# CST 2018/19 Part II Denotational Semantics Exercise Sheet

## Lectures 1–3

**Exercise 1** Let  $\Omega$  de the domain of "vertical natural numbers" pictured in Figure 1 of the lecture notes.

- (i) Is every monotone function from  $\Omega$  to  $\Omega$  continuous?
- (ii) Does every monotone function from  $\Omega$  to  $\Omega$  have a least prefixed point?

**Exercise 2** For a partially ordered set  $(P, \sqsubseteq)$ , let  $(Ch(P), \sqsubseteq_{ptw})$  be the partially ordered set of chains in *P* ordered pointwise. That is

$$Ch(P) \stackrel{\text{def}}{=} \{x = \{x_n\}_{n \in \mathbb{N}} \mid \text{for all } i \le j \text{ in } \mathbb{N}, x_i \sqsubseteq x_j \text{ in } P\}$$

and

$$x \sqsubseteq_{ptw} x' \stackrel{\text{der}}{\Leftrightarrow} x_n \sqsubseteq x'_n$$
 for all  $n \in \mathbb{N}$ 

Show that if *P* is a domain then so is Ch(P).

#### **Exercise 3**

- (i) Let  $D_1$ ,  $D_2$  and E be domains. Show that a function  $f : D_1 \times D_2 \to E$  is continuous if it is continuous is each argument separately, that is, if for all  $d_1 \in D_1$  and  $d_2 \in D_2$ , the functions  $f(d_1, -) : D_2 \to E$  and  $f(-, d_2) : D_1 \to E$  are continuous.
- (ii) Let  $\mathbb{O}$  be the domain with two elements  $\bot \sqsubseteq \top$ . For a domain *E* and  $e \in E$ , define the function  $g_e : E \to \mathbb{O}$  by

$$g_e(x) = \begin{cases} \bot & \text{if } x \sqsubseteq e \\ \top & \text{if } x \not\sqsubseteq e \end{cases}$$

Show that  $g_e$  is continuous.

- (iii) As an example of the definition in part (ii), let  $E = \mathbb{B}_{\perp} \times \mathbb{B}_{\perp}$ , where  $\mathbb{B} = \{true, false\}$ , and consider  $g_{(false, false)} : E \to \mathbb{O}$ . Show that  $g_{(false, false)}(x, y) = \top$  iff x = true or y = true.
- (iv) Let  $f : D \to E$  be a function between domains *D* and *E*. Show that *f* is continuous iff  $\forall e \in E. g_e \circ f$  is continuous.

**Exercise 4** Let O be the domain in Exercise 3(ii).

- (i) Draw a diagram which represents the elements of the function domain  $O \rightarrow O$  and shows their ordering.
- (ii) Any set X can be considered as a flat domain X<sub>⊥</sub> by adding a bottom element. Show that the strict continuous functions X<sub>⊥</sub> → O are in 1-1 correspondence with the subsets of X.

### Lectures 4–6

**Exercise 5** Let *D* be a domain and  $k : D \to D$  a continuous function. Let  $\mathbb{B} = \{true, false\}$ . Define the conditional function  $if : \mathbb{B}_{\perp} \times D \times D \to D$  by

$$if(b, d, d') = \begin{cases} d & \text{if } b = true \\ d' & \text{if } b = false \\ \bot & \text{if } b = \bot \end{cases}$$

Let  $h : D \to \mathbb{B}_{\perp}$  be a continuous function which is strict (so  $h(\perp) = \perp$ ). The function  $f^*$  is the least continuous function from D to D such that

$$\forall x \in D. f^*(x) = if(h(x), x, f^*(k(x)))$$

- (i) Show that  $\forall x \in D$ .  $h(f^*(x)) = if(h(x), h(x), h(f^*(k(x))))$ .
- (ii) Prove that the property  $Q(f) \stackrel{\text{def}}{\Leftrightarrow} \forall x \in D$ .  $h(f(x)) \sqsubseteq true$  is admissible.
- (iii) Prove  $Q(f^*)$  by fixed point induction.

**Exercise 6** Suppose that *D* is a domain and  $f : D \times D \rightarrow D$  is a continuous function satisfying the property  $\forall d, e \in D$ . f(d, e) = f(e, d). Let  $g : D \times D \rightarrow D \times D$  be defined by

$$g(d_1, d_2) = (f(d_1, f(d_1, d_2)), f(f(d_1, d_2), d_2))$$

Let  $(u_1, u_2) = fix(g)$ . Show that  $u_1 = u_2$  using Scott induction.

**Exercise 7** Let *D* and *E* be domains and let  $f : D \to D$  and  $g : E \to E$  be continuous functions.

- (i) Define  $f \times g : D \times E \to D \times E$  to be the continuous function given by  $(f \times g)(d, e) = (f(d), g(e))$  and let  $\pi_1 : D \times E \to D$  and  $\pi_2 : D \times E \to E$  respectively denote the first and second projection functions. Show that  $fix(f \times g) \sqsubseteq (fix(f), fix(g))$  and that  $fix(f) \sqsubseteq \pi_1(fix(f \times g))$  and  $fix(g) \sqsubseteq \pi_2(fix(f \times g))$ .
- (ii) It follows from part (i) that  $fix(f \times g) = (fix(f), fix(g))$ . Use this and Scott's Fixed Point Induction Principle to show that, for all strict continuous functions  $h : D \to E$ ,

$$h \circ f = g \circ h \Rightarrow h(fix(f)) = fix(g)$$

#### Lectures 7–10

**Exercise 8** Which of the following statements are true or false, and why?

- (i) For all PCF-types  $\tau$  and closed PCF-terms  $M_1$  and  $M_2$  of type  $\tau$ , if  $\llbracket M_1 \rrbracket = \llbracket M_2 \rrbracket$  in  $\llbracket \tau \rrbracket$  then  $M_1 \cong_{ctx} M_2 : \tau$ .
- (ii) For all PCF-types  $\tau$  and closed PCF-terms  $M_1$  and  $M_2$  of type  $\tau$ , if  $M_1 \cong_{ctx} M_2 : \tau$ , then  $[M_1] = [M_2]$  in  $[\tau]$ .
- (iii) For all closed PCF-terms  $M_1$  and  $M_2$  of type  $nat \rightarrow nat$ , if  $M_1 \cong_{ctx} M_2 : nat \rightarrow nat$ , then  $\llbracket M_1 \rrbracket = \llbracket M_2 \rrbracket$  in  $\mathbb{N}_{\perp} \rightarrow \mathbb{N}_{\perp}$ .

**Exercise 9** Consider the following two statements for PCF terms  $M_1$  and  $M_2$  for which the typings  $\Gamma \vdash M_1 : \tau$  and  $\Gamma \vdash M_2 : \tau$  hold for some type environment  $\Gamma$  and type  $\tau$ .

(1) For all PCF contexts C[-] for which  $C[M_1]$  : *bool* and  $C[M_2]$  : *bool*,

 $\mathcal{C}[M_1] \Downarrow_{bool} \Leftrightarrow \mathcal{C}[M_2] \Downarrow_{bool}$ 

where, for a closed term *M* of type  $\tau$ , the notation  $M \Downarrow_{\tau}$  stands for the existence of a value *V* :  $\tau$  for which  $M \Downarrow_{\tau} V$ .

(2) For all PCF contexts C[-] for which  $C[M_1]$  : *bool* and  $C[M_2]$  : *bool*,

 $\mathcal{C}[M_1] \Downarrow_{bool} \operatorname{true} \Leftrightarrow \mathcal{C}[M_2] \Downarrow_{bool} \operatorname{true}$ 

- (i) Show that (1) implies (2).
- (ii) Show that (2) implies that  $M_1$  and  $M_2$  are contextually equivalent.

### Extra

**Exercise 10** Suppose that *D* is a domain and that  $lam : (D \rightarrow D) \rightarrow D$  and  $app : D \rightarrow (D \rightarrow D)$  are continuous functions. Use this data to give a denotational semantics for the terms of the untyped  $\lambda$ -calculus, by answering question 5 from paper 7 of the 1998 CS Tripos.