

Probability and Computation: Problem sheet 1

Question 1. Recall that a permutation $\sigma : [n] \rightarrow [n]$ is a bijection from $[n]$ to $[n]$. A cycle $c = (c_1, c_2, \dots, c_k)$ is a sequence such that $\sigma(a_i) = a_{i+1 \bmod k}$. Example: if $\sigma(1) = 3, \sigma(2) = 2, \sigma(3) = 4$ and $\sigma(4) = 1$ (this is also denoted $\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 3 & 2 & 4 & 1 \end{pmatrix}$) then $(1, 3, 4)$ is a cycle (as are $(3, 4, 1)$ and $(4, 1, 3)$).

- (i) How many permutations on $[n]$ are there?
- (ii) Show that the number of permutations on $[n]$ which contain a cycle of length ℓ is $n!/\ell$.
- (iii) Let \mathcal{E} be the event that a uniformly distributed random permutation contains no cycle longer than $\lfloor n/2 \rfloor$. Show that $\mathbf{P}[\mathcal{E}] = 1 - \left(\frac{1}{\lfloor n/2 \rfloor + 1} + \dots + \frac{1}{n} \right)$.
- (iv) Approximate the expression for $\mathbf{P}[\mathcal{E}]$ above by integrals to give $\mathbf{P}[\mathcal{E}] = 1 - \ln 2 + o(1) \approx 0.31183$.

Question 2. Recall that a probability vector (distribution) is a non-negative real vector whose elements sum to 1. A stochastic matrix is a real square matrix, where each row is a probability vector. Observe every Stochastic matrix gives rise to a Markov chain and visa versa.

- (i) Let $\nu \in \mathbb{R}_+^n$ be a probability vector and $M \in \mathbb{R}_+^{n \times n}$ be a stochastic matrix. Show that νM is a probability vector.

A doubly stochastic matrix is a real square matrix, where each row and column is a probability vector.

- (ii) Prove that the uniform distribution is stationary for any Markov chain whose transition matrix is doubly stochastic.

Question 3. In this question we shall consider two probability measures μ_1, μ_2 on the common state space $[12] = \{1, \dots, 12\}$. The first measure is the uniform measure $\mu_1(x) = 1/12$ for all $x \in [12]$, this is a “12 sided fair die”. The second μ_2 is the measure generated by the sum of two independent, fair, 6-sided die i.e. $\mu_2(1) = 0, \mu_2(2) = 1/6^2, \dots$. Calculate $\|\mu_1 - \mu_2\|$.

Question 4. Recall Jerry’s Coupling Z_t for two Simple Random walks T_t, J_t on the 3-cycle. This coupling is given by Markov chain $Z_t = (T_t, J_t)$ on $\{0, 1, 2\} \times \{0, 1, 2\}$:

- Run the Cat T_t as normal
- Mouse J_t moves according to the rule: $J_{t+1} = \begin{cases} J_t + 1 \bmod 3 & \text{if } T_{t+1} = T_t + 1 \bmod 3 \\ J_t - 1 \bmod 3 & \text{if } T_{t+1} = T_t - 1 \bmod 3. \end{cases}$

This coupling lets the “Mouse” J_t avoid the “Cat” T_t forever.

1. Write down the rule for “Tom’s Coupling” - the coupling where T_t and J_t meet as soon as possible and show this is a valid coupling.
2. Draw a diagram of the Markov Chain Z_t on $\{0, 1, 2\} \times \{0, 1, 2\}$ generated by Tom’s Coupling.

Question 5. In lecture 1 we saw the balls in bins experiment where one assigns m balls to n bins uniformly and independently. By considering the number of balls assigned to bins between the first time i bins are empty and $i - 1$ bins are empty, show that the the expected number of balls one must assign before there is no empty bin is $n \log n + \Theta(n)$.

Question 6. Let X_n be the sum of n independent rolls of a fair die. Show that, for any $k \geq 2$,

$$\lim_{n \rightarrow \infty} \mathbf{P}[X_n \text{ is divisible by } k] = \frac{1}{k}.$$

Hint. At face value X_n is an (infinite) Markov chain on \mathbb{N} . We would like to consider it as a finite Markov chain, reduction $\text{mod } m$ (for some suitable m) will help us achieve this.

Question 7. State j is accessible from state i if, for some integer $n \geq 0$, $P_{i,j}^n > 0$. If two states i and j are accessible from each other, we say that they communicate and we write $i \sim j$. Prove that communicating relation \sim defines an equivalence relation.

Question 8. Prove the following Lemma from class: For any probability distributions μ and η on a countable state space Ω

$$\|\mu - \eta\|_{tv} = \frac{1}{2} \sum_{\omega \in \Omega} |\mu(\omega) - \eta(\omega)|.$$

Hint. Recall the sets Ω^\pm from the Coupling Lemma.

Question 9. This question asks you to prove lower bounds on the mixing time of some lazy random walks on graphs.

1. Let $G = (V_1 \cup V_2)$ be a graph made of two disjoint complete graphs of n vertices, supported respectively on V_1 and V_2 , connected by a single edge. This is called the Barbell graph. Consider a lazy random walk on G . Prove that $t_{mix}(G) = \Omega(n^2)$ (recall from Lecture 3 that $t_{mix} = \tau(1/4)$).
2. Suppose now we add $s < n$ edges to the Barbell graph, where each edge has one endpoint in V_1 and the other endpoint in V_2 . What happens to $t_{mix}(G)$?
3. Consider now a version of the Barbell graph where $|V_1| = n$, $|V_2| = \lfloor \log(n) \rfloor$ and there exists only an edge between V_1 and V_2 . What is the mixing time of this graph?