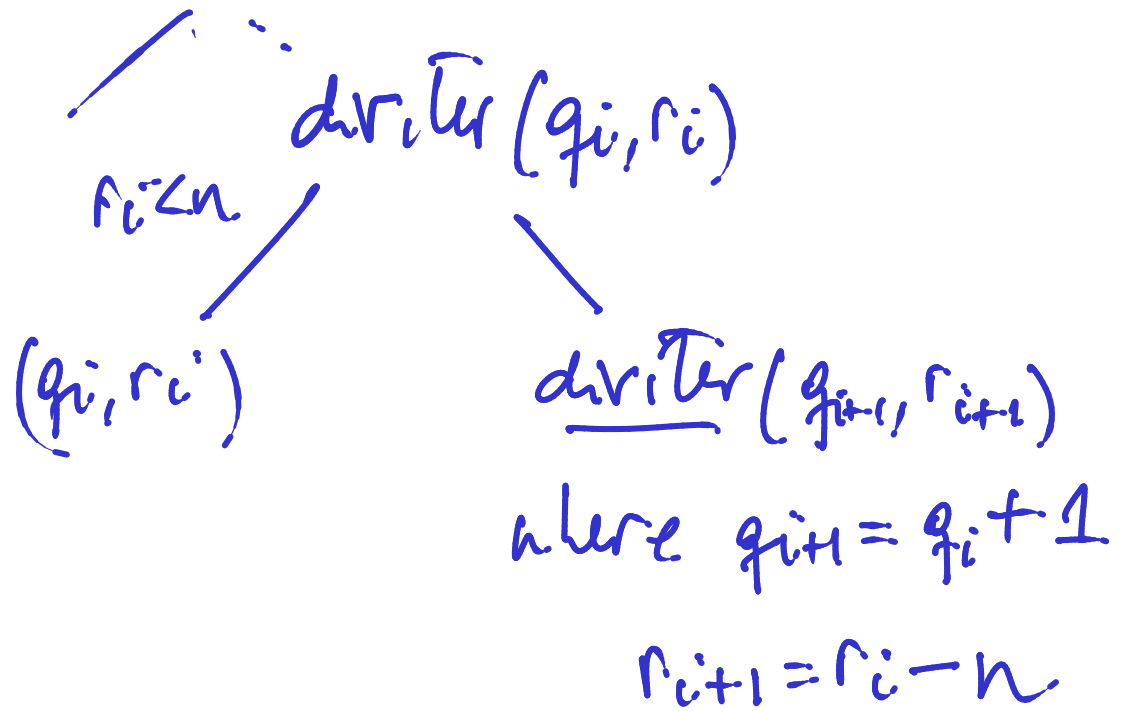
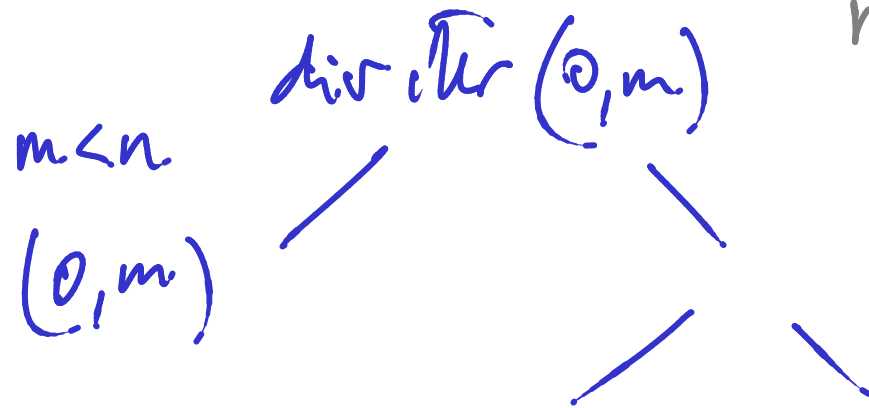
**Definition 44** *The natural numbers  $q$  and  $r$  associated to a given pair of a natural number  $m$  and a positive integer  $n$  determined by the Division Theorem are respectively denoted  $\text{quo}(m, n)$  and  $\text{rem}(m, n)$ .*

PROOF OF Theorem 43:



# Partial Correctness

$$m = 0 \cdot n + m \quad \checkmark$$



Assume  
 $m = q_i \cdot n + r_i$

RTP:  
 $?$   
 $m = q_{i+1} \cdot n + r_{i+1} \quad \checkmark$

$$(q_{i+1}) \cdot n + (r_i - n)$$

||

$$q_i \cdot n + n + r_i - n = q_i \cdot n + r_i = m$$

If the computation terminates in step  $k$  then  
 $m = q_k \cdot n + r_k$  and  $r_k < n$

Total correctness: We need show the algorithm terminates. At each call of divider( $q, r$ ) either  $r < n$  or we are done or we have a call divider( $q+1, r-n$ ) with smaller, but positive, second argument.



The Division Algorithm in ML:

```
fun divalg( m , n )
```

```
  = let
```

```
    fun diviter( q , r )
```

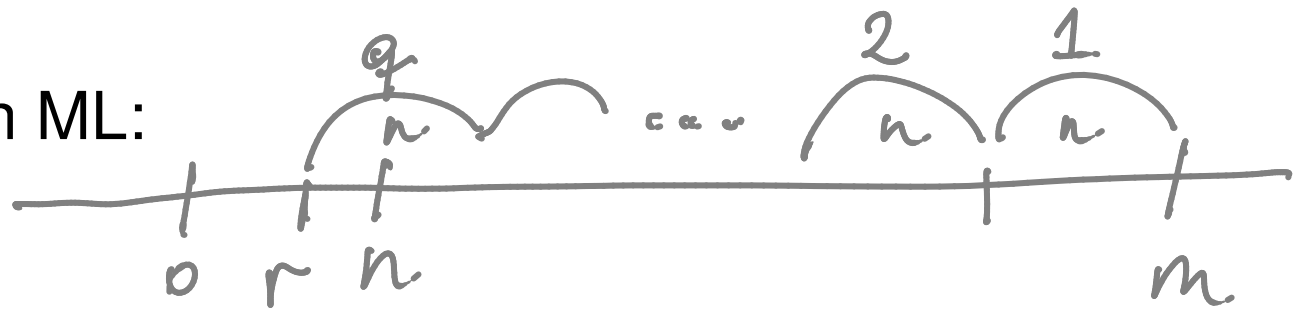
```
      = if r < n then ( q , r )
```

```
        else diviter( q+1 , r-n )
```

```
  in
```

```
    diviter( 0 , m )
```

```
  end
```



$$\boxed{\begin{array}{l} 0 \leq q, \quad 0 \leq r < n \\ m = q \cdot n + r \end{array}}$$

```
fun quo( m , n ) = #1( divalg( m , n ) )
```

```
fun rem( m , n ) = #2( divalg( m , n ) )
```



**Theorem 45** *For every natural number  $m$  and positive natural number  $n$ , the evaluation of  $\text{divalg}(m, n)$  terminates, outputting a pair of natural numbers  $(q_0, r_0)$  such that  $r_0 < n$  and  $m = q_0 \cdot n + r_0$ .*

PROOF:

$$k = \underline{\text{quo}}(k, m) \cdot m + \underline{\text{rem}}(k, m)$$

**Proposition 46** Let  $m$  be a positive integer. For all natural numbers  $k$  and  $l$ ,

$$l = \underline{\text{quo}}(l, m) \cdot m + \underline{\text{rem}}(l, m)$$

$$k \equiv l \pmod{m} \iff \underline{\text{rem}}(k, m) = \underline{\text{rem}}(l, m) .$$

PROOF: Let  $m$  be a positive integer. Let  $k, l$  be nat.

$(\Leftarrow)$  Easy.

$(\Rightarrow)$  Assume  $k - l = i \cdot m$  for some  $i \in \mathbb{Z}$ .

$$\begin{array}{c} \parallel \\ (\underline{\text{quo}}(k, m) - \underline{\text{quo}}(l, m)) \cdot m + (\underline{\text{rem}}(k, m) - \underline{\text{rem}}(l, m)) \end{array}$$

Then  $\underline{\text{rem}}(k, m) - \underline{\text{rem}}(l, m) = 0$

So  $\underline{\text{rem}}(k, m) = \underline{\text{rem}}(l, m)$  .



**Corollary 47** Let  $m$  be a positive integer.

1. For every natural number  $n$ ,

$$n \equiv \text{rem}(n, m) \pmod{m} .$$

Since

$$n = \text{quo}(n, m) \cdot m + \text{rem}(n, m)$$

we have  $n - \text{rem}(n, m)$  is a multiple of  $m$   $\square$

PROOF:

**Corollary 47** Let  $m$  be a positive integer.

1. For every natural number  $n$ ,

$$n \equiv \text{rem}(n, m) \pmod{m} .$$

2. For every integer  $k$  there exists a unique integer  $[k]_m$  such that

$$0 \leq [k]_m < m \text{ and } k \equiv [k]_m \pmod{m} .$$

PROOF:

